

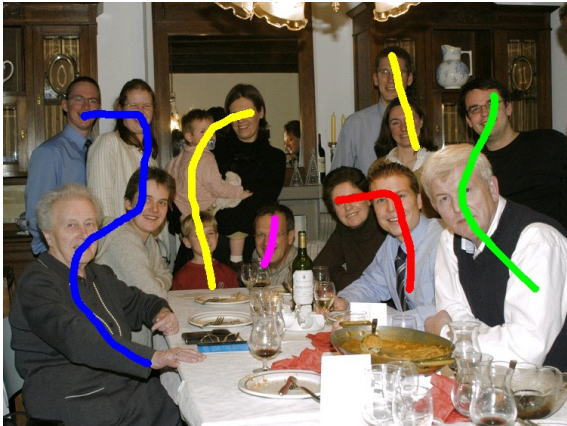
Graph Cut

Digital Visual Effects

Yung-Yu Chuang

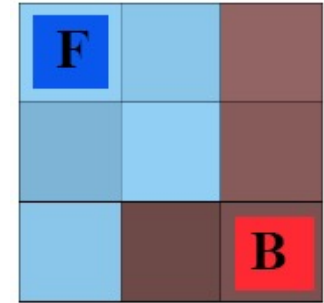
with slides by Fredo Durand, Ramesh Raskar

Graph cut



Graph cut

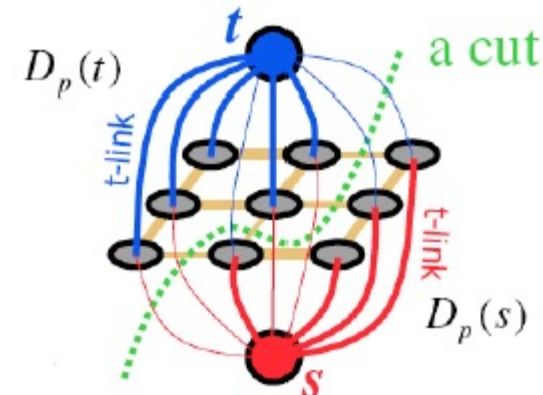
- Interactive image segmentation using graph cut
- Binary label: foreground vs. background
- User labels some pixels
 - similar to trimap, usually sparser
- Exploit
 - Statistics of known Fg & Bg
 - Smoothness of label
- Turn into discrete graph optimization
 - Graph cut (min cut / max flow)



F F B

F F B

F B B



Energy function

- Labeling: one value per pixel, F or B
- Energy(labeling) = data + smoothness

F		
		B

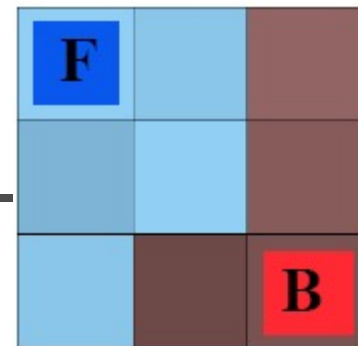
F	B	B
F	B	B
F	B	B

**One labeling
(ok, not best)**

- Very general situation
 - Will be minimized
- Data: for each pixel
 - Probability that this color belongs to F (resp. B)
 - Similar in spirit to Bayesian matting
- Smoothness (aka regularization): per neighboring pixel pair
 - Penalty for having different label
 - Penalty is downweighted if the two pixel colors are very different
 - Similar in spirit to bilateral filter

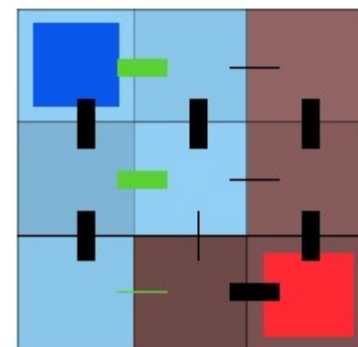
Data term

- A.k.a regional term
(because integrated over full region)
- $D(L) = \sum_i -\log h[L_i](C_i)$
- Where i is a pixel
 L_i is the label at i (F or B),
 C_i is the pixel value
 $h[L_i]$ is the histogram of the observed Fg
(resp Bg)
- Note the minus sign



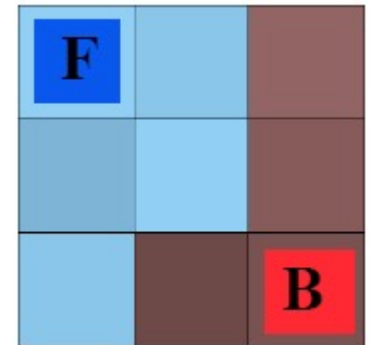
F	B	B
F	B	B
F	B	B

F		B
F	B	B
F	B	B



Hard constraints

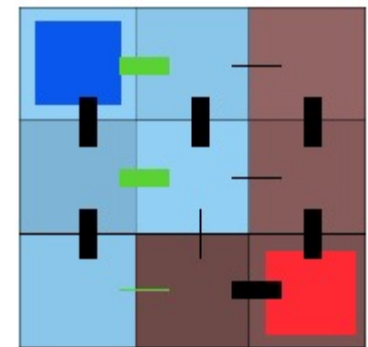
- The user has provided some labels
- The quick and dirty way to include constraints into optimization is to replace the data term by a huge penalty if not respected.
- $D(L_i)=0$ if respected
- $D(L_i)=K$ if not respected
 - e.g. $K=-\text{\#pixels}$



