

* probabilistic pigeon hole principle
 $0 < \delta < 1$ if $H_\delta(\bar{X}) > H_0(Y)$

there is no encoder ν & decoder μ
 s.t. $\Pr(\mu(\nu(x)) = x) \geq 1 - \delta$

ν : deterministic $\bar{X} \rightarrow Y$

μ : deterministic $Y \rightarrow \bar{X}$

cannot name $\bar{X}$ w/ fewer than $H_\delta(\bar{X})$ bits
 (specify) (deterministically) to achieve error $\leq \delta$

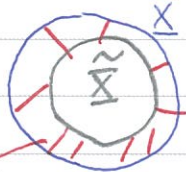
proof: (by contradiction)

assume ν & μ exist while $H_\delta(\bar{X}) > H_0(Y)$

$\bar{X} \rightarrow \nu(\bar{X}) \rightarrow \mu(\nu(\bar{X}))$ sufficiently correct

usually $(\geq 1 - \delta)$ $x \rightarrow \nu(x) \rightarrow \mu(\nu(x)) = x$

call such $x \in \tilde{\bar{X}}$ always correctly specified



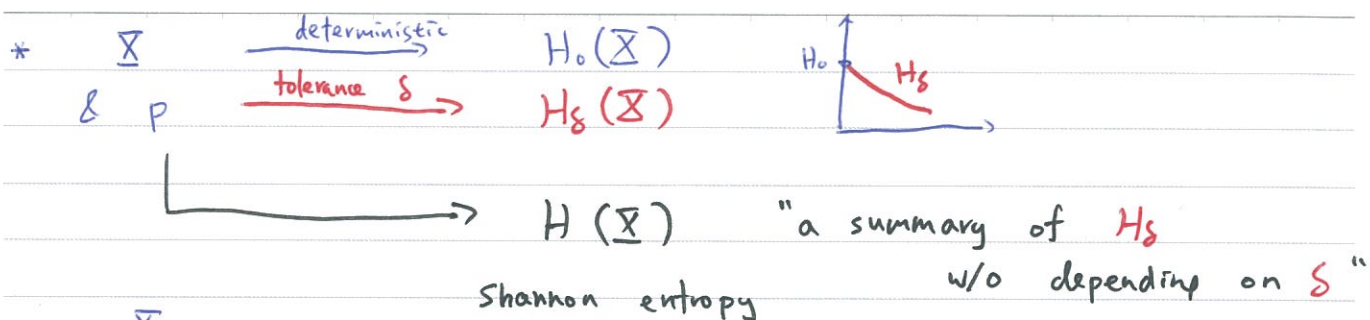
$\leq \delta$

by assumption

$$H_\delta(\bar{X}) \leq \log_2(\tilde{\bar{X}})$$

$|\tilde{\bar{X}}| \leq |Y|$ by deterministic pigeon hole

$$H_\delta(\bar{X}) \leq H_0(Y) \quad \nrightarrow$$



* $\underbrace{\mathcal{X}}_{a_1}$ $p(a_1) \rightarrow p_1$ as if needing $\log \frac{1}{p(a_1)}$ bits of information
 a_2 $p(a_2) \rightarrow p_2$ as if needing $\log \frac{1}{p(a_2)}$ bits of information
 $\vdots$ $\vdots$ $\vdots$ $\vdots$
 $\underbrace{a_k}$ $p(a_k) \rightarrow p_k$ $\log \frac{1}{p(a_k)}$

$$H(\mathcal{X}) = \underbrace{\sum p(a_k)}_{\text{self-avg.}} \underbrace{\log \frac{1}{p(a_k)}}_{\text{ind. info.}}$$

will take $0 \log \frac{1}{0} = 0$

* similar to H_0 , certainty = 0 information
 $H_0(\{a_i\}) = 0$
 $H(\mathcal{X}) = 0$ iff some $p_i = 1$

* similar to H_δ , $H(\mathcal{X}) \leq H_\delta(\mathcal{X})$

$$\begin{aligned} \sum p_k \log \frac{1}{p_k} - \log k &= \frac{1}{\ln 2} \sum p_k \boxed{\ln \frac{1}{p_k k}} \\ &\leq \frac{1}{\ln 2} \sum p_k \left(\frac{1}{p_k k} - 1 \right) \\ &= 0 \end{aligned}$$

) "=" if $\frac{1}{p_k k} = 1$

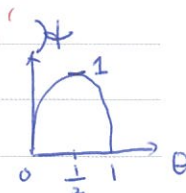
* binary ensemble $\mathcal{X} = \{0, 1\}$, $p_0 = \theta$, $p_1 = (1-\theta)$

$$H(\mathcal{X}) = \theta \log \frac{1}{\theta} + (1-\theta) \log \frac{1}{1-\theta} = \mathcal{H}(\theta)$$

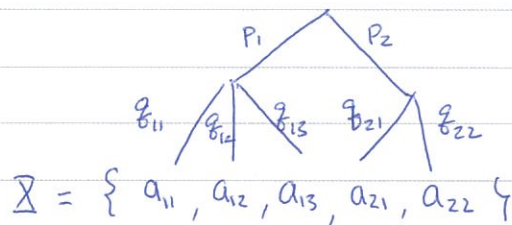
$$\mathcal{H}(0) = 0 + 0 = 0$$

$$\mathcal{H}(1) = 0$$

$$\max_{\theta} \mathcal{H}(\theta) = \text{"find } \mathcal{H}'(\theta) = 0 \text{"} = \mathcal{H}(0.5) = 1 = H_0(\mathcal{X})$$



* additivity of H



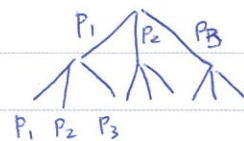
$$\begin{aligned}
 H(\mathcal{X}) &= \sum_i P_i \sum_j q_{ij} \log(P_i q_{ij}) \\
 &= \sum_i P_i \sum_j q_{ij} \log q_{ij} + \sum_i P_i \sum_j q_{ij} \log P_i \\
 &= \sum_i P_i \underbrace{H(\mathcal{X}_i)}_{\text{sub-level choices}} + \underbrace{H(\mathcal{X}_0)}_{\text{top-level choice}}
 \end{aligned}$$

* multiple independent experiments

$$H(\mathcal{X} \cdot \mathcal{X})$$

space: $\{(a_1, a_1), (a_1, a_2), \dots, (a_k, a_k)\}$

prob: $P_1 \cdot P_1 \quad P_1 \cdot P_2 \quad P_k \cdot P_k$



$$= H(\mathcal{X}) + H(\mathcal{X})$$

$$\Rightarrow H(\mathcal{X}^N) = N H(\mathcal{X})$$