Tree Eccentricities, diameters, and radii

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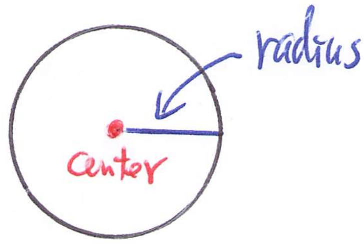
$G=(V,E,w)$



w: positive edge weights

$U \subseteq V$

$$D_G(v,U) = \max_{u\in U}\{d_G(v,u)\}$$

The eccentricity of v is the maximum of the distance to any vertex in the graph, i.e. $\max_{u\in V}\{d_G(v,u)\}$ or $D_G(v,V)$.

Diameter: $\max_{v\in V}\left\{\max_{u\in V}\{d_G(v,u)\}\right\} = \max_{v,u} d_G(v,u)$

also referred to as a path

Radius: $\min_{v\in V}\left\{\max_{u\in V}\{d_G(v,u)\}\right\}$

1 Center: the vertex with the minimum eccentricity

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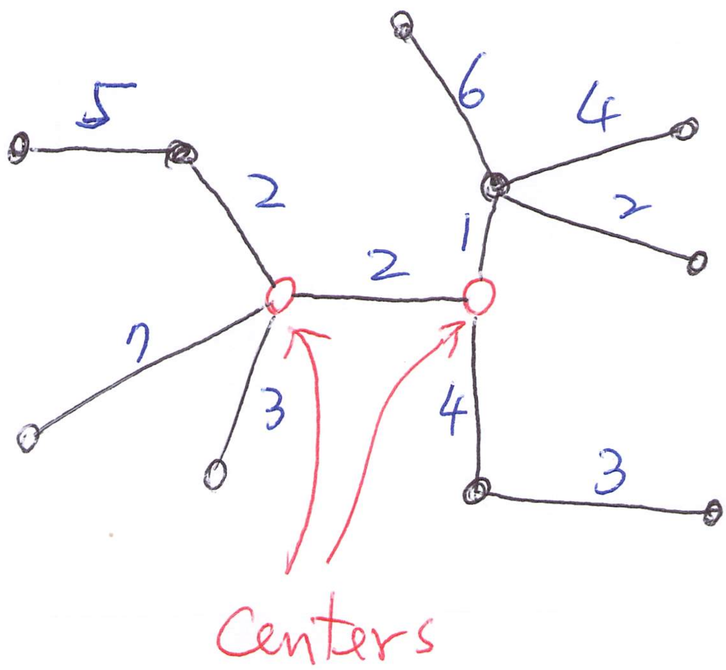
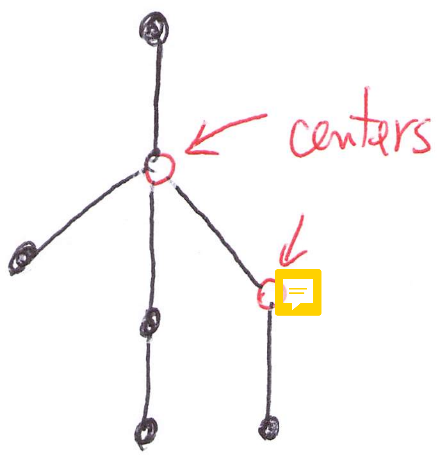


