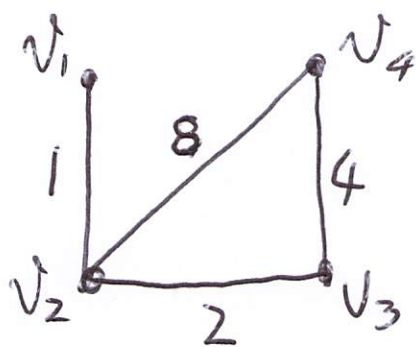


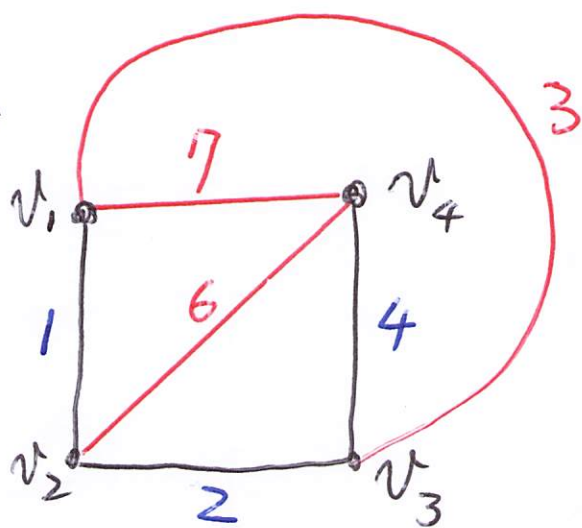
# A reduction to the metric case Kan-Mao Chao @2019

▲ The metric closure of a graph  $G=(V,E,w)$  is the complete graph  $\bar{G}=(V,V \times V,\bar{w})$ , where  $\bar{w}(u,v) = d_G(u,v)$  for all  $u,v \in V$ .

$G =$



$\bar{G} =$



Since  $\bar{w}(u,v) \leq w(u,v)$  for all  $u,v \in V$ ,

$$C(\text{mrcf}(\bar{G})) \leq C(\text{mrcf}(G)).$$

Surprisingly, it can be shown that

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$$C(\text{mrct}(\bar{G})) \geq C(\text{mrct}(G)).$$

Therefore,  $C(\text{mrct}(\bar{G})) = C(\text{mrct}(G))$ .

An edge  $(a, b)$  is called a **bad edge** if

$$(a, b) \notin E \text{ or } w(a, b) > \bar{w}(a, b).$$

executed at most  $O(n^2)$  times.

Algorithm: REMOVE\_BAD

Input: A spanning tree  $T$  of  $\bar{G}$ .

Output: A spanning tree  $Y$  of  $G$  such that  $C(Y) \leq C(T)$ .

Compute all-pairs shortest paths of  $G$ .

(I) while there exists a bad edge in  $T$

Pick a bad edge  $(a, b)$ . Root  $T$  at  $a$ .

/\* assume  $SP_G(a, b) = (a, x, \dots, b)$  and  $y$  is the parent of  $x$  \*/  
if  $b$  is not an ancestor of  $x$  then

$$Y^* \leftarrow T \cup (x, b) - (a, b); Y^{**} \leftarrow Y^* \cup \{(a, x)\} - \{(x, y)\};$$

else

$$Y^* \leftarrow T \cup (a, x) - (a, b); Y^{**} \leftarrow Y^* \cup \{(b, x)\} - \{(x, y)\};$$

