DSA Homework #2 Reference Solution

April 11, 2014

2.1

	Algorithm 1 My algorithm		
(1)	1:	function BinSearch(recordarrayarr, integerlen, integervalue)	
	2:	$left \leftarrow 0, right \leftarrow len$	
	3:	$ans \leftarrow -1$	
	4:	while $left < right$ do	
	5:	$mid \leftarrow \lfloor \frac{left+right}{2} \rfloor$	
	6:	$\mathbf{if} \ arr[mid] = value \ \mathbf{then}$	
	7:	$ans \leftarrow mid, left \leftarrow mid + 1$	
	8:	else if $arr[mid] > value$ then	
	9:	$right \leftarrow mid$	
	10:	else if $arr[mid] < value$ then	
	11:	$left \leftarrow mid + 1$	
	12:	if $ans = -1$ then return NOTFOUND	
	13:	return arr[ans]	

(2) They will appear in the same relative order, because that the new element will always be sorted to the end of the elements with the same value due to the while condition A[j] > cur in the code.

2.2

- (1) For $n \in \mathbb{N}$, $f(n) = \lceil g(n) \rceil \leq g(n) + 1 \leq g(n) + g(n)/\epsilon = (1 + 1/\epsilon)g(n)$ since $g(n) \geq \epsilon$.
- (2) There exist $n_1, n_2 \in \mathbb{N}$ and $c_1, c_2 > 0$ such that for i = 1, 2 and $n \ge n_i$, $f_i(n) \le c_i g_i(n)$. Then, for $n \ge \max\{n_1, n_2\}, f_1(n) f_2(n) \le c_1 c_2 g_1(n) g_2(n)$.

- (3) Let $n_0 = 1$ and c = c'. Then, $f(n) \le cg(n)$ for all $n \ge n_0$.
- (4) There exist $n_0 \in \mathbb{N}$ and c > 0 such that $f(n) \leq cg(n)$ for $n \geq n_0$. Let $c' = \max_{n < =n_0} \{f(n)/g(n), c\}$. Then, if $n < n_0$, we have $f(n) = f(n)/g(n) \cdot g(n) \leq c'g(n)$, and if $n \geq n_0$, $f(n) \leq cg(n) \leq c'g(n)$.
- (5) The Horner's method requires n times of addition and n times of multiplication, so the total number of arithmetic operations this method executes is 2n = O(n).
- (6) If $f(n) = n \mod 2$ and $g(n) = 1 (n \mod 2)$, then neither f(n) = O(g(n)) or g(n) = O(f(n)) is true.
- (7) Let f(n) = 1/n and suppose that $f(n) = O(f(n)^2)$. Then, there exists c > 0 such that $f(n) \le cf(n)^2$ for all $n \in \mathbb{N}$, i.e., $n \le c$ for all n, a contradiction.