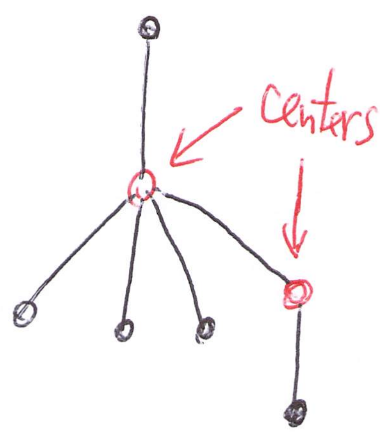
radius = 9

diameter = 16



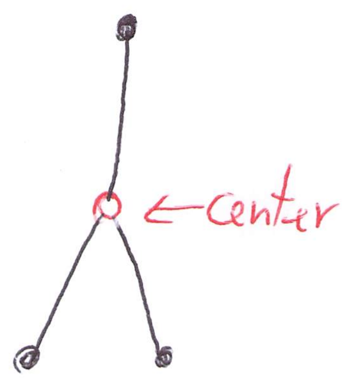
radius = 2

diameter = 4



radius = 2

diameter = 3



radius = 1

diameter = 2

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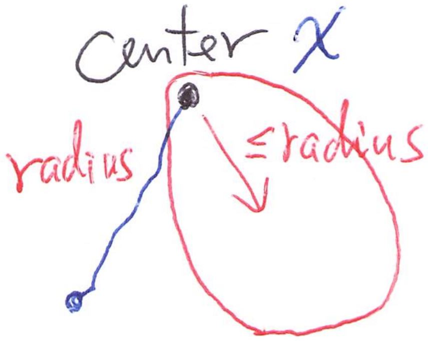
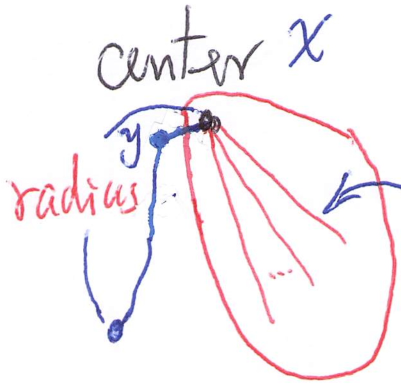
For an unweighted tree T,

$$2 \times \text{radius} - 1 \le \text{diameter} \le 2 \times \text{radius}.$$



$$d_T(u,v) \le d_T(u,x) + d_T(x,v)$$
$$\le \text{radius} + \text{radius}$$
$$= 2 \times \text{radius}$$

$\Rightarrow$ diameter $\le 2 \times$ radius
($\max_{u,v} d_T(u,v)$)



At least one branch $\ge$ radius -1

(Otherwise, we have a smaller radius by selecting y as the center.)

$$\Rightarrow \text{diameter} \ge \text{radius} + \text{radius} - 1$$
$$= 2 \times \text{radius} - 1$$

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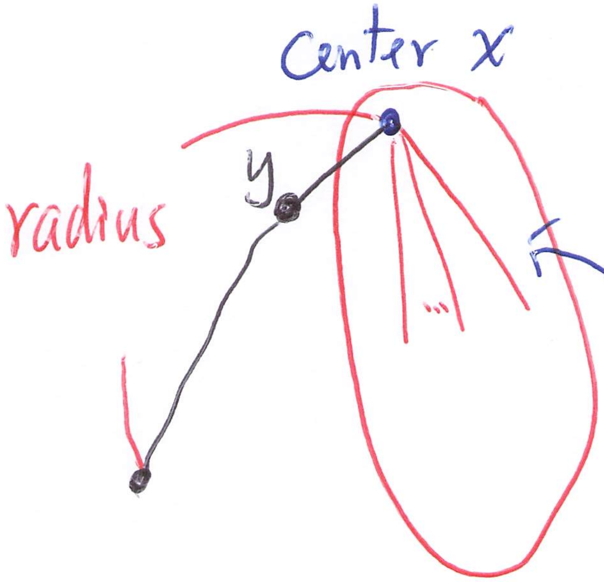
For positive weighted tree T,

$$2 \times \text{radius} - \max_{e \in E}\{w(e)\} \le \underline{\text{diameter} \le 2 \times \text{radius}}$$

$$d_T(u,v) \le d_T(u,\underset{\text{center}}{x}) + d_T(x,v) \le 2 \times \text{radius}$$



At least one branch $\ge \text{radius} - \max_{e \in E}\{w(e)\}$

(Otherwise, we have a smaller radius by selecting $\underline{y}$ as the center.)

