Data Structures and Algorithms

(資料結構與演算法)

Lecture 1: Algorithm

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Roadmap

1 the one where it all began

Lecture 1: Algorithm

- definition of algorithm
- pseudo code of algorithm
- criteria of algorithm
- correctness proof of algorithm
- 2 the data structures awaken
- 3 fantastic trees and where to find them
- 4 the search revolutions
- 5 sorting: the final frontier

definition of algorithm

Name Origin of Algorithm

O NO TIA O O CCCPO

Muhammad ibn Mūsā al-Kwārizmī on a Soviet Union stamp

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algorithm

- named after al-Kwārizmī (780–850), Persian mathematician and father of algebra
- algebra: rules to calculate with symbols
- algorithm: instructions to compute with variables

algorithm: recipe-like instructions for computing

Recipe for Cooking Dish

Cookbook/Hamburger - Wikibooks

http://ex.wikibooks.org/wiki/Cookbook/Ha

Cookbook:Hamburger

From Wikibooks

Cookbook | Recipe Index | Mest recipes

A hamburger (or, loss frequently, a hamburg, or in the United Kingdoos, a beetburger) is a statiant on a materials involving a party of ground must that is

ingredients

• 500g (1.1 Ib) minced (ground) heef

NO(g (i.i. lb) mineral (ground) herd is brefs and spiers (spinned) is shown and spiers (spinned) is shore (spinned) is should (strace, spinnels, alfalfa sprouts, tomato, onion etc. - optional) is hamburger but for each burger.



Procedure

- Add the beef to a food processor for approximately 10 seconds.
 Now add your herbs under space is toxic. Depending on the quality of your leads beef, you may shall an all some beef shock in superve for these.
 The best of the processor for mobile 20 seconds or said fully mixed.
 This has been for the processor for mobile 20 seconds or said fully mixed.
 This has been for the processor for mobile 20 seconds or said fully mixed.
 This has been for the food flowed remaind, saids were your mixed to your.
- 3. Most in the final processor or another 2 of focusine or treat easy strategy.

 4. If you brought the best disorably promise were you mis in your someoning with. From may which is the light you. The final the filter, we you will be first the strategy with the many which is the light you had use lapsals, test the growth bett source for some bick. Notice If you ald use lapsals, test the growth bed will then superce out the centra jack when forming paties.

 5. Remove the better for some for body necessor and shape by had into live lapsacs. You should get between 4-6 bargers from 500g (3.).
- iterative the entire that the stood processor and stage of man and outgots, two anomal got notices are trapped train 740g (1.18) of beef.
 The bugger, can be fitted (about 5 mins on each side for buggers which aten't too thick), grilled (same times as for frying), or barbogot.
- backopied.

 T. Beamer your bragers are fully cooked through before serving. If your bragers are quite thick or if you are stream, you can ceit one eyes to ensure the mode are browned. If the insides are rul, there is a charer that the mod is not fully cooked, one eyes to ensure the small in set fully cooked, and the cooking the

For further serving suggestions, see the Wikipedia article on hamburgers.

Notes, tips and variations

• Yes can use almost are tree of miscod (resemb) must be make hamburers, including next, chicken, butter, lamb, big

Ver care or almost any type of miscod groundy most to make handwagers, including peek, chickers, makey, lanth, bious, revisions, centrick, over no most midmales and a Queen.

1 if your beapers full apert, adding an egg sold, will help drop it superfor. Dissipa four ground bod's will also help.

1 if your beapers full apert, adding a negg sold, will help drop it to grounder of you begree before contain.

1 if you may vish to perspections with ackeding devote in the control of you begree before contain.

2 if you which may work well in handwagers include that, paper, chill (nither front or provietr). Worcestenber Succe and soy asser, Department or full good combinations.

succ. Expediment to fing good commutation.

Almost stay below that well, including boul, engages and gualety.

Burgers can also be moded on a grill, Smoked burgers will appear red and glazed on the conside, but between on the inside.

Sociologia to be specified profile grill as accordate way to said in the fluorieth judges.

Variation, Adding most and spaces together in a lovel and mixing by hand send the spices are distributed may produce better rooths. This will also stay your despire them falling query after fluoriething and the spices are distributed may produce better rooths. This will allow stay your despire them falling query.