Smoothness term

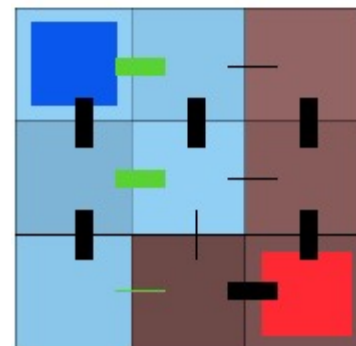
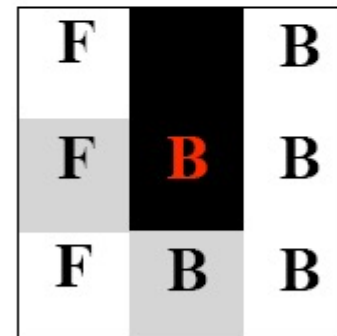
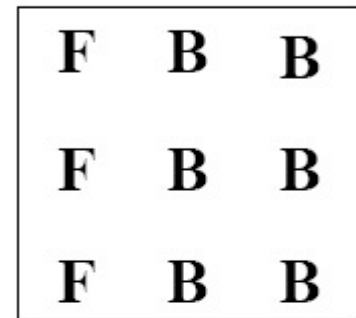
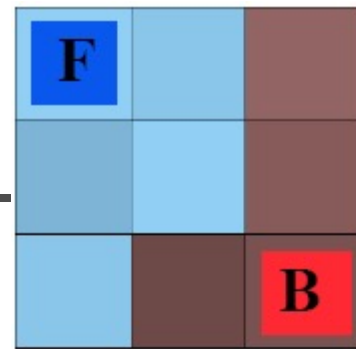
- a.k.a boundary term, a.k.a. regularization
- $S(L) = \sum_{\{j, i\} \in N} B(C_i, C_j) \delta(L_i - L_j)$
- Where i, j are neighbors
 - e.g. 8-neighborhood
(but I show 4 for simplicity)
- $\delta(L_i - L_j)$ is 0 if $L_i = L_j$, 1 otherwise
- $B(C_i, C_j)$ is high when C_i and C_j are similar, low if there is a discontinuity between those two pixels
 - e.g. $\exp(-||C_i - C_j||^2 / 2\sigma^2)$
 - where σ can be a constant or the local variance
- Note positive sign

F	B	B
F	B	B
F	B	B



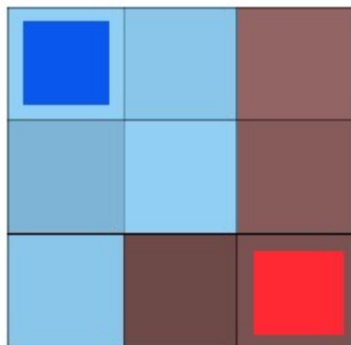
Optimization

- $E(L) = D(L) + \lambda S(L)$
- λ is a black-magic constant
- Find the labeling that minimizes E
- In this case, how many possibilities?
 - 2^9 (512)
 - We can try them all!
 - What about megapixel images?