$$\Rightarrow \text{diameter} \ge \text{radius} + \text{radius} - \max_{e \in E}\{w(e)\}$$

$$= 2 \times \text{radius} - \max_{e \in E}\{w(e)\}$$

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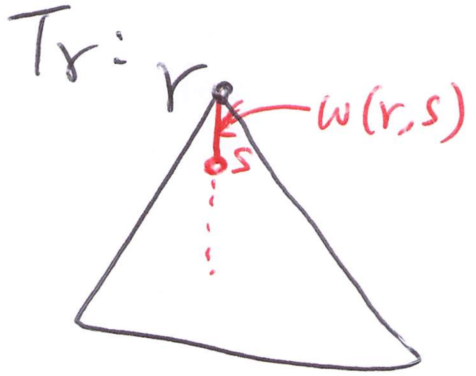
T_r : a subtree rooted at r



$$D_{T_r}(r, V(T_r)) = \max_{s \in child(r)} \{ D_{T_r}(s, V(T_s)) + w(r,s) \}$$

We can root the tree at any given vertex v and compute v's eccentricity in linear time by the above recurrence relation. (bottom-up Postorder traversal or recursive algorithm)

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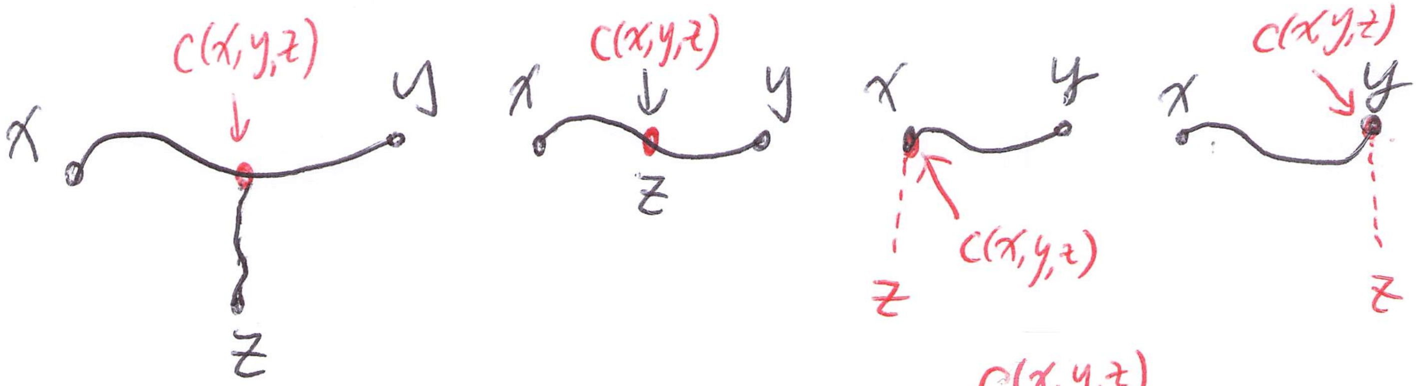
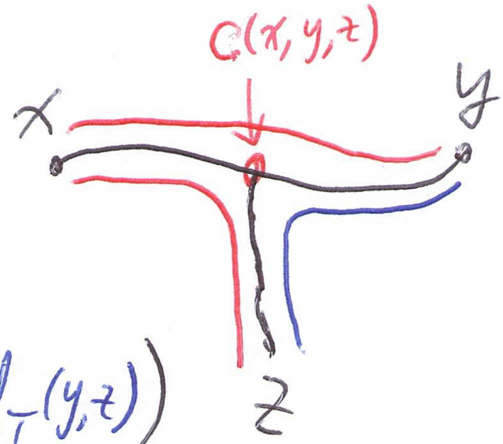
$$C(x,y,z) = SP_T(x,y) \cap SP_T(x,z) \cap SP_T(y,z)$$