Links

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a recipe for hamburger on Wikibooks

figure by Gentgeen,

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recipe

Wikipedia: a set of instructions that describes how to prepare or make something, especially a dish of prepared food

recipe: instructions to complete a (cooking) task

Sheet Music for Playing Instrument



first page of manuscript of Bach's lute suite in G minor figure licensed

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sheet music

Wikipedia: handwritten or printed form of musical notation ... to indicate the pitches, rhythms or chords of a song

sheet music: instructions to play instrument (well)

Kifu for Playing Go



a Japanese kifu

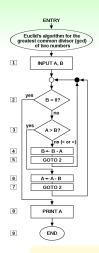
figure by Velobici,

licensed under CC BY-SA 4.0 via Wikimedia Commons

kifu

go game record of steps that describe how the game had been played

kifu: instructions to mimic/learn to play go (professionally)



Algorithm for Computing

flowchart of Euclid's algorithm for calculating the greatest common divisor (g.c.d.) of two numbers

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algorithm

Wikipedia: algorithm is a finite sequence of well-defined, computer-implementable instructions, typically to solve a class of problems or to perform a computation

algorithm ~ computing recipe:

(computable) instructions to solve a computing task

efficiently/correctly

Which of the following in the kitchen is the best metaphor for an algorithm?

- 1 recipe
- 2 chef
- garbage
- 4 meat

Which of the following in the kitchen is the best metaphor for an algorithm?

- 1 recipe
- 2 chef
- garbage
- 4 meat

Reference Answer: (1)

algorithm \sim computing recipe:

(computable) instructions to solve a computing task efficiently/correctly

pseudo code of algorithm

Pseudo Code for Get-Min-Index

C Version

```
/* return index to min. element
    in arr[0] ... arr[len-1] */
int getMinIndex
        (int arr[], int len){
    int i;
    int m=0;
    for(i=0;i<len;i++){
        if (arr[m] > arr[i]){
            m = i;
        }
    }
    return m;
}
```

Pseudo Code Version

```
Get-Min-Index(A)

1 m = 1

2 for i = 2 to A. length

3 // update if i-th element smaller

4 if A[m] > A[i]

5 m = i

6 return m
```

pseudo code: spoken language of programming

Bad Pseudo Code: Too Detailed

Unnecessarily Detailed

```
Get-Min-Index(A)

1 m = 1

2 for i = 2 to A. length

3  // update if i-th element smaller

4  Am = A[m]

5  Ai = A[i]

6  if Am > Ai

7  m = i

8  else

9  m = m

10 return m
```

Concise

```
Get-Min-Index(A)

1 m = 1

2 for i = 2 to A. length

3  // update if i-th element smaller

4  if A[m] > A[i]

5  m = i

6 return m
```

goal of pseudo code: communicate efficiently

Bad Pseudo Code: Too Mysterious

Unnecessarily Mysterious

```
Get-Min-Index(A)

1 x = 1

2 for xx = 2 to A. length

3 4 if A[x] > A[xx]

5 xx = x

6 return xx
```

Clear

```
Get-Min-Index(A)

1 m = 1 // store current min. index

2 for i = 2 to A. length

3 // update if i-th element smaller

4 if A[m] > A[i]

5 m = i

6 return m
```

goal of pseudo code: communicate correctly

Bad Pseudo Code: Too Abstract

Unnecessarily Abstract

Get-Min-Index(A)

- 1 m = 1 // store current min. index
- 2 run a loop through *A* that updates *m* in every iteration
- 3 return m

Concrete

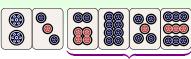
Get-Min-Index(A)

- 1 m = 1 // store current min. index
- 2 for i = 2 to A. length
- 3 // update if i-th element smaller
- 4 if A[m] > A[i]
- 5 m=i
- 6 return m

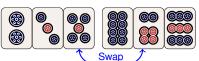
goal of pseudo code: communicate effectively

From Get-Min-Index to Selection-Sort

Get-Min-Index(A, ℓ, r) $m = \ell$ // store current min. index for $i = \ell + 1$ to r // update if i-th element smaller if A[m] > A[i]m = ireturn m



Get-Min-Index



Good Pseudo Code

- modularize, just like coding
- depends on speaker/listener
- usually no formal definition

Selection-Sort(A)

```
for i = 1 to A. length
     m = \text{Get-Min-Index}(A, i, A. length))
     if i \neq m
           \operatorname{Swap}(A[i], A[m])
return A // which has been sorted in place
```

follow any textbook if you really need a definition

Quick Demo of Selection Sort

instructions	A[1]	A[2]	A[3]	A[4]	A[5]	<i>A</i> [6]
$m \stackrel{2}{\leftarrow} \text{Get-Min-Index}(A, 1, 6)$ Swap(A[1], A[2])						
$m \stackrel{2}{\leftarrow} \text{Get-Min-Index}(A, 2, 6)$ $\frac{\text{Swap}(A[2], A[2])}{\text{Swap}(A[2], A[2])}$						
$m \stackrel{5}{\leftarrow} \text{Get-Min-Index}(A, 3, 6)$ Swap(A[3], A[5])					00	
$m \stackrel{5}{\leftarrow} \text{Get-Min-Index}(A, 4, 6)$ Swap(A[4], A[5])					88	
$m \stackrel{5}{\leftarrow} \text{Get-Min-Index}(A, 5, 6)$ $\frac{\text{Swap}(A[5], A[5])}{\text{Swap}(A[5], A[5])}$						
$m \stackrel{6}{\leftarrow} \text{Get-Min-Index}(A, 6, 6)$ $\frac{\text{Swap}(A[6], A[6])}{\text{Swap}(A[6], A[6])}$						

suggestion: draw, don't just watch

Which of the following can be used to describe good pseudo code?

