Computation That Counts

Counting Problems

- Counting problems are concerned with the number of solutions.

  - #SAT: the number of satisfying truth assignments to a boolean formula.

  - #HAMILTONIAN PATH: the number of Hamiltonian paths in a graph.

- They cannot be easier than their decision versions.

  - The decision problem has a solution if and only if the solution count is larger than 0.

- But they can be harder than their decision versions.

Decision and Counting Problems

- FP is the set of polynomial-time computable functions $f : \{0, 1\}^* \to \mathbb{Z}$.

  - GCD, LCM, matrix-matrix multiplication, etc.

- If $\#\text{SAT} \in \text{FP}$, then $\text{P} = \text{NP}$.

  - Given boolean formula ϕ, calculate its number of satisfying truth assignments, k, in polynomial time.

  - Declare "$\phi \in \text{SAT}$" if and only if $k \geq 1$.

- The validity of the reverse direction is open.

A Counting Problem Harder than Its Decision Version

- CYCLE asks if a directed graph contains a cycle.

- #CYCLE counts the number of cycles in a directed graph.

- CYCLE is in P by a simple greedy algorithm.

- But #CYCLE is hard unless P = NP.

Counting Class #P

A function f is in #P (or $f \in$ #P) if

- There exists a polynomial-time NTM M.

- $M(x)$ has $f(x)$ accepting paths for all inputs x.

- $f(x) =$ number of accepting paths of $M(x)$.

Some #P Problems

- $f(\phi)$ = number of satisfying truth assignments to ϕ.

  - The desired NTM guesses a truth assignment T and accepts ϕ if and only if $T \models \phi$.

  - Hence $f \in \#\mathrm{P}$.

  - f is also called #SAT.

- #HAMILTONIAN PATH.

- #3-COLORING.

#P Completeness

- Function f is #P-complete if

  - $f \in$ #P.

  - #P $\subseteq$ FP^f.

    * Every function in #P can be computed in polynomial time with access to a black box[a] or **oracle** for f.

  - Of course, oracle f will be accessed only a polynomial number of times.

  - #P is said to be **polynomial-time Turing-reducible to** f.

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[a]Think of it as a subroutine

#SAT Is #P-Complete[a]

- First, it is in #P (p. 735).

- Let $f \in$ #P compute the number of accepting paths of M.

- Cook's theorem uses a *parsimonious* reduction from M on input x to an instance ϕ of SAT (p. 280).

  - Hence the number of accepting paths of $M(x)$ equals the number of satisfying truth assignments to ϕ.

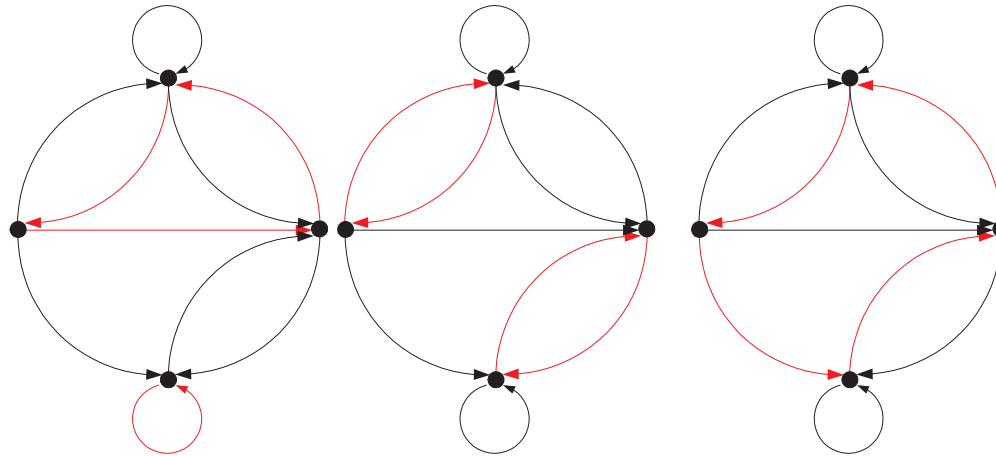- Call the oracle #SAT with ϕ to obtain the desired answer regarding $f(x)$.

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[a]Valiant (1979); in fact, #2SAT is also #P-complete.

Leslie G Valiant (1949–)

CYCLE COVER

- A set of node-disjoint cycles that cover all nodes in a directed graph is called a **cycle cover**.


- There are 3 cycle covers (in red) above.

CYCLE COVER and BIPARTITE PERFECT MATCHING

Proposition 92 CYCLE COVER *and* BIPARTITE PERFECT MATCHING *(p. 440) are parsimoniously reducible to each other.*

- A polynomial-time algorithm creates a bipartite graph G' from any directed graph G.

- Moreover, the number cycle covers for G equals the number of bipartite perfect matchings for G'.

- And vice versa.

Corollary 93 CYCLE COVER $\in P$.

Illustration of the Proof
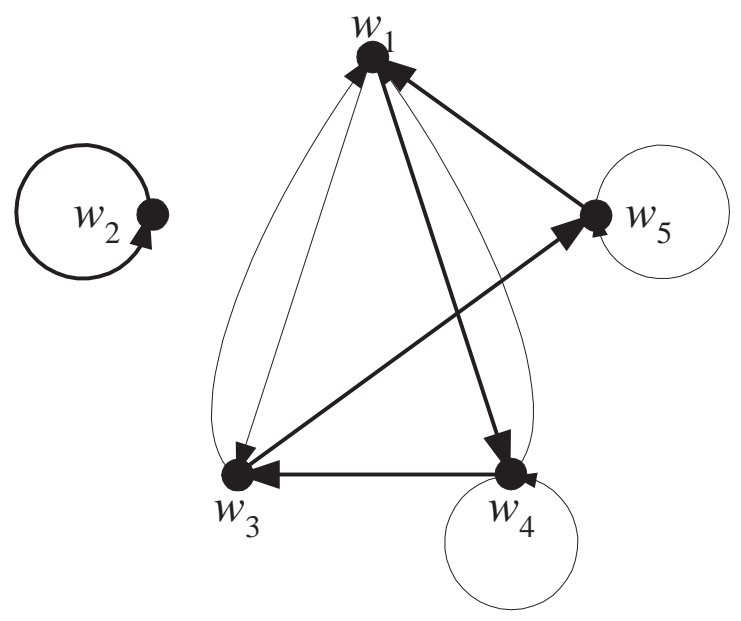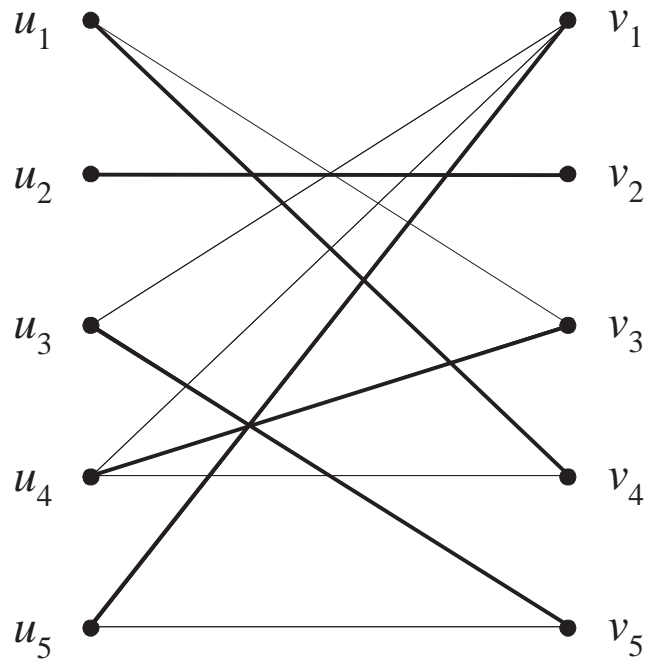
Permanent

