

$$* \quad P_*(S_t | \Sigma_{t-1}) = \frac{P_*(\Sigma_t)}{P_*(\Sigma_{t-1})} \in (0, 1)$$

$$P_*(\lambda | \Sigma_{t-1}) = 1 - P_*(0 | \Sigma_{t-1}) - P_*(1 | \Sigma_{t-1})$$

\* if do probabilistic prediction

$$\hat{S}_t = \begin{cases} 1 \\ 0 \\ \lambda \end{cases} \quad \text{with prob.} \quad \begin{matrix} P_*(1 | \Sigma_{t-1}) \\ P_*(0 | \Sigma_{t-1}) \\ P_*(\lambda | \Sigma_{t-1}) \end{matrix}$$

$$\boxed{\hat{S}_t \neq S_t} \quad \text{w/ prob.} \quad \underbrace{1 - P_*(S_t | \Sigma_{t-1})}_{q_t}$$

error  $e_t$

$$\begin{aligned} \underbrace{E\left(\sum_{t=1}^T e_t\right)}_{\text{total error}} &= \sum_{t=1}^T q_t \\ &= \sum_{t=1}^T \left(1 - \frac{P_*(\Sigma_t)}{P_*(\Sigma_{t-1})}\right) \\ &\leq \sum_{t=1}^T -\ln \frac{P_*(\Sigma_t)}{P_*(\Sigma_{t-1})} \\ &= -\sum_{t=1}^T (\ln P_*(\Sigma_t) - \ln P_*(\Sigma_{t-1})) \end{aligned}$$

$$= -\ln P_*(\Sigma_T) + \ln P_*(\Sigma_0)$$

$$\leq -\ln 2^{-\underbrace{K_*(\Sigma_T)}} \downarrow$$

(prefix) Kolmogorov complexity  
of generating  $\Sigma_T^*$

$$= K_*(\Sigma_T) \cdot \ln 2$$

$$\leq \underbrace{|\text{actual program that generates } \Sigma_\infty|}_{\text{"constant"}} \ln 2$$

$$\text{avg. error} \leq \frac{1}{T} \cdot \text{constant} \Rightarrow O(\text{complexity of } \Sigma_\infty) \quad \text{"examples" to learn (Solomonoff. Inference)}$$

Subject : .....

\* if do deterministic prediction

$$\tilde{s}_t = \operatorname{argmax}_{s \in \{1,0,2\}} P_*(s | \Sigma_{t-1})$$

$$\text{if } [\tilde{s}_t \neq s_t] \Rightarrow P_*(s_t | \Sigma_{t-1}) < \frac{1}{2}$$

error  $\tilde{e}_t$

$$\Leftrightarrow q_t \geq \frac{1}{2}$$

$$\sum_{t=1}^T \tilde{e}_t \leq \sum_{t=1}^T [q_t \geq \frac{1}{2}]$$

$$\leq \sum_{t=1}^T 2 \cdot q_t$$

⋮

\* learning :

consider a hypothesis set  $\mathcal{H} = \{h\}$

and examples  $D = \{(x_n, y_n = h_*(x_n))\}_{n=1}^N$

generated iid from some  $P(x)$  and  $h_* \in \mathcal{H}$

called realizable

$\Sigma \rightarrow$  examples  $(x_n, y_n)$

programs  $\rightarrow$  hypotheses  $h$

program that generates  $\Sigma_0 \rightarrow$  generating hypothesis  $h_*$

error on future bits  $\rightarrow$  error on future  $x \sim P(x)$  and  $y = h_*(x)$

$$e(h) = \mathbb{E}_{x \sim P(x)} ([h(x) \neq h_*(x)])$$

goal: find  $g \in \mathcal{H}$  s.t.  $e(g) \leq \epsilon$  w/ high probability  
how many examples ( $N$ ) needed?

\* conclusion

if always choosing one  $g \in \mathcal{H}$  s.t.  $g(x_n) = y_n$  for all  $n$   
 (one must exist because  $h_x \in \mathcal{H}$ )

and have a preference  $p(h)$  on  $h \in \mathcal{H}$ w/  $p(h) \geq 0$ ,  $\sum_h p(h) = 1$  "prior"then for any given  $0 < \delta \leq 1$ 

$$\Pr_{D \sim P^N} \left( \underbrace{\forall h \text{ s.t. } e_D(h) = 0}_{\text{i.e. "g"}}, e(h) \leq \frac{1}{N} \left( \ln \frac{1}{\delta} + \ln \frac{1}{p(h)} \right) \right) \geq 1 - \delta$$

\* proof: consider one  $h$  with  $e(h) = \lambda$ 

$$\Pr(e_D(h) = 0) = (1 - \lambda)^N = e^{N \ln(1 - \lambda)} \leq e^{-N\lambda}$$

$$\text{want: } \Pr(e_D(h) = 0) < \delta \cdot p(h)$$

 $\Uparrow$ 

$$e^{-N\lambda} < \delta \cdot p(h)$$

 $\Uparrow$ 

$$\lambda > \frac{1}{N} \left( \ln \frac{1}{\delta} + \ln \frac{1}{p(h)} \right)$$

So w/ prob  $\geq 1 - \delta \cdot p(h)$ 

$$e(h) \leq \frac{1}{N} \left( \ln \frac{1}{\delta} + \ln \frac{1}{p(h)} \right)$$

 $\Downarrow$  union bound over  $e_D(h) = 0$ 

$$\text{w/ prob } \geq 1 - \delta \sum_{e_D(h)=0} p(h) \geq 1 - \delta$$

$$\forall h \text{ w/ } e_D(h) = 0, e(h) \leq \frac{1}{N} \left( \ln \frac{1}{\delta} + \ln \frac{1}{p(h)} \right)$$

# examples (N)

proportion to

self information  $\ln \frac{1}{p(h)}$ to reach certain  $e(h)$  guarantee