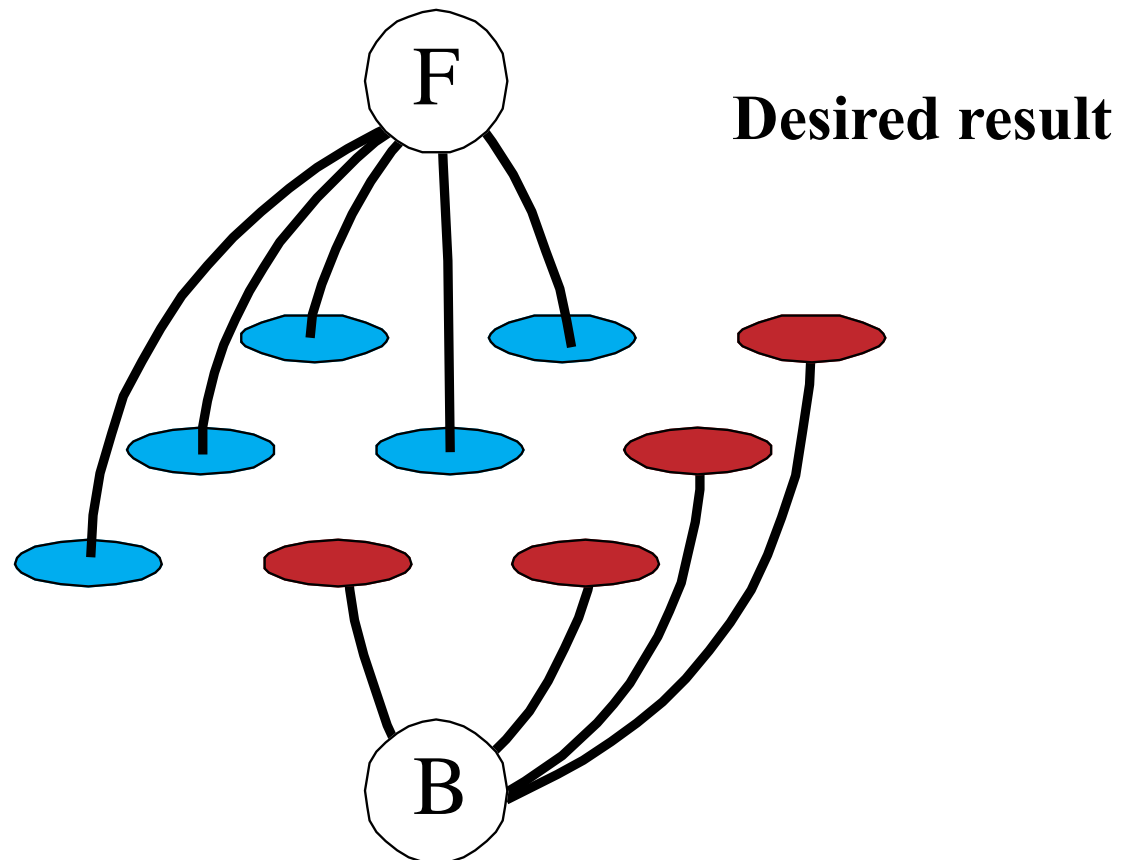


Labeling as a graph problem

- Each pixel = node
- Add two nodes F & B
- Labeling: link each pixel to either F or B