- The **permanent** of an $n \times n$ integer matrix A is

$$\operatorname{perm}(A) = \sum_\pi \prod_{i=1}^n A_{i,\pi(i)}.$$

  - π ranges over all permutations of n elements.

- 0/1 PERMANENT computes the permanent of a 0/1 (binary) matrix.

  - The permanent of a binary matrix is at most $n!$.

- Simpler than determinant (5) on p. 443: no signs.

- Surprisingly, much harder to compute than determinant!

Permanent and Counting Perfect Matchings

- BIPARTITE PERFECT MATCHING is related to determinant (p. 444).

- #BIPARTITE PERFECT MATCHING is related to permanent.

Proposition 94 0/1 PERMANENT *and* BIPARTITE PERFECT MATCHING *are parsimoniously reducible to each other.*

The Proof

- Given a bipartite graph G, construct an $n \times n$ binary matrix A.

  - The (i, j)th entry A_{ij} is 1 if $(i, j) \in E$ and 0 otherwise.

- Then $\text{perm}(A) =$ number of perfect matchings in G.

Illustration of the Proof Based on p. 741 (Left)

$$A = \begin{bmatrix} 0 & 0 & 1 & \boxed{1} & 0 \\ 0 & \boxed{1} & 0 & 0 & 0 \\ 1 & 0 & 0 & 0 & \boxed{1} \\ 1 & 0 & \boxed{1} & 1 & 0 \\ \boxed{1} & 0 & 0 & 0 & 1 \end{bmatrix}.$$

- $\mathrm{perm}(A) = 4$.

- The permutation corresponding to the perfect matching on p. 741 is marked.

Permanent and Counting Cycle Covers

Proposition 95 $0/1$ PERMANENT *and* CYCLE COVER *are parsimoniously reducible to each other.*

- Let A be the adjacency matrix of the graph on p. 741 (right).

- Then $\text{perm}(A) =$ number of cycle covers.

Three Parsimoniously Equivalent Problems

We summarize Propositions 92 (p. 740) and 94 (p. 743) in the following.

Lemma 96 0/1 PERMANENT, BIPARTITE PERFECT MATCHING, *and* CYCLE COVER *are* **parsimoniously equivalent**.

We will show that the counting versions of all three problems are in fact #P-complete.

WEIGHTED CYCLE COVER

- Consider a directed graph G with integer weights on the edges.

- The weight of a cycle cover is the product of its edge weights.

- The **cycle count** of G is sum of the weights of all cycle covers.

  - Let A be G's adjacency matrix but $A_{ij} = w_i$ if the edge (i, j) has weight w_i.

  - Then $\text{perm}(A) = G$'s cycle count (same proof as Proposition 95 on p. 746).

- #CYCLE COVER is a special case: All weights are 1.

An Example[a]
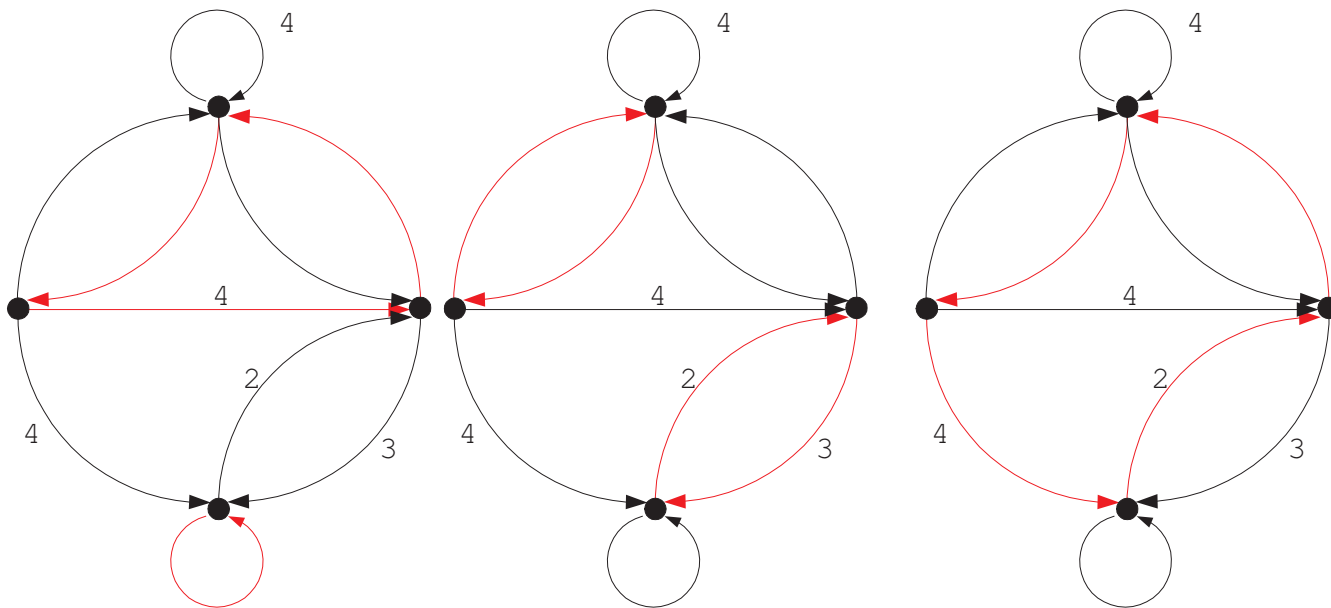
There are 3 cycle covers, and the cycle count is

$$(4 \cdot 1 \cdot 1) \cdot (1) + (1 \cdot 1) \cdot (2 \cdot 3) + (4 \cdot 2 \cdot 1 \cdot 1) = 18.$$

[a]Each edge has weight 1 unless stated otherwise.

