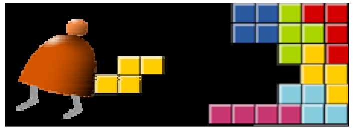
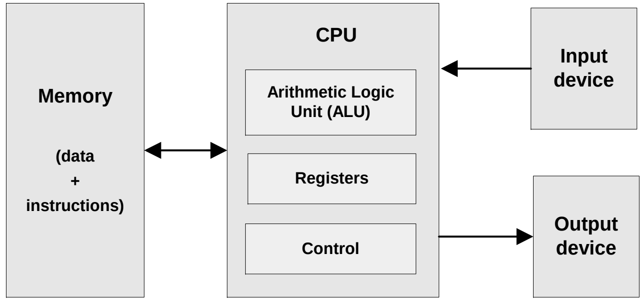
Computer Architecture

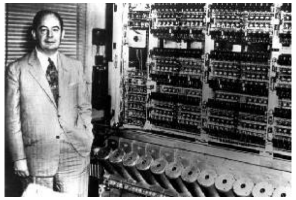


Building a Modern Computer From First Principles

www.nand2tetris.org

Von Neumann machine (circa 1940)



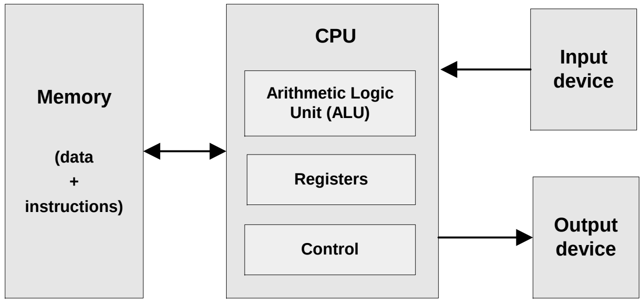


Stored program concept!

John Von Neumann (and others) ... made it possible

Andy Grove (and others) ... made it small and fast.

Processing logic: fetch-execute cycle



Executing the *current instruction* involves one or more of the following micro-tasks:

- Have the ALU compute some function out = *f* (register values)
- Write the ALU output to selected registers
- As a side-effect of this computation, figure out which instruction to fetch and execute next.

The Hack chip-set and hardware platform

Elementary logic gates

- `Nand`
- `Not`
- `And`
- `Or`
- `Xor`
- `Mux`
- `Dmux`
- `Not16`
- `And16`
- `Or16`
- `Mux16`
- `Or8Way`
- `Mux4Way16`
- `Mux8Way16`
- `DMux4Way`
- `DMux8Way`

done

Combinational chips

- `HalfAdder`
- `FullAdder`
- `Add16`
- `Inc16`
- `ALU`

done

Sequential chips

- `DFF`
- `Bit`
- `Register`
- `RAM8`
- `RAM64`
- `RAM512`
- `RAM4K`
- `RAM16K`
- `PC`

done

Computer Architecture

- `Memory`
- `CPU`
- `Computer`

this lecture

The Hack computer

- A 16-bit Von Neumann platform
- The *instruction memory* and the *data memory* are physically separate
- Screen: 512 rows by 256 columns, black and white
- Keyboard: standard
- Designed to execute programs written in the Hack machine language
- Can be easily built from the chip-set that we built so far in the course

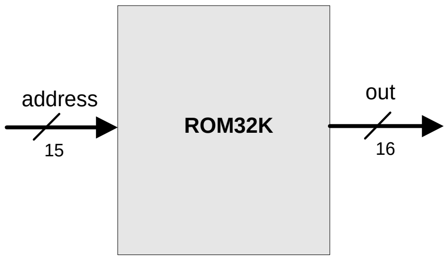
Main parts of the Hack computer:

- Instruction memory (ROM)
- Memory (RAM):
  - Data memory
  - Screen (memory map)
  - Keyboard (memory map)
- CPU
- Computer (the logic that holds everything together).

Lecture / construction plan

- Instruction memory
- Memory:
  - Data memory
  - Screen
  - Keyboard
- CPU
- Computer

Instruction memory



Function:

- The ROM is pre-loaded with a program written in the Hack machine language
- The ROM chip always emits a 16-bit number:

  `out = ROM32K[address]`

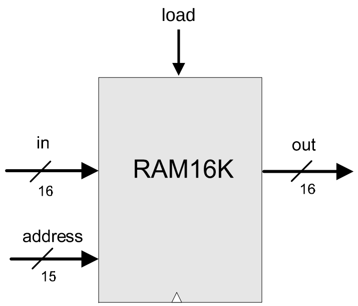
- This number is interpreted as the *current instruction.*

Data memory

Low-level (hardware) read/write logic:

To read RAM[k]: set address to k,
probe out

To write RAM[k]=x: set address to k,
set in to x,
set load to 1,
run the clock



High-level (OS) read/write logic:

To read RAM[k]: use the OS command out = peek(k)

To write RAM[k]=x: use the OS command poke(k,x)

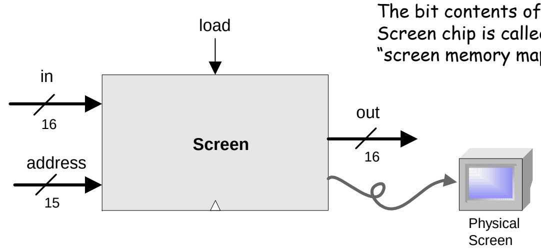
peek and poke are OS commands whose implementation should effect the same behavior as the low-level commands

More about peek and poke this later in the course, when we'll write the OS.

Lecture / construction plan

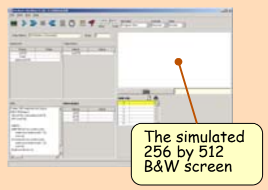
- Instruction memory ✓
- Memory:
  - Data memory ✓
  - Screen
  - Keyboard
- CPU
- Computer

Screen



The bit contents of the Screen chip is called the "screen memory map"

Simulated screen:