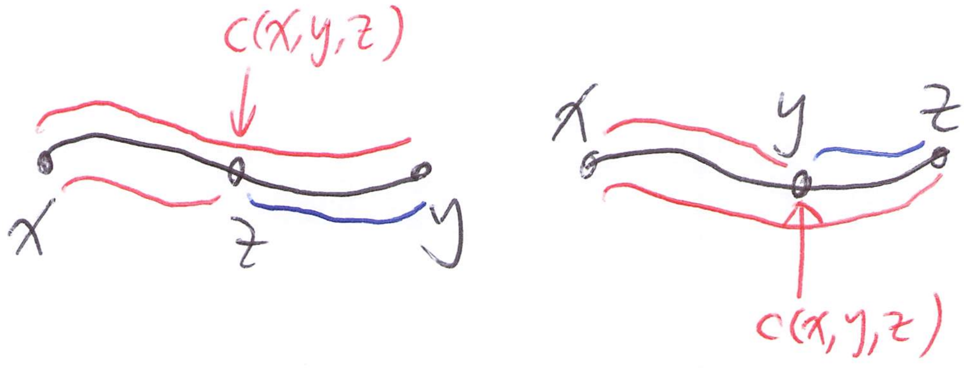


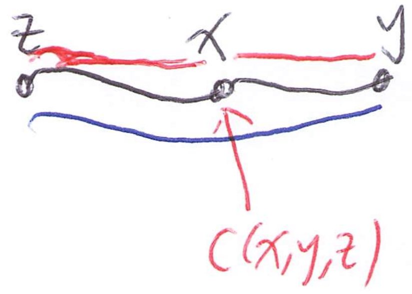


$$d_T(x, c(x,y,z))$$

$$= \frac{1}{2}\left(d_T(x,y) + d_T(x,z) - d_T(y,z)\right)$$

$$= \frac{1}{2}\left(d_T(x,y) + d_T(x,z) + d_T(y,z)\right) - d_T(y,z)$$





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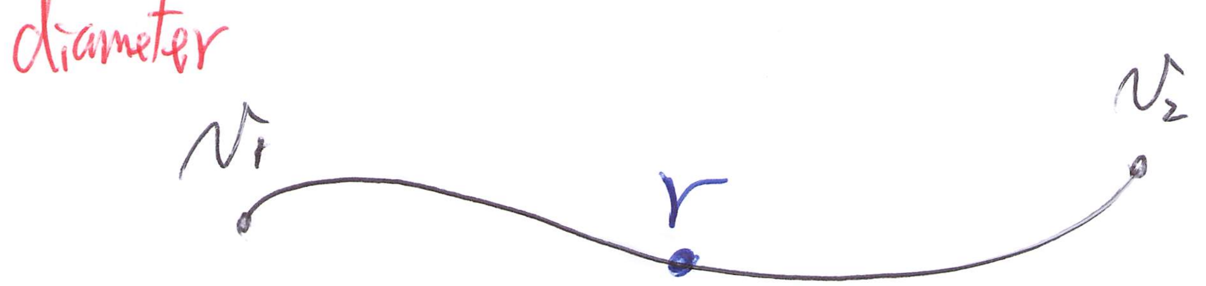
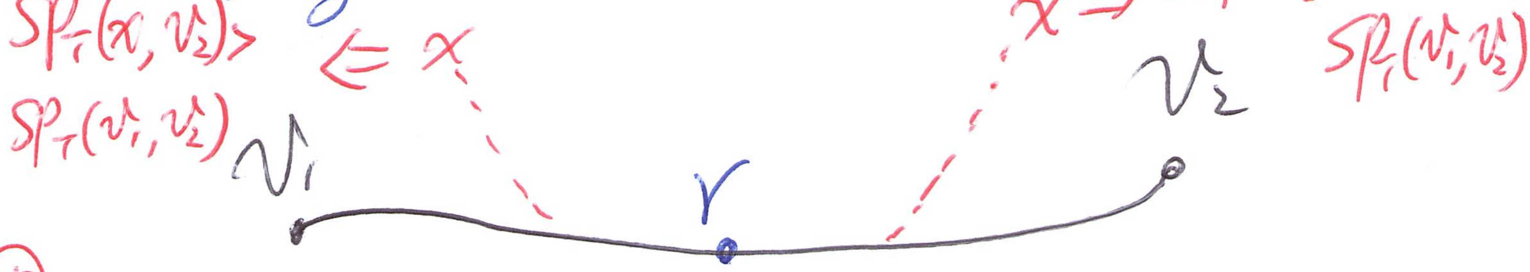
$SP_T(v_1, v_2)$: diameter

r : a vertex on $SP_T(v_1, v_2)$



For any vertex x, $d_T(x, r) \le \max\{d_T(r, v_1), d_T(r, v_2)\}$.

pf. If $d_T(x, r) > \max\{d_T(r, v_1), d_T(r, v_2)\}$, then $SP_T(x, v_1)$ or $SP_T(x, v_2)$ is a path longer than the diameter.