Three #P-Complete Counting Problems

Theorem 97 (Valiant (1979)) 0/1 PERMANENT, #BIPARTITE PERFECT MATCHING, *and* #CYCLE COVER *are #P-complete.*

- By Lemma 96 (p. 747), it suffices to prove that #CYCLE COVER is #P-complete.

- #SAT is #P-complete (p. 737).

- #3SAT is #P-complete because it and #SAT are parsimoniously equivalent (p. 283).

- We shall prove that #3SAT is polynomial-time Turing-reducible to #CYCLE COVER.
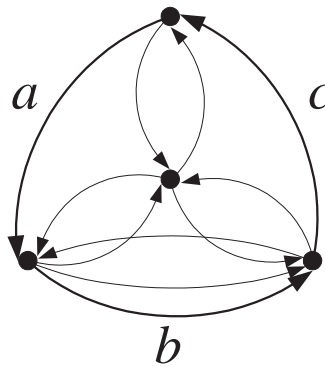
The Proof (continued)

- Let ϕ be the given 3SAT formula.

  - It contains n variables and m clauses (hence $3m$ literals).

  - It has $\#\phi$ satisfying truth assignments.

- First we construct a *weighted* directed graph H with cycle count

$$\#H = 4^{3m} \times \#\phi.$$

- Then we construct an unweighted directed graph G.

- We make sure $\#H$ (hence $\#\phi$) is polynomial-time Turing-reducible to G's number of cycle covers (denoted $\#G$).

The Proof: the Clause Gadget (continued)

- Each clause is associated with a **clause gadget**.

- Each edge has weight 1 unless stated otherwise.

- Each bold edge corresponds to one literal in the clause.

- There are not *parallel* lines as bold edges are schematic only (preview p. 765).

The Proof: the Clause Gadget (continued)

- Following a bold edge means making the literal false (0).

- A cycle cover cannot select *all* 3 bold edges.

  – The interior node would be missing.

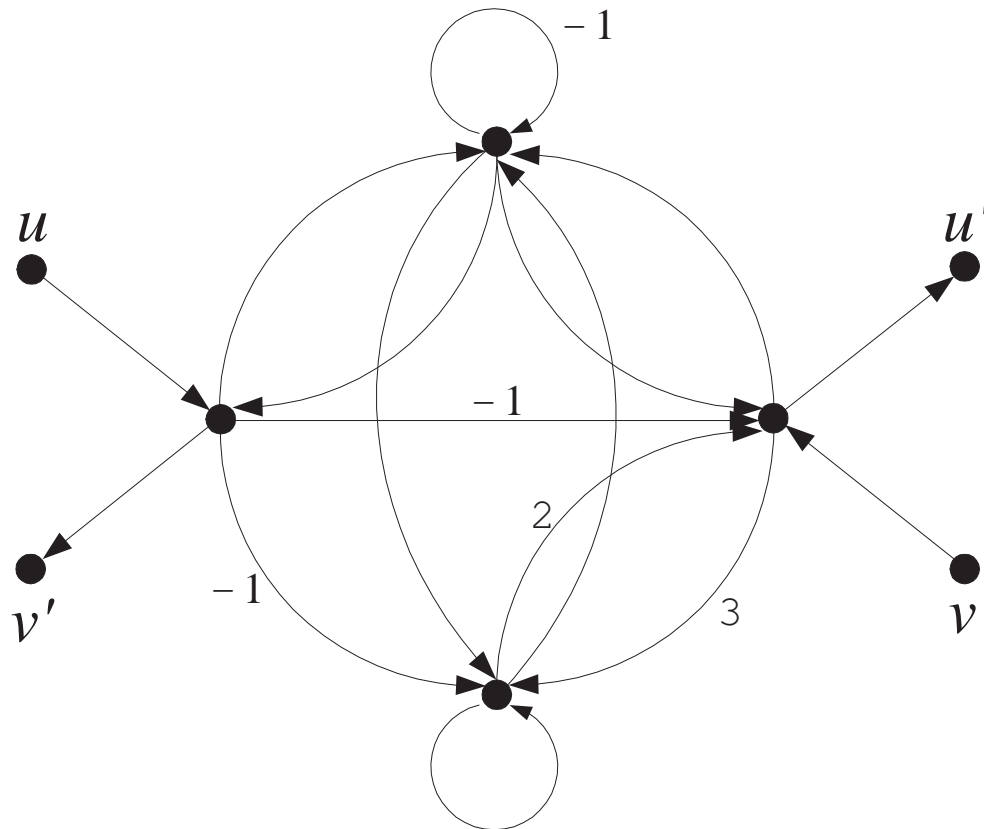- Every proper nonempty subset of bold edges corresponds to a unique cycle cover of weight 1 (see next page).

The Proof: the Clause Gadget (continued)

7 possible cycle covers, one for each satisfying assignment:
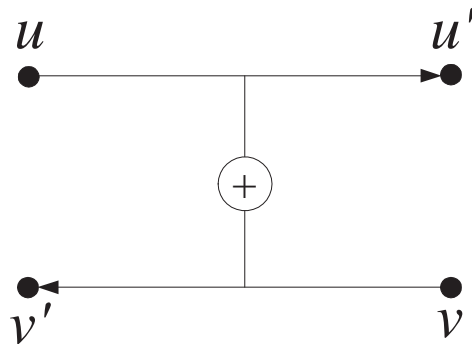(1) $a = 0, b = 0, c = 1$, (2) $a = 0, b = 1, c = 0$, etc.