- clear
- 2 concise
- 3 concrete
- 4 all of the above

Which of the following can be used to describe good pseudo code?

- clear
- 2 concise
- 3 concrete
- 4 all of the above

Reference Answer: 4

Have fun communicating with other programmers using good pseudo code! :-)

criteria of algorithm

Criteria of Recipe



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Cocktail Recipe:

Screwdriver (from Wikipedia)

inputs: 5 cl vodka, 10 cl orange juice

- mix inputs in a highball glass with ice
- 2 garnish with orange slice and serve

output: a glass of delicious cocktail

- input: ingredients
- definiteness: clear instructions
- effectiveness: feasible instructions
- finiteness: completable instructions
- output: delicious drink

algorithm \sim recipe: same five criteria for algorithm

(Knuth, The Art of Computer Programming)

Input of Algorithm

... quantities which are given to it initially before the algorithm begins.

These inputs are taken from specified sets of objects. (Knuth, TAOCP)

```
Get-Min-Index(A)

1 m = 1 // store current min. index

2 for i = 2 to A. length

3 // update if i-th element smaller

4 if A[m] > A[i]

5 m = i

6 return m
```

one algorithm, many uses (on different legal inputs)

Definiteness of Algorithm

Each step of an algorithm must be precisely defined; the actions to be carried out must be rigorously & unambiguously specified. (Knuth, TAOCP)

Clear

```
Get-Min-Index(A)

1 m = 1 // store current min. index

2 for i = 2 to A. length

3 // update if i-th element smaller

4 if A[m] > A[i]

5 m = i
```

Ambiguous

```
Get-Zero-Index(A)

1
2 for i = 1 to A. length
3
4 if A[m] is almost zero
5 return m
6 // what to return here?
```

definiteness: clarity of algorithm

Effectiveness of Algorithm

... all of the operations to be performed in the algorithm must be sufficiently basic that they can in principle be done exactly and in a finite length of time by a man using paper and pencil. (Knuth, TAOCP)

Effective

```
Get-Min-Index(A)

1 m = 1 // store current min. index

2

3 for i = 2 to A. length

4 // update if i-th element smaller

5

6 if A[m] > A[i]

7 m = i

8

9 return m
```

Ineffective

```
Get-Min-Index-Robot(A)

1 m = 1 // store current min. index

2 form a ball Bm of weight A[1]

3 for i = 2 to A. length

4 // update if i-th element smaller

5 form a ball Bi of weight A[i]

6 if Bi is lighter than Bm on a balance

7 m = i

8 replace Bm with Bi

9 return m
```

forming a ball (& other actions) are arguably ineffective on typical computers

Finiteness of Algorithm

An algorithm must always terminate after a finite number of steps . . . a very finite number, a reasonable number. (Knuth, TAOCP)

```
Get-Min-Index(A)

1 m = 1 // store current min. index

2 for i = 2 to A. length

3 // update if i-th element smaller

4 if A[m] > A[i]

5 m = i

6 return m
```

finiteness (& efficiency): often need analysis for sophisticated algorithms (to be taught later)

Output of Algorithm

... quantities which have a specified relation to the inputs

(Knuth, TAOCP)

```
Get-Min-Index(A)

1 m = 1 // store current min. index

2 for i = 2 to A. length

3 // update if i-th element smaller

4 if A[m] > A[i]

5 m = i

6 return m
```

output (correctness): need proving with respect to requirements

What best describes the input/output relationship of the selection sort algorithm below?

- input: an ascending array; output: the same array sorted in descending order
- input: an arbitrary array; output: the same array sorted in descending order
- input: an arbitrary array; output: the same array sorted in ascending order
- 4 none of the other choices

```
Selection-Sort(A)

1 for i = 1 to A. length

2 m = \text{Get-Min-Index}(A, i, A. \text{length}))

3 if i \neq m

4 \text{Swap}(A[i], A[m])

5 return A / \!\!/ which has been sorted in place
```

What best describes the input/output relationship of the selection sort algorithm below?