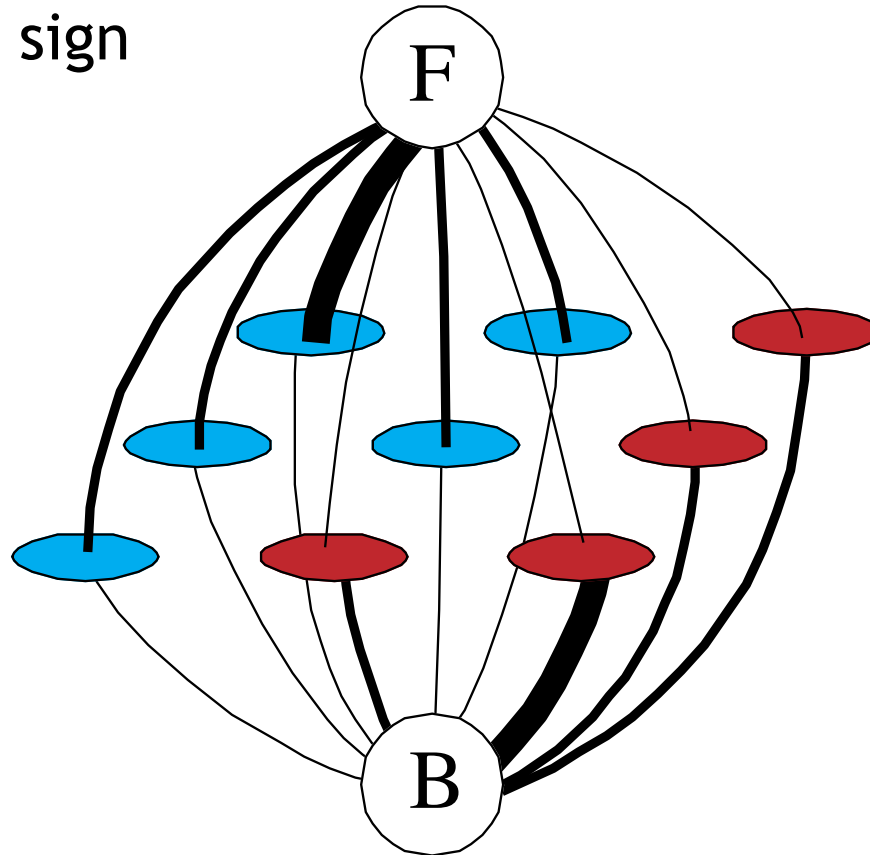


F	F	B
F	F	B
F	B	B



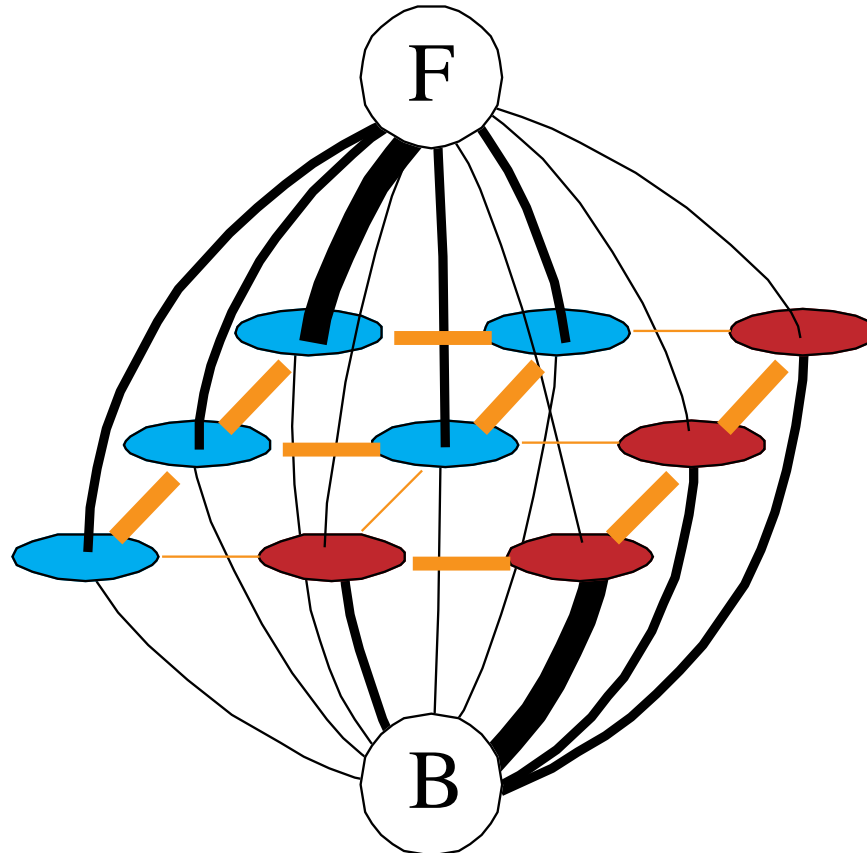
Data term

- Put one edge between each pixel and F & G
- Weight of edge = minus data term
 - Don't forget huge weight for hard constraints
 - Careful with sign