(1)   (2)   (3)   (4)   (5)   (6)   (7)
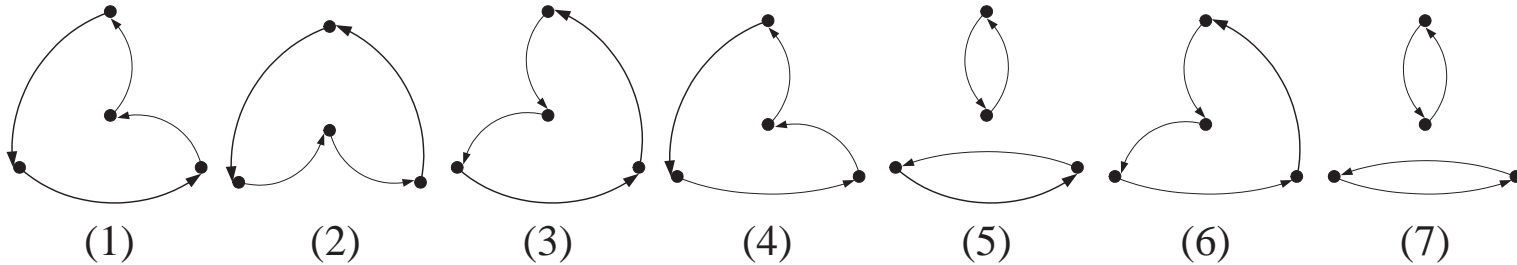
The Proof: the XOR Gadget (continued)

The Proof: Properties of the XOR Gadget (continued)

- The XOR gadget schema:

$$
\begin{array}{ccc}
u & & u' \\
\bullet \longrightarrow & \!\!\!\!\!\!\! \vert \!\!\!\!\!\!\! & \longrightarrow \!\!\!\bullet \\
& \oplus & \\
\bullet \longleftarrow & \!\!\!\!\!\!\! \vert \!\!\!\!\!\!\! & \longleftarrow \!\!\!\bullet \\
v' & & v
\end{array}
$$

- *At most one* of the 2 schematic edges will be included in a cycle cover.

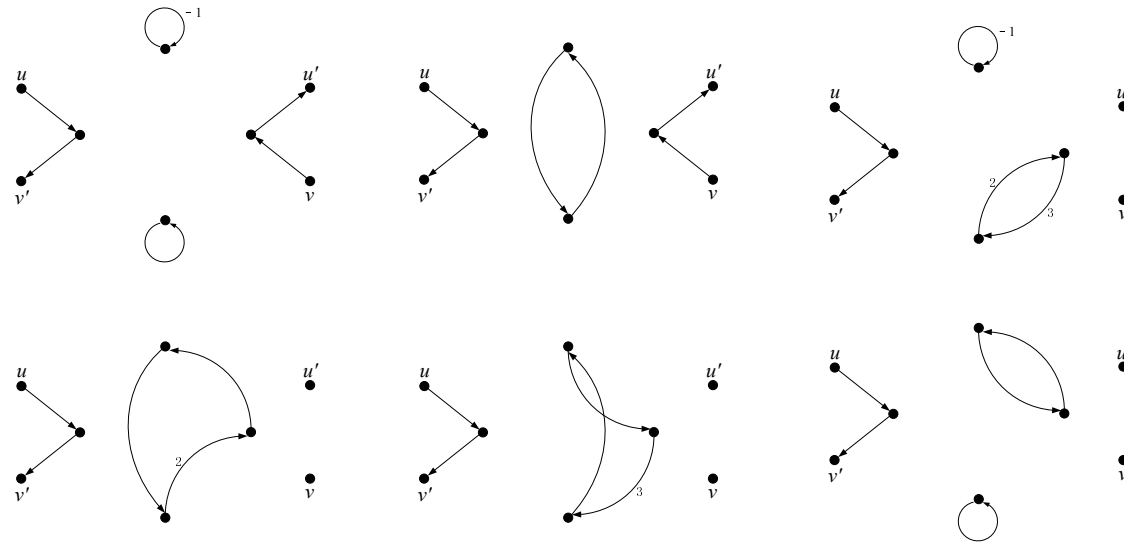- There will be $3m$ XOR gadgets, one for each literal.

The Proof: Properties of the XOR Gadget (continued)

Total weight of $-1 - 2 + 6 - 3 = 0$ for cycle covers not
entering or leaving it.
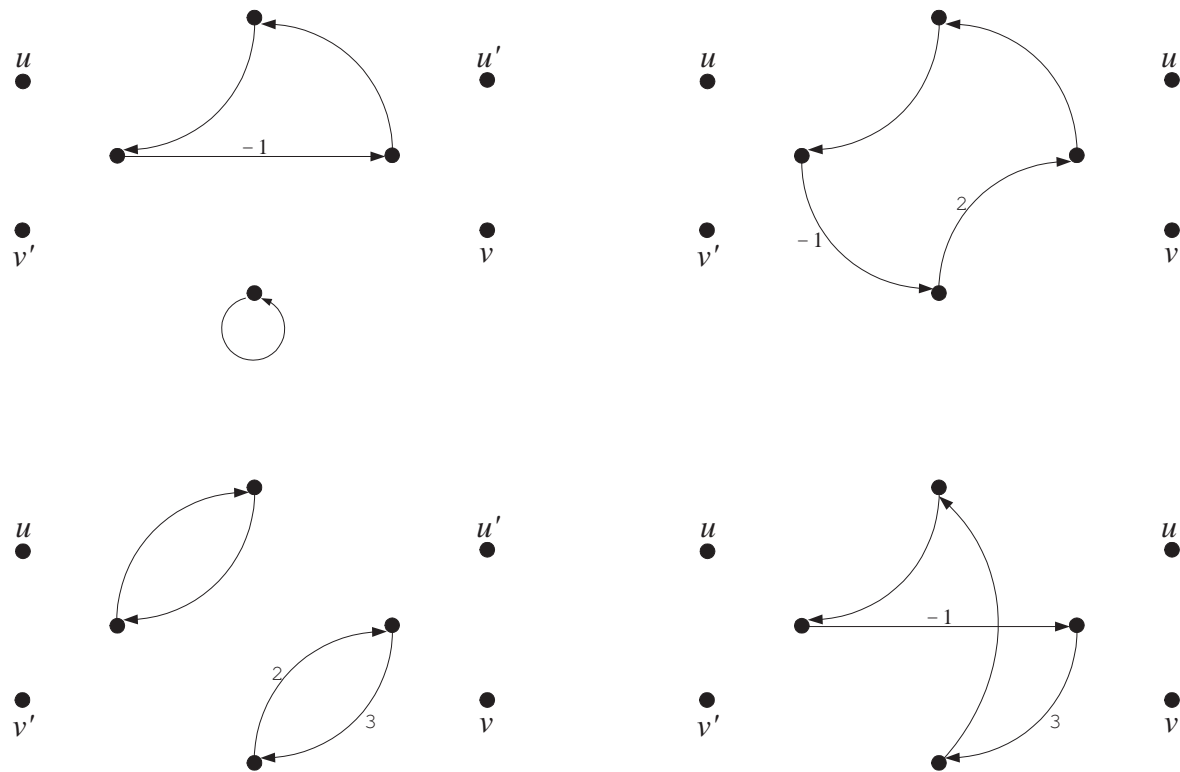
The Proof: Properties of the XOR Gadget (continued)