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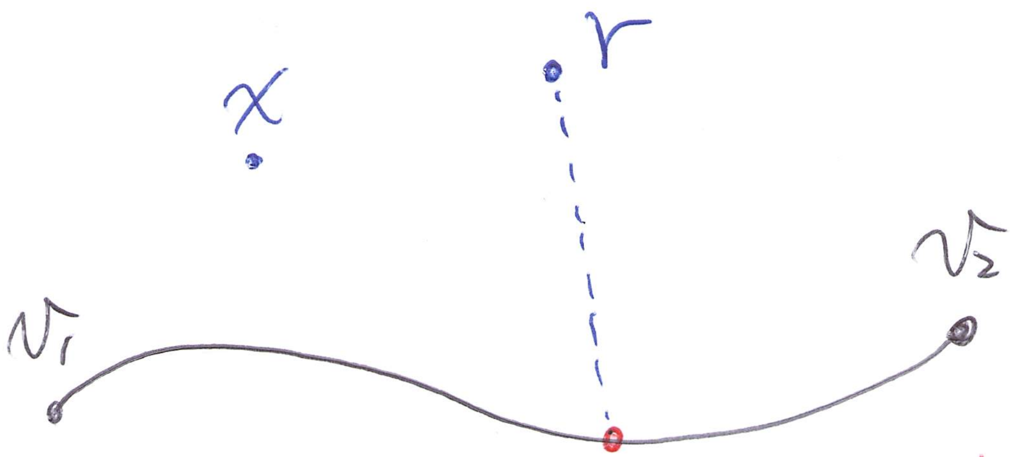
$SP_T(v_1, v_2)$: diameter

r : a vertex not on the diameter



$u = c(r, v_1, v_2)$

Wlog. let $d_T(u, v_1) \geq d_T(u, v_2)$.

$$\begin{aligned} d_T(x, r) &\leq d_T(x, \underbrace{u}_{\text{on the diameter}}) + d_T(u, r) \\ &\leq d_T(v_1, u) + d_T(u, r) \\ &= d_T(v_1, r) \end{aligned}$$

For any vertex, either on the diameter or not, one of the endpoints of a diameter

8 must be the farthest vertex.

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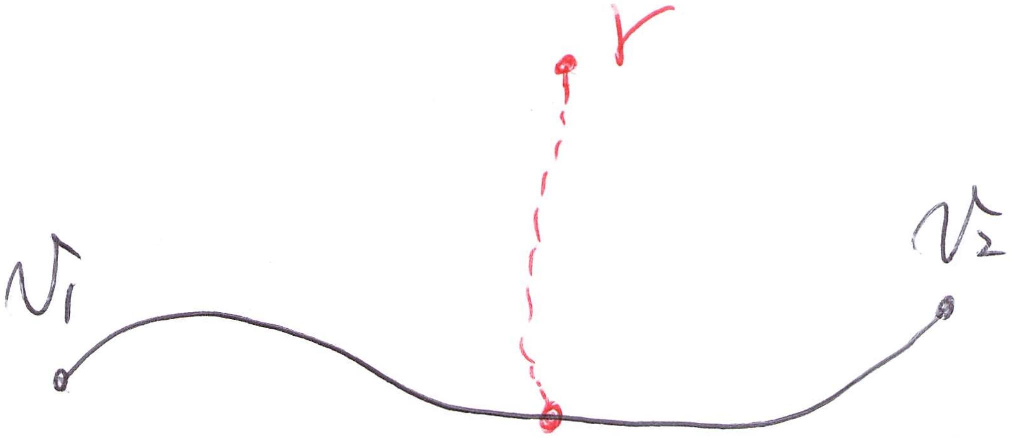
$SP_T(v_1, v_2)$: diameter