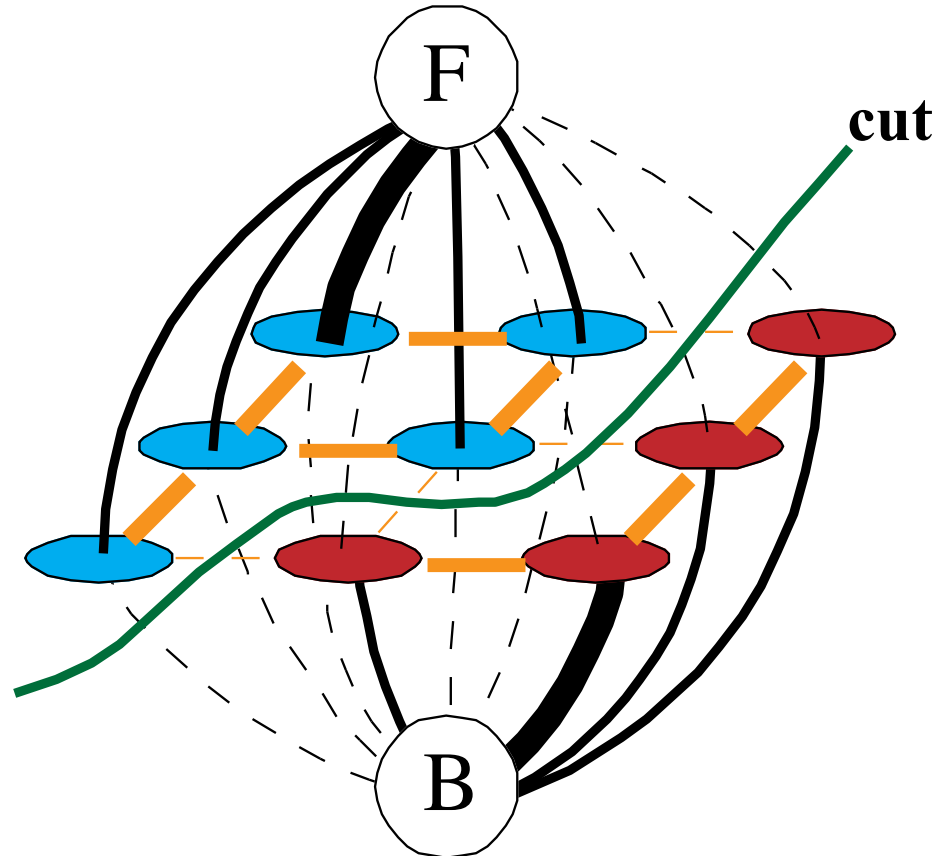
Smoothness term

- Add an edge between each neighbor pair
- Weight = smoothness term



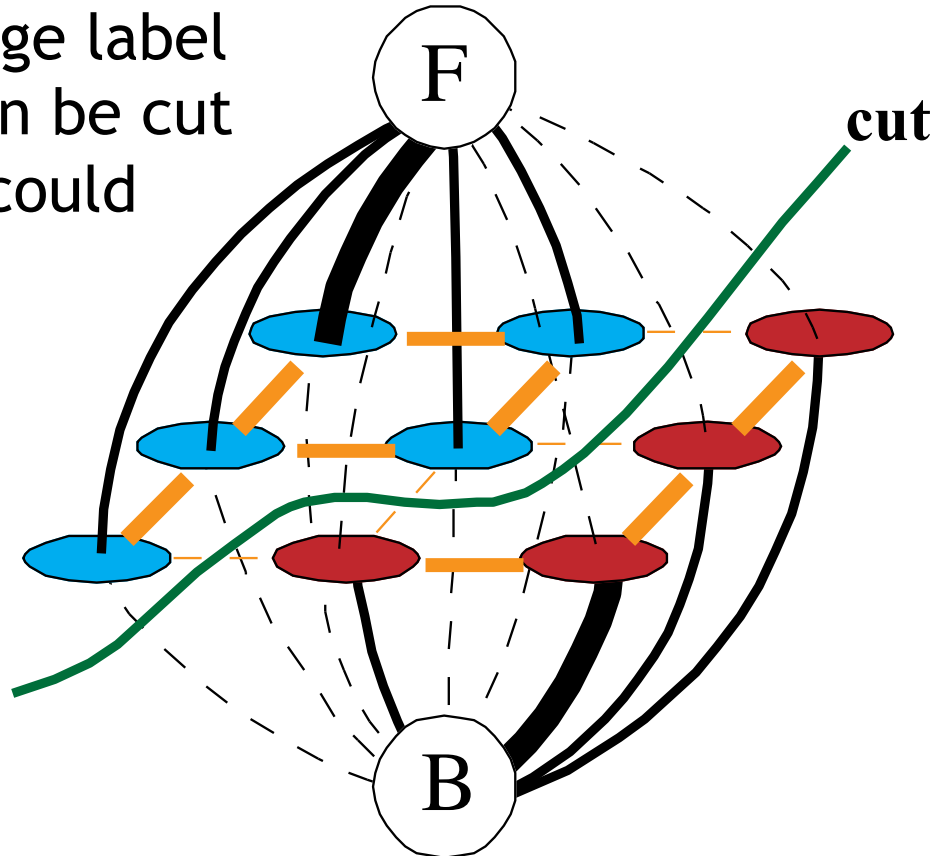
Min cut

- Energy optimization equivalent to min cut
- Cut: remove edges to disconnect F from B
- Minimum: minimize sum of cut edge weight