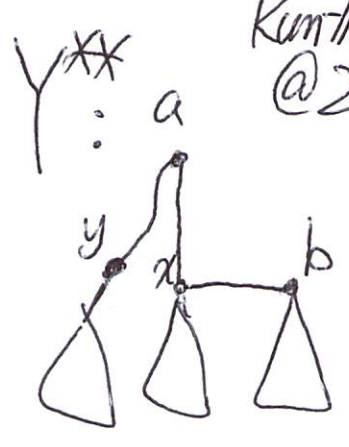
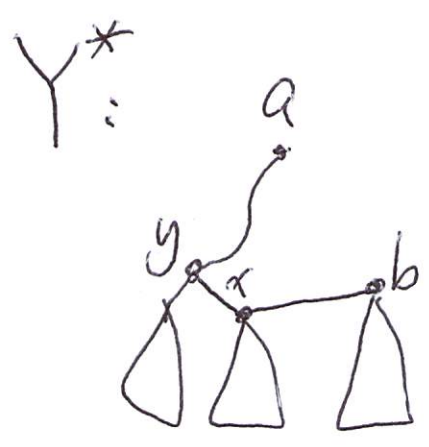
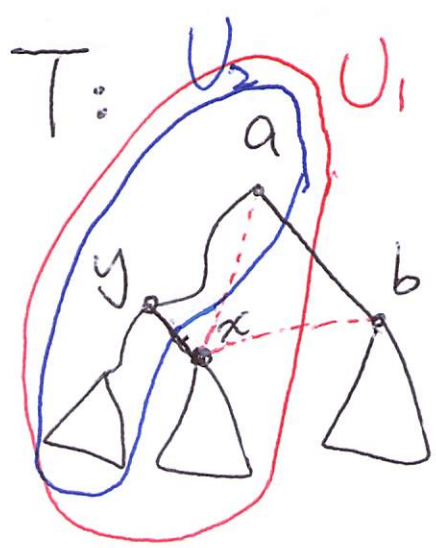
if  $C(Y^*) < C(Y^{**})$  then  $Y \leftarrow Y^*$  else  $Y \leftarrow Y^{**}$

(II)  $T \leftarrow Y$

$$C(Y) \leq C(T)$$

$$\min\{C(Y^*), C(Y^{**})\}$$

(2)



$$U_1 = V(T_a) - V(T_b)$$

If  $U_2 = V(T_a) - V(T_b) - V(T_x)$   
 $C(Y^*) > C(T)$

$$\Rightarrow (C(Y^*) - C(T)) / 2 = \sum_{u \in U_1} \sum_{v \in V(T_b)} (d_T(u, x) + \bar{w}(x, b) - (d_T(u, a) + \bar{w}(a, b))) > 0$$

$$\Rightarrow \sum_{u \in U_1} \sum_{v \in V(T_b)} (d_T(u, x) - d_T(u, a) - \bar{w}(a, x)) > 0$$

$$\Rightarrow |V(T_b)| \sum_{u \in U_1} (d_T(u, x) - d_T(u, a) - \bar{w}(a, x)) > 0$$

$$\Rightarrow \sum_{u \in U_1} (d_T(u, x) - d_T(u, a)) > |U_1| \cdot \bar{w}(a, x)$$

$$\Rightarrow \sum_{u \in U_2} (d_T(u, x) - d_T(u, a)) > |U_1| \bar{w}(a, x)$$

$$\textcircled{3} \Rightarrow \sum_{u \in U_2} (d_T(u, a) - d_T(u, x)) < -|U_1| \bar{w}(a, x)$$



$$(C(Y^{**}) - C(T))/2$$

$$= \sum_{u \in U_2} \sum_{v \in V(T_x)} (d_T(u, a) + \bar{w}(a, x) - d_T(u, x))$$

$$+ \sum_{u \in V(T_x)} \sum_{v \in V(T_b)} (d_T(u, x) + \bar{w}(x, b) - (d_T(u, a) + \bar{w}(a, b)))$$

 $\leq 0$ 

$$\leq \sum_{u \in U_2} \sum_{v \in V(T_x)} (d_T(u, a) - d_T(u, x) + \bar{w}(a, x))$$

$$= |U(T_x)| \left( \sum_{u \in U_2} (d_T(u, a) - d_T(u, x)) + \bar{w}(a, x) \right)$$

$$< -|U_1| \cdot \bar{w}(a, x)$$

$$+ |U_2| \bar{w}(a, x)$$

$$< \underbrace{|U(T_x)|}_{>0} \cdot \underbrace{(|U_2| - |U_1|)}_{<0} \cdot \underbrace{\bar{w}(a, x)}_{\geq 0}$$

$$\leq 0$$

It has been shown that

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$$\min \{ C(Y^*), C(Y^{**}) \}$$

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$$\leq C(T).$$

\* For any spanning tree  $T$  of  $\bar{G}$ ,

it can be transformed into a spanning

tree  $Y$  of  $G$  <sup>in  $O(n^3)$  time</sup> such that  $C(Y) \leq C(T)$ .

\* If there is a  $(HE)$ -approximation algorithm

for  $\Delta$ MRCT (for  $\bar{G}$ ), then there is a

$(HE)$ -approximation algorithm for MRCT (for  $G$ ).

\* The transformation from  $G$  to  $\bar{G}$

(metric closure of  $G$ ) can be done in

time  $O(n \log n + m)$  \*  $\mathcal{N} = O(n^2 \log n + mn)$ .

Dijkstra's algorithm

(5)