r: any vertex

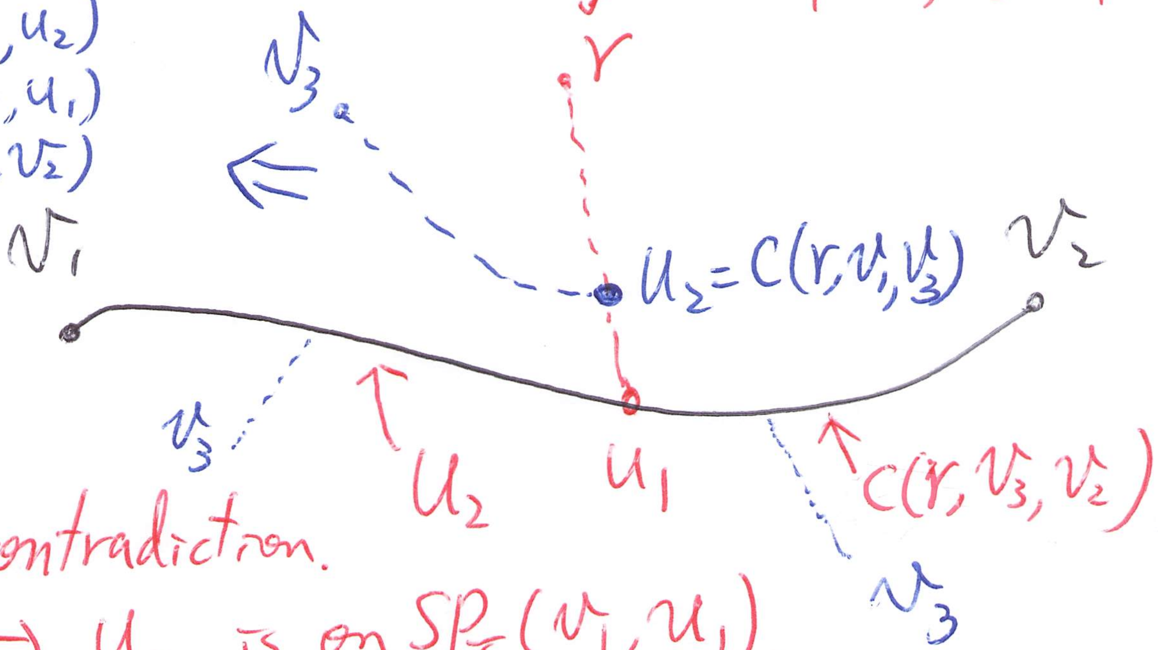
v_3: the vertex farthest to r $\Rightarrow$ v_3 must be an endpoint of a diameter.



$u_1 = c(r, v_1, v_2)$

Wlog, let $d_T(u_1, v_1) \geq d_T(u_1, v_2)$.

$d_T(v_1, v_3)$
$= d_T(v_1, u_2) + d_T(u_2, v_3)$
$= 2d_T(v_1, u_2)$
$> 2d_T(v_1, u_1)$
$\geq d_T(v_1, v_2)$



A contradiction.

$\Rightarrow u_2$ is on $SP_T(v_1, u_1)$.

If $c(r, v_3, v_2) \neq u_1$, then $SP_T(v_1, v_3)$ is a diameter.

Otherwise, $SP_T(v_3, v_2)$ is a diameter.

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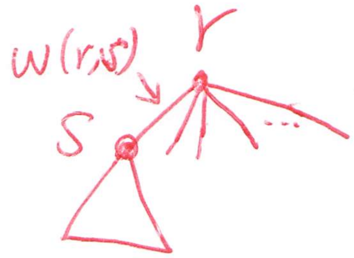
Algorithm: ECCENT(T_r)
Input: A tree $T_r = (V, E, w)$ rooted at r.
Output: The eccentricity of r in T_r.