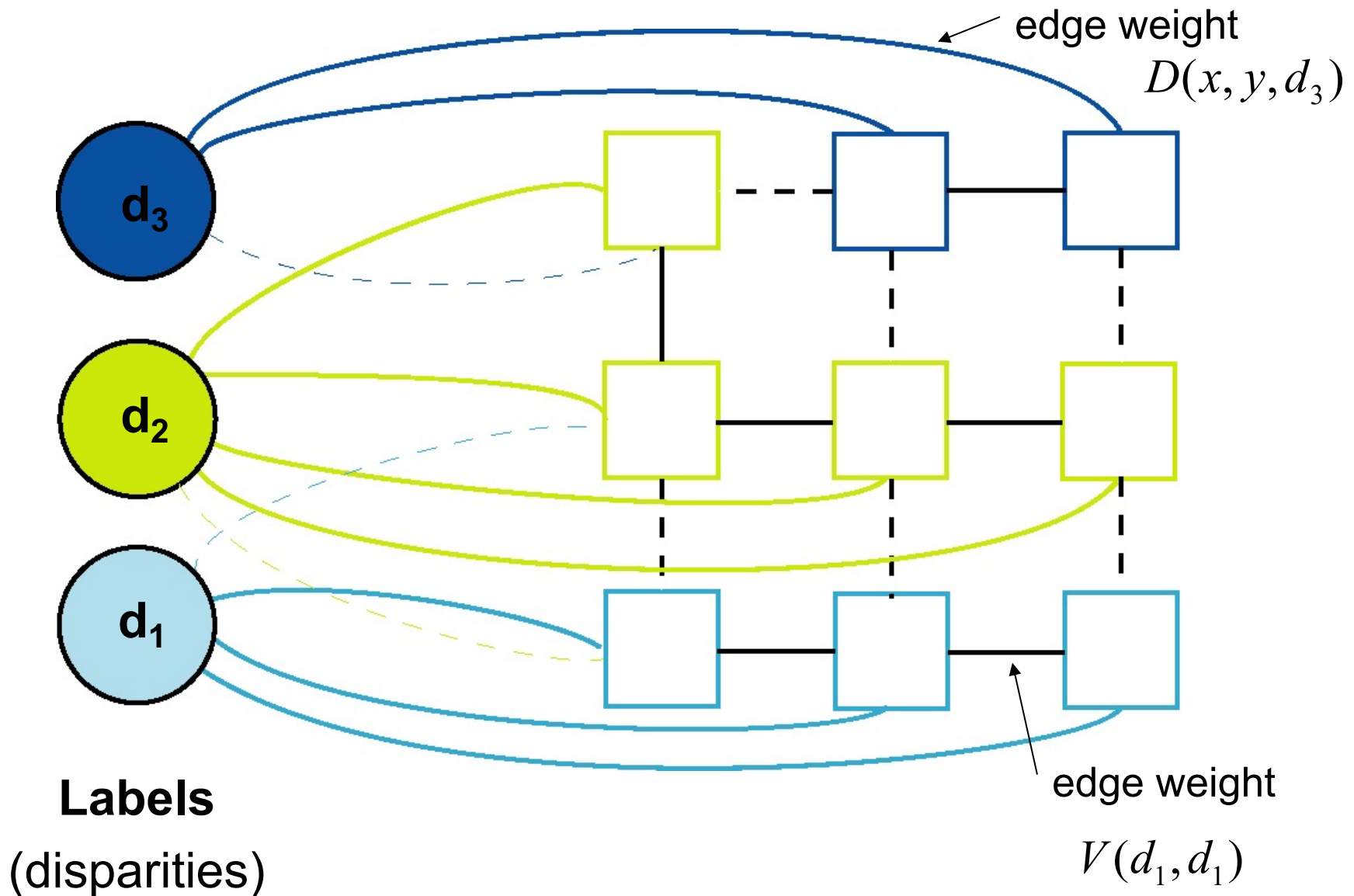


Min cut \Leftrightarrow labeling

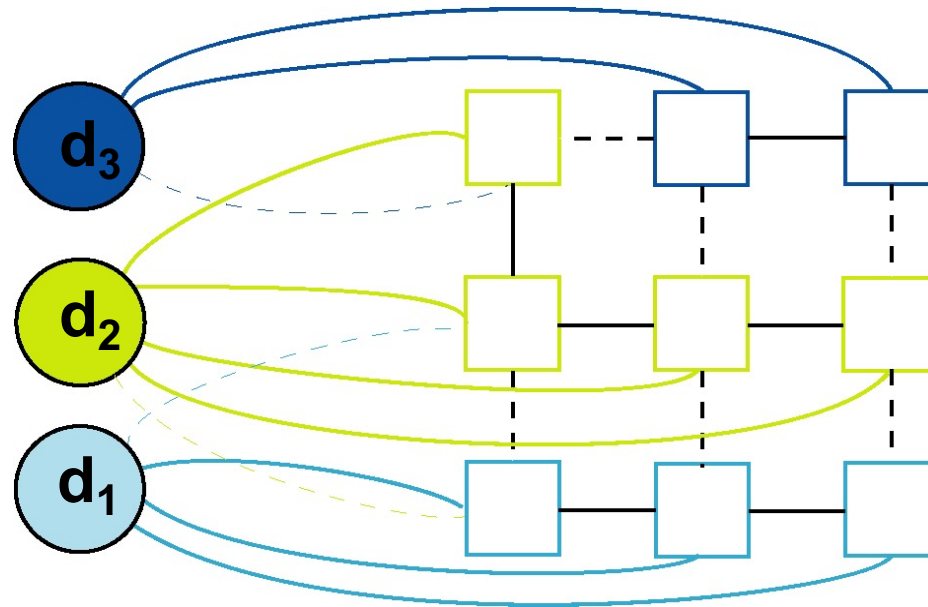
- In order to be a cut:
 - For each pixel, either the F or G edge has to be cut
- In order to be minimal
 - Only one edge label per pixel can be cut (otherwise could be added)



Energy minimization via graph cuts

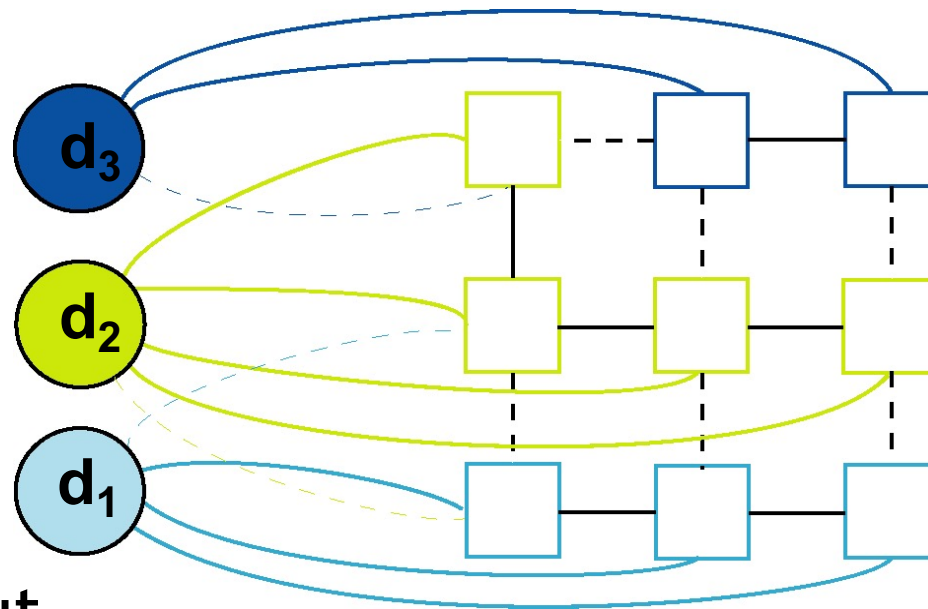


Energy minimization via graph cuts



- Graph Cost
 - Matching cost between images
 - Neighborhood matching term
 - Goal: figure out which labels are connected to which pixels

Energy minimization via graph cuts



- Graph Cut

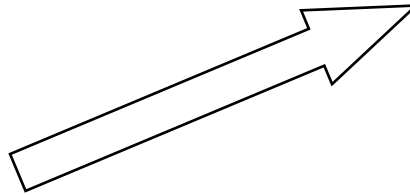
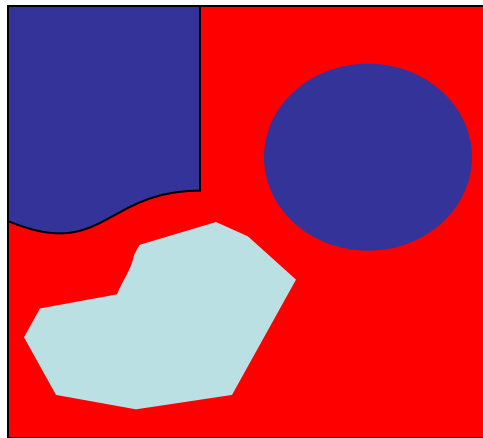
- Delete enough edges so that
 - each pixel is (transitively) connected to exactly one label node
- Cost of a cut: sum of deleted edge weights
- Finding min cost cut equivalent to finding global minimum of energy function