- input: an ascending array; output: the same array sorted in descending order
- input: an arbitrary array; output: the same array sorted in descending order
- input: an arbitrary array; output: the same array sorted in ascending order
- 4 none of the other choices

```
Selection-Sort(A)

1 for i = 1 to A. length

2 m = \text{Get-Min-Index}(A, i, A. \text{length})
```

3 if $i \neq m$ 4 Swap(A[i], A[m]) 5 return A # which has been sorted in place

Reference Answer: (3)

The selection sort algorithm re-arranges an arbitrary array into ascending order.

correctness proof of algorithm

Claim



figure by Nick Youngson, licensed CC BY-SA 3.0 via Picpedia.Org

```
Get-Min-Index(A)

1 m = 1 // store current min. index

2 for i = 2 to A. length

3 // update if i-th element smaller

4 if A[m] > A[i]

5 m = i

6 return m
```

Correctness of Get-Min-Index

Upon exiting Get-Min-Index(A),

$$A[m] = \min_{1 \le j \le n} A[j]$$

with n = A. length

claim: math, statement that declares correctness

Invariant

invariants when constructing fractals figures by Johannes Rössel,

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Correctness of Get-Min-Index

Upon exiting Get-Min-Index(A),

$$A[m] = \min_{1 \le j \le n} A[j]$$

with n = A. length



Get-Min-Index(A)

return m

```
    1 m = 1 // store current min. index
    2 for i = 2 to A. length
    3 // update if i-th element smaller
    4 if A[m] > A[i]
    5 m = i
```

Invariant within Get-Min-Index

Upon finishing the loop with i = k, denote m by m_k ,

$$A[m_k] < A[j] \text{ for } j = 1, 2, ..., k$$

(loop) invariant: property that algorithm maintains

Proof of Loop Invariant

Mathematical Induction

Base

when i = 2, invariant true because

- assume invariant true for i = t 1
- when i = t.

o if
$$A[m_{t-1}] > A[t] \Rightarrow m_t = t$$

$$\begin{array}{ll}
A[m_t] &= A[t] &\leq A[t] \\
&< A[m_{t-1}] &\leq A[j] \text{ for } j < t
\end{array}$$

$$o if A[m_{t-1}] \leq A[t] \Rightarrow m_t = m_{t-1}$$

$$\begin{array}{ll}
A[m_t] &= A[m_{t-1}] & \leq A[t] \\
&= A[m_{t-1}] & \leq A[j] \text{ for } j < t
\end{array}$$

—by mathematical induction, invariant true for i = 2, 3, ..., k

Get-Min-Index(A)

1 m = 1// store current min. index

2 for i = 2 to A. length

3 // update if i-th element smaller

4 if A[m] > A[i]5 m = i

6 return m

Correctness of Get-Min-Index



Invariant within Get-Min-Index

Upon finishing the loop with i = k, denote m by m_k ,

$$A[m_k] \le A[j] \text{ for } j = 1, 2, \dots, k$$

proof of (loop) invariants \Rightarrow correctness claim of algorithm

 \Rightarrow

Which of the following is a loop invariant to selection sort?

```
Selection-Sort(A)

1 for i = 1 to A. length

2 m = \text{Get-Min-Index}(A, i, A. \text{length}))

3 if i \neq m

4 \text{Swap}(A[i], A[m])

5 return A // which has been sorted in place
```

- Upon finishing the loop with i = k, $A[1] \ge A[2] \ge ... \ge A[k]$.
- 2 Upon finishing the loop with i = k, $A[1] \le A[2] \le ... \le A[k]$.
- **3** Upon finishing the loop with i = k, $A[k + 1] \ge ... \ge A[A.length]$.
- 4 Upon finishing the loop with i = k, $A[k + 1] \le ... \le A[A. length]$.

Which of the following is a loop invariant to selection sort?

```
Selection-Sort(A)

1 for i = 1 to A. length

2 m = \text{Get-Min-Index}(A, i, A. \text{length}))

3 if i \neq m

4 \text{Swap}(A[i], A[m])

5 return A /\!\!/ which has been sorted in place
```

- **1** Upon finishing the loop with i = k, $A[1] \ge A[2] \ge ... \ge A[k]$.
- 2 Upon finishing the loop with i = k, $A[1] \le A[2] \le ... \le A[k]$.
- **3** Upon finishing the loop with i = k, $A[k + 1] \ge ... \ge A[A. length]$.
- 4 Upon finishing the loop with i = k, $A[k + 1] \le ... \le A[A. length]$.

Reference Answer: 2

The selection sort algorithm essentially picks the smallest element, the 2nd-smallest, and so on, and place them orderly. You can prove the loop invariant by mathematical induction.

Summary

Lecture 1: Algorithm

- definition of algorithm instructions to complete a task by computer
- pseudo code of algorithm communicate efficiently/correctly/effectively
- criteria of algorithm input, definite, effective, finite, output
- correctness proof of algorithm from (loop) invariants to claims
- next: data structures and their connections to algorithms