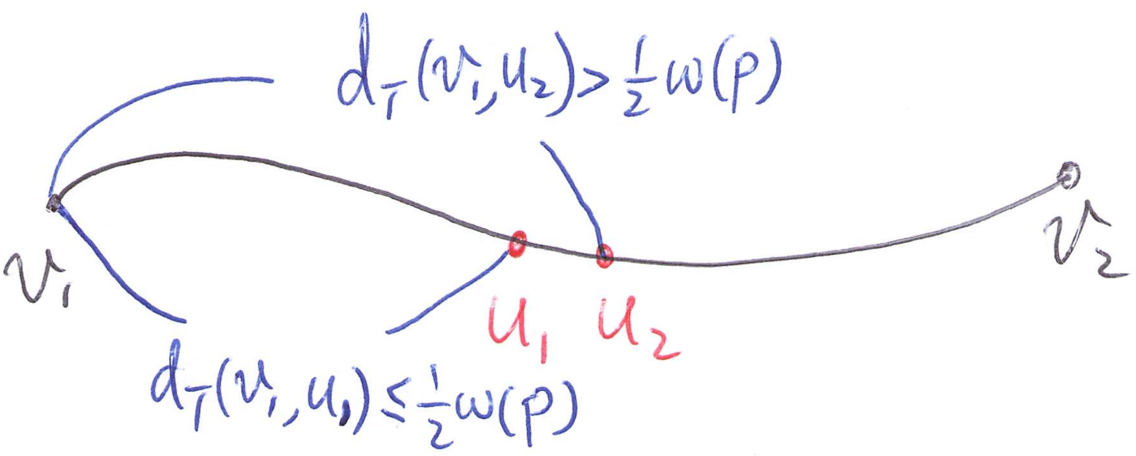
1: **if** r is a leaf **then**
 return 0;
2: **for** each child s of r **do**
 compute ECCENT(T_s) recursively;
3: **return** $\max_{s \in child(r)}\{\text{ECCENT}(s) + w(r,s)\}$.

Algorithm: TREEDIAMETER
Input: A tree $T = (V, E, w)$.
Output: The diameter of T.

an endpoint of a diameter

1: Root T at an arbitrary vertex r.
2: Use ECCENT to find the farthest vertex v to r.
3: Root T at v.
4: Use ECCENT to find the eccentricity of v.
5: Output the eccentricity of v as the diameter of T.

$P = SP_T(v_1, v_2)$: diameter



The eccentricity of u_1 is $d_T(u_1, v_2)$.
The eccentricity of u_2 is $d_T(u_2, v_1)$.

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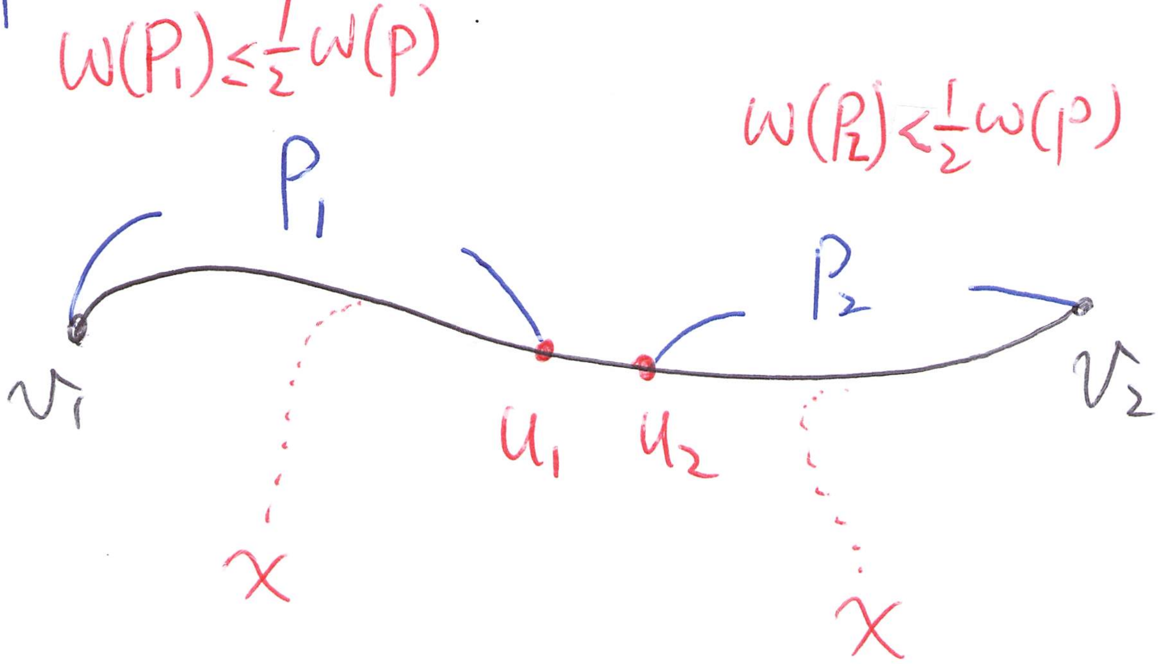
P:



For any vertex connected to P at a vertex in P_1, $d_T(x, v_2) > \underline{d_T(u_1, v_2)}$.

u_1's eccentricity

For any vertex connected to P at a vertex in P_2, $d_T(x, v_1) > \underline{d_T(u_2, v_1)}$.

u_2's eccentricity

$\Rightarrow \min\{d_T(\underline{\underline{u_1}}, v_2), d_T(\underline{\underline{u_2}}, v_1)\}$ is the radius.

One of them must be a center.

We have proved the following theorem.

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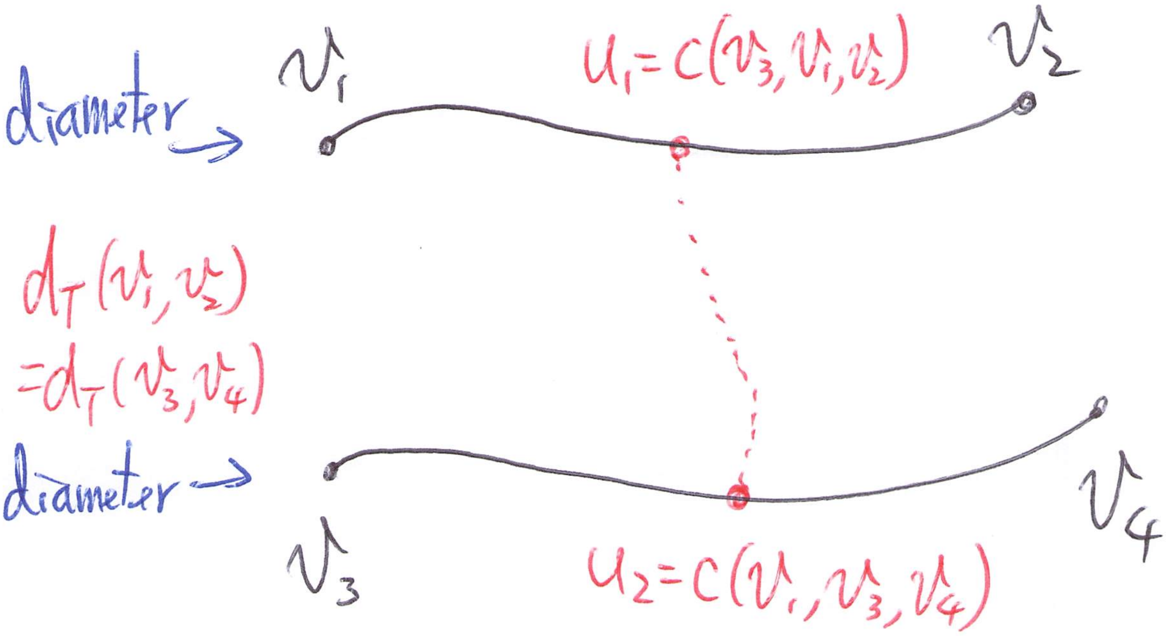
☆ **Thm.** The diameter, radius, and center of a tree can be computed in linear time.

Two diameters of a tree cannot be disjoint.



$$d_T(v_1, v_4) + d_T(v_2, v_3)$$

$$= d_T(v_1, v_2) + d_T(v_3, v_4) + \overset{>0}{2d_T(u_1, u_2)}$$

$$> 2d_T(v_1, v_2)$$

$$\Rightarrow d_T(v_1, v_4) > d_T(v_1, v_2) \text{ or } d_T(v_2, v_3) > d_T(v_1, v_2)$$

One of them is longer than the diameter. A contradiction!

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All diameters of a tree share at least one common vertex.