Computing a multiway cut

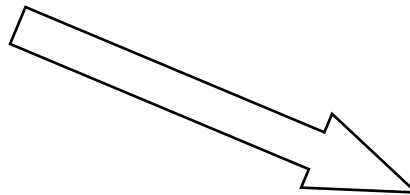
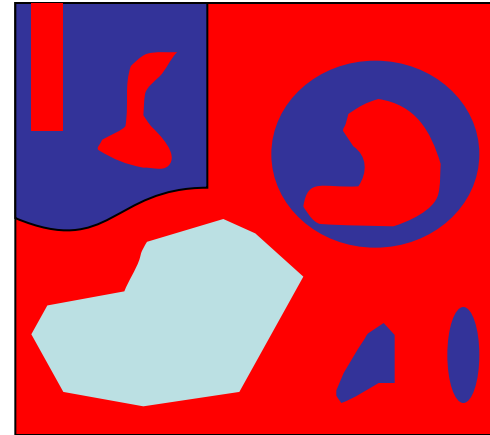
- With 2 labels: classical min-cut problem
 - Solvable by standard flow algorithms
 - polynomial time in theory, nearly linear in practice
 - More than 2 terminals: NP-hard
 - [Dahlhaus *et al.*, STOC '92]
- Efficient approximation algorithms exist
 - Within a factor of 2 of optimal
 - Computes local minimum in a strong sense
 - even very large moves will not improve the energy
 - Yuri Boykov, Olga Veksler and Ramin Zabih, [Fast Approximate Energy Minimization via Graph Cuts](#), International Conference on Computer Vision, September 1999.

Move examples

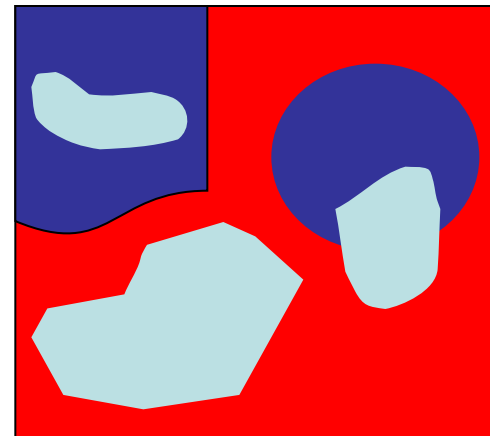
Starting point



Red-blue swap move



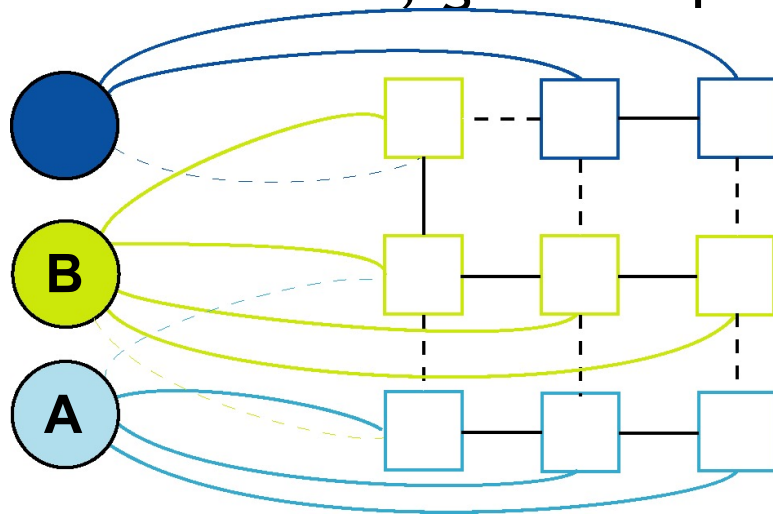
Green expansion move



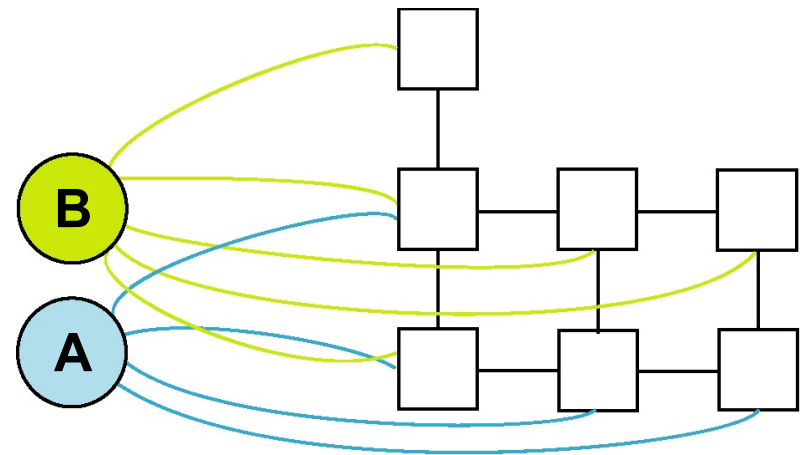
The swap move algorithm

1. Start with an arbitrary labeling
2. Cycle through every label pair (A,B) in some order
 - 2.1 Find the lowest E labeling within a single AB -swap
 - 2.2 Go there if E is lower than the current labeling
3. If E did not decrease in the cycle, we're done

Otherwise, go to step 2



Original graph



AB subgraph
(run min-cut on this graph)

The expansion move algorithm

1. Start with an arbitrary labeling
2. Cycle through every label A in some order
 - 2.1 Find the lowest E labeling within a single A -expansion
 - 2.2 Go there if it E is lower than the current labeling
3. If E did not decrease in the cycle, we're done
Otherwise, go to step 2