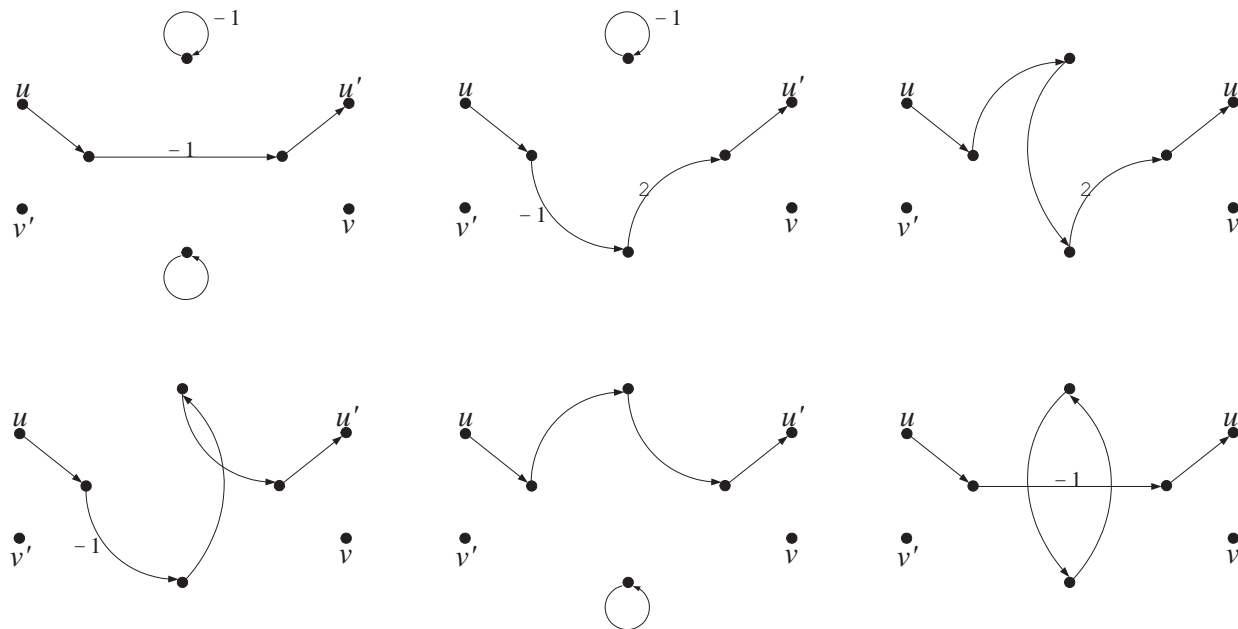- Total weight of $-1 + 1 - 6 + 2 + 3 + 1 = 0$ for cycle covers entering at u and leaving at v'.[a]

- Same for cycle covers entering at v and leaving at u'.

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[a]Corrected by Mr. Yu-Tsung Dai (**B91201046**) and Mr. Che-Wei Chang (**R95922093**) on December 27, 2006.

The Proof: Properties of the XOR Gadget (continued)

- Total weight of $1 + 2 + 2 - 1 + 1 - 1 = 4$ for cycle covers entering at u and leaving at u'.

- Same for cycle covers entering at v and leaving at v'.

The Proof: Summary (continued)

- Cycle covers not entering *all* of the XOR gadgets contribute 0 to the cycle count.

    – Let x denote an XOR gadget not entered for a cycle cover c.

    – Now, the said cycle covers' total contribution is

$$
\begin{aligned}
&= \sum_{\text{cycle cover } c \text{ for } H} \text{weight}(c) \\
&= \sum_{\text{cycle cover } c \text{ for } H-x} \text{weight}(c) \sum_{\text{cycle cover } c \text{ for } x} \text{weight}(x) \\
&= \sum_{\text{cycle cover } c \text{ for } H-x} \text{weight}(c) \cdot 0 \\
&= 0.
\end{aligned}
$$

The Proof: Summary (continued)

- Cycle covers entering *any* of the XOR gadgets and leaving illegally contribute 0 to the cycle count.

- For every XOR gadget entered and exited legally, the total weight of a cycle cover is multiplied by 4.
  - With an XOR gadget x entered and exited legally fixed,

  $$\text{contributions of such cycle covers to the cycle count}$$

  $$\sum_{\text{cycle cover } c \text{ for } H} \text{weight}(c)$$

  $$= \sum_{\text{cycle cover } c \text{ for } H - x} \text{weight}(c) \sum_{\text{cycle cover } c \text{ for } x} \text{weight}(x)$$

  $$= \sum_{\text{cycle cover } c \text{ for } H - x} \text{weight}(c) \cdot 4.$$

The Proof: Summary (continued)

- Hereafter we consider only cycle covers which enter every XOR gadget and leaves it legally.

