

Homework 2

October 20, 2004

Due date: November 3, 2004

1. (15%) [Sipser 2.2] Given the languages $A = \{a^m b^n c^n \mid m, n \geq 0\}$ and $C = \{a^n b^n c^m \mid m, n \geq 0\}$,
 - a. Design pushdown automata to recognize each of the languages A and C .
 - b. Use the languages A and C together with Example 2.20 (in which, we prove that $B = \{a^n b^n c^n \mid n \geq 0\}$ is not context free) to show that the class of context-free languages is not closed under intersection.
 - c. Use part (b) and DeMorgan's law ($\overline{A \cup B} = \overline{A} \cap \overline{B}$) to show that the class of context-free languages is not closed under complementation.
2. (15%) [Sipser 2.4] Give context-free grammars that generate the following languages over $\{0, 1\}$.
 - a. $\{w \mid w \text{ contains at least three 1s}\}$
 - b. $\{w \mid \text{the length of } w \text{ is odd and its middle symbol is a } 0\}$
 - c. $\{w \mid w = w^R\}$
3. (10%) Convert the following CFG to an equivalent PDA,
 $A \rightarrow BAB \mid B \mid \epsilon$
 $B \rightarrow 00 \mid \epsilon$
4. (10%) [Sipser 2.18b] Using the pumping lemma to show that the language over $\{0, \#\}$, $\{0^n \# 0^{2n} \# 0^{3n} \mid n \geq 0\}$, is not context free.
5. (12%) [Sipser 3.5] Examine the formal definition of a Turing machine to answer the following questions, and explain your reasoning.
 - a. Can a Turing machine ever write the blank symbol \sqcup on its tape?
 - b. Can the tape alphabet Γ be the same as the input alphabet Σ ?
 - c. Can a Turing machine's head *ever* be in the same location in two successive steps?
 - d. Can a Turing machine contain just a single state?
6. (20%) [Sipser 3.2] Given TM M_1 whose description and state diagram appear in Example 3.5 of the textbook (the Turing machine shown in the class to recognize the language $\{w \# w \mid w \in \{0, 1\}^*\}$), for each of the following inputs, show the sequence of configurations that M_1 enters.
 - a. 11
 - b. 1#1
 - c. 1##1
 - d. 10#11
 - e. 10#10
7. (18%) [Sipser 3.8] Give implementation-level descriptions of Turing Machines that decide the following languages over the alphabet $\{0, 1\}$:
 - a. $\{w \mid w \text{ contains an equal number of 0s and 1s}\}$
 - b. $\{w \mid w \text{ contains twice as many 0s as 1s}\}$ (that is, (the number of 0s)=2*(the number of 1s))
 - c. $\{w \mid w \text{ does not contain twice as many 0s as 1s}\}$