GrabCut

Interactive Foreground Extraction using Iterated Graph Cuts

Carsten Rother
Vladimir Kolmogorov
Andrew Blake

Microsoft Research Cambridge-UK



GrabCut

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Interactive Digital Photomontage



- Combining multiple photos
- Find seams using graph cuts
- Combine gradients and integrate















set of originals

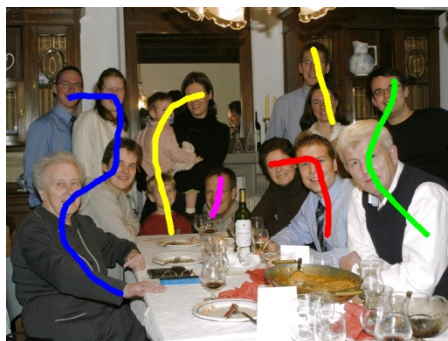


photo montage

Source images



Brush strokes



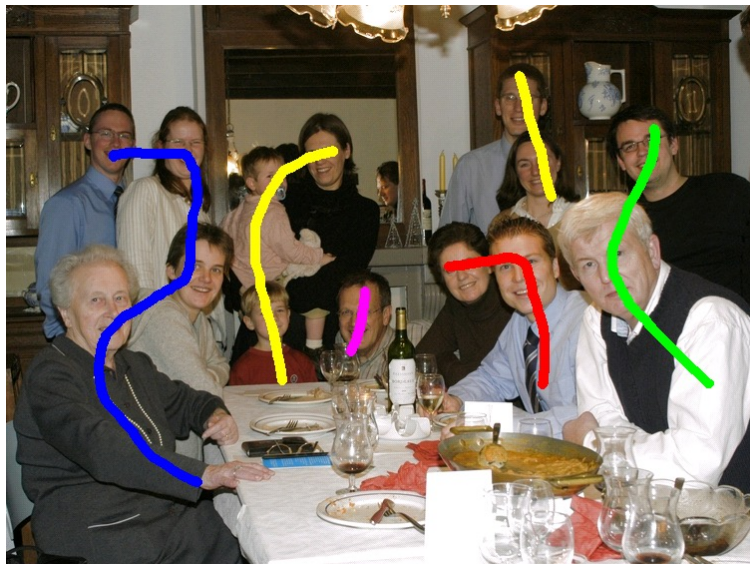
Computed labeling



Composite



Brush strokes

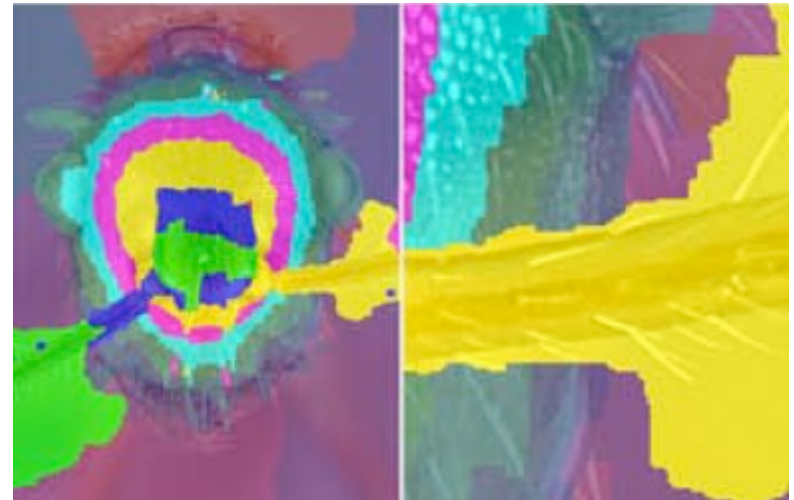


Computed labeling



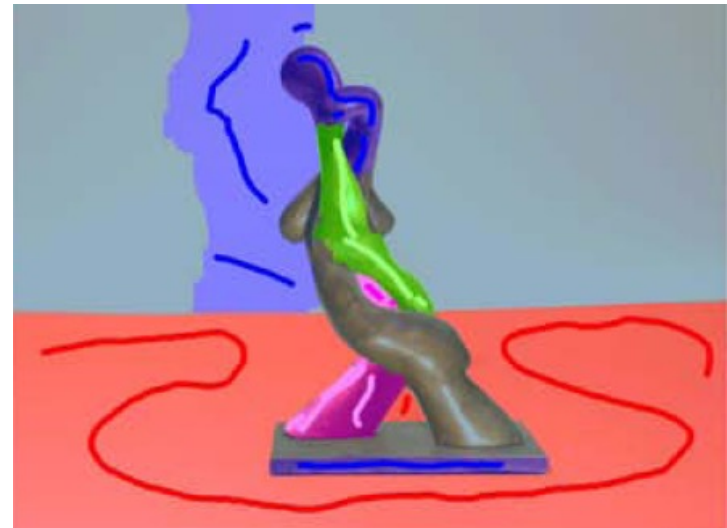
Interactive Digital Photomontage

- Extended depth of field



Interactive Digital Photomontage

- Relighting



Interactive Digital Photomontage



Interactive Digital Photomontage

Aseem Agarwala, Mira Dontcheva
Maneesh Agrawala, Steven Drucker, Alex Colburn
Brian Curless, David Salesin, Michael Cohen