  - Only these cycle covers contribute nonzero weights to the cycle count.
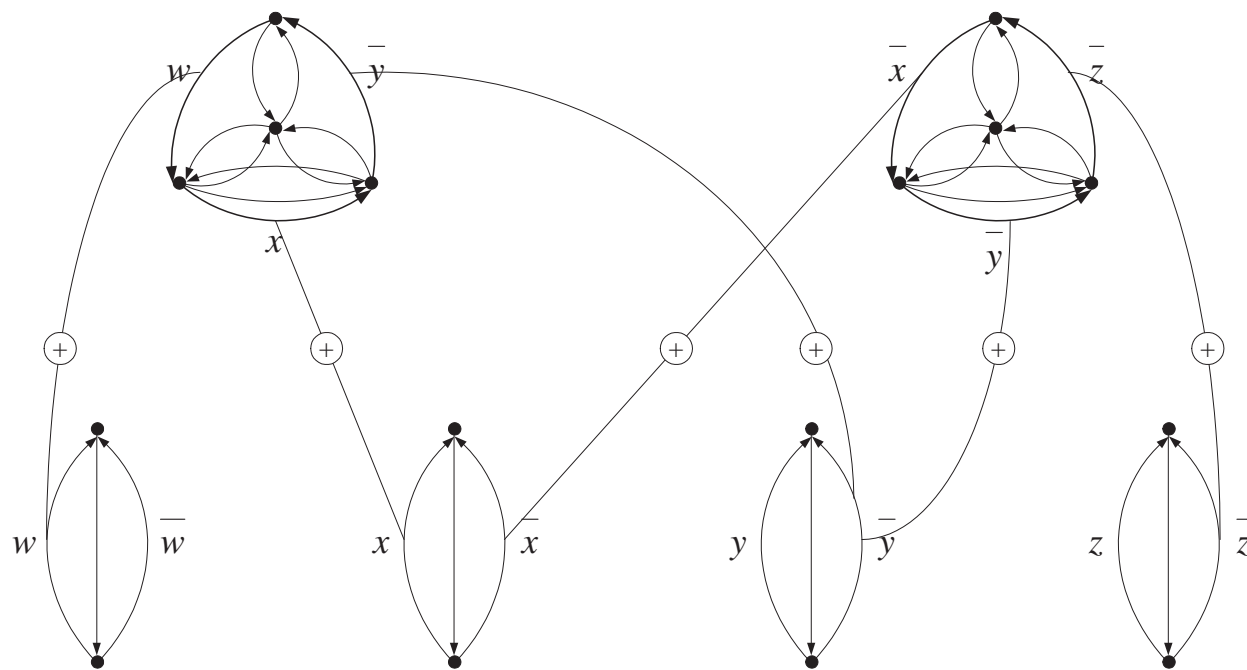
- They are said to **respect** the XOR gadgets.

The Proof: the Choice Gadget (continued)

- One choice gadget (a schema) for each variable.
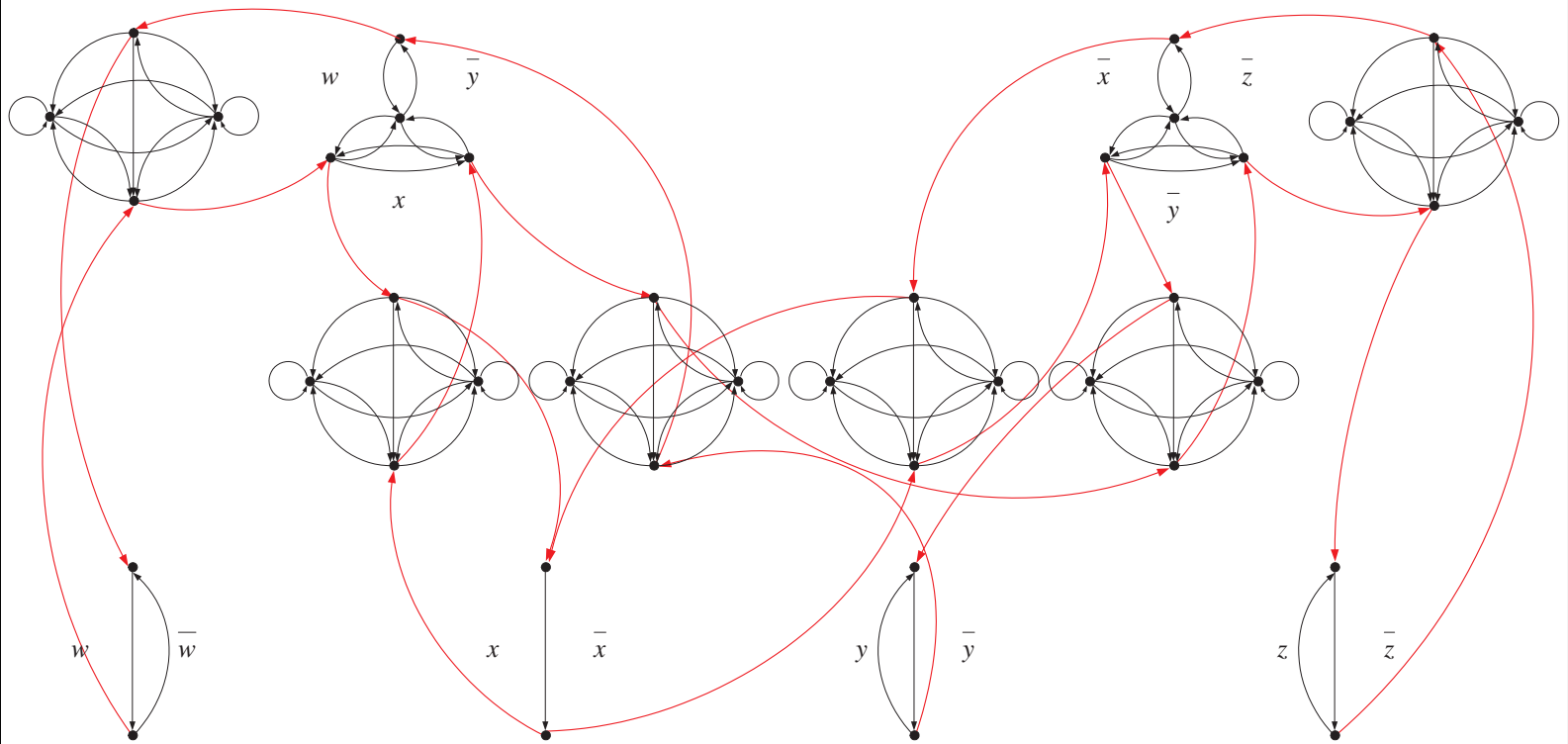
$$x \qquad \overline{x}$$
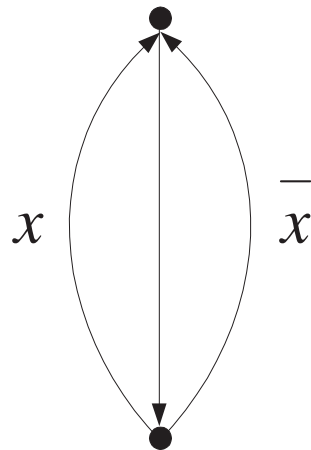
- It gives the truth assignment for the variable.

- Use it with the XOR gadget to enforce consistency.

Schema for $(w \vee x \vee \bar{y}) \wedge (\bar{x} \vee \bar{y} \vee \bar{z})$
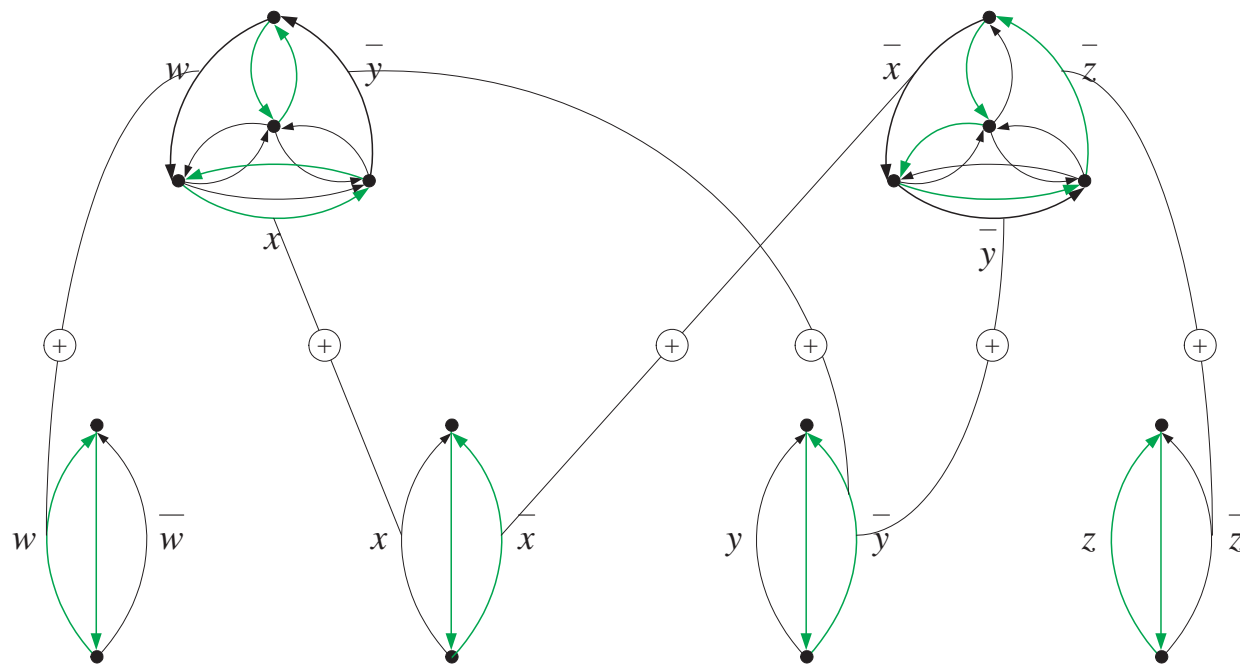
Full Graph $(w \vee x \vee \bar{y}) \wedge (\bar{x} \vee \bar{y} \vee \bar{z})$

The Proof: a Key Observation (continued)

Each satisfying truth assignment to ϕ corresponds to a schematic cycle cover that respects the XOR gadgets.

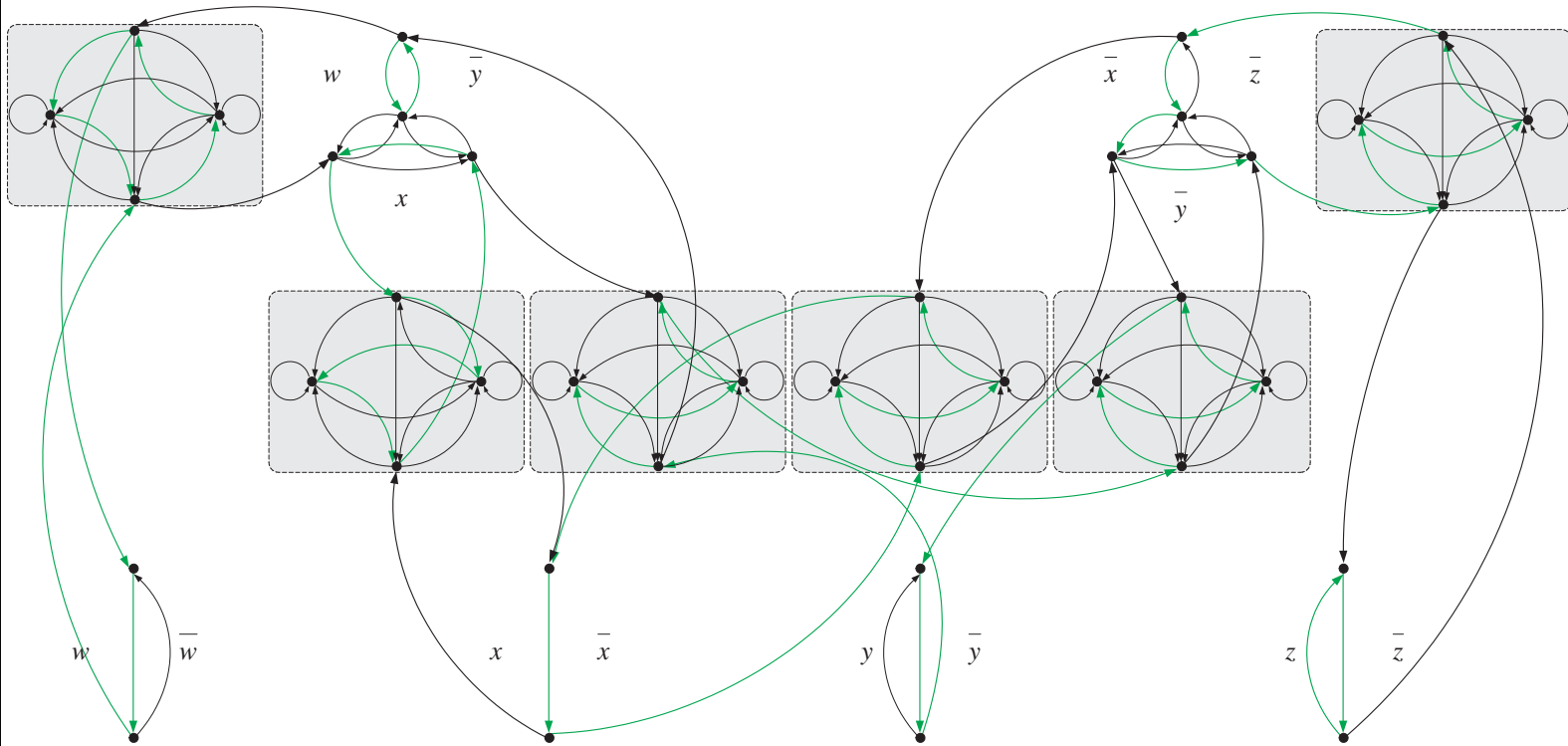$$w = 1, x = 0, y = 0, z = 1 \Leftrightarrow \text{One Cycle Cover}$$

The Proof: a Key Corollary (continued)

- Recall that there are $3m$ XOR gadgets.

- Each satisfying truth assignment to ϕ contributes 4^{3m} to the cycle count $\#H$.

- Hence
$$\#H = 4^{3m} \times \#\phi,$$
as desired.

# "$w = 1, x = 0, y = 0, z = 1$" Adds 4^6 to Cycle Count

The Proof (continued)

- We are almost done.

- The weighted directed graph H needs to be *efficiently* replaced by some unweighted graph G.

- Furthermore, knowing $\#G$ should enable us to calculate $\#H$ *efficiently*.

  - This done, $\#\phi$ will have been Turing-reducible to $\#G$.[a]

- We proceed to construct this graph G.

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[a]By way of $\#H$ of course.

The Proof: Construction of G (continued)

- Replace edges with weights 2 and 3 as follows (note that the graph cannot have parallel edges):
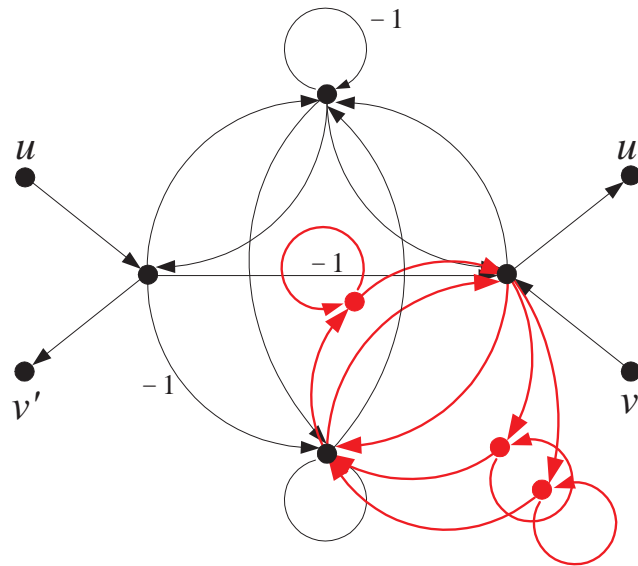
- The cycle count $\#H$ remains *unchanged*.

The Proof: Construction of G (continued)

- We move on to edges with weight -1.

- First, we count the number of nodes, M.

- Each clause gadget contains 4 nodes (p. 752), and there are m of them (one per clause).

- Each revised XOR gadget contains 7 nodes (p. 771), and there are $3m$ of them (one per literal).

- Each choice gadget contains 2 nodes (p. 763), and there are $n \le 3m$ of them (one per variable).

- So

$$M \le 4m + 21m + 6m = 31m.$$

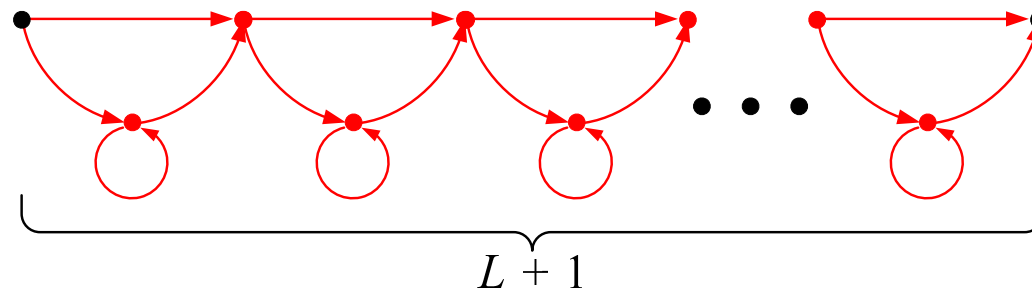# The Proof: Construction of G (continued)

- $\#H \leq 2^L$ for some $L = O(m \log m)$.

  - The maximum absolute value of the edge weight is 1.

  - Hence each term in the permanent is at most 1.

  - There are $M! \leq (31m)!$ terms.

  - Hence

  $$
  \begin{aligned}
  \#H &\leq \sqrt{2\pi(31m)} \left(\frac{31m}{e}\right)^{31m} e^{\frac{1}{12 \times (31m)}} \\
  &= 2^{O(m \log m)}
  \end{aligned}
  \tag{13}
  $$

  by a refined Stirling's formula.

The Proof: Construction of G (continued)

- Replace each edge with weight -1 with the following:

$$L + 1$$

- Each increases the number of cycle covers 2^{L+1}-fold.

- The desired unweighted G has been obtained.

The Proof (continued)

- $\#G$ equals $\#H$ after replacing each appearance -1 in $\#H$ with 2^{L+1}:

$$\#H \quad = \quad \cdots + \overbrace{(-1) \cdot 1 \cdot\cdots\cdot 1}^{\text{a cycle cover}} + \cdots ,$$

$$\#G \quad = \quad \cdots + \overbrace{2^{L+1} \cdot 1 \cdot\cdots\cdot 1}^{\text{a cycle cover}} + \cdots .$$

- Let $\#G = \sum_{i=0}^{n} a_i \times (2^{L+1})^i$, where $0 \le a_i < 2^{L+1}$.

- As $\#H \le 2^L$ even if we replace -1 by 1 (p. 773), each a_i equals the number of cycle covers with i edges of weight -1.

The Proof (concluded)

- We conclude that

$$\#H = a_0 - a_1 + a_2 - \cdots + (-1)^n a_n,$$

  indeed easily computable from $\#G$.

- We know $\#H = 4^{3m} \times \#\phi$ (p. 768).

- So

$$\#\phi = \frac{a_0 - a_1 + a_2 - \cdots + (-1)^n a_n}{4^{3m}}.$$

  – More succinctly,

$$\#\phi = \frac{\#G \bmod (2^{L+1} + 1)}{4^{3m}}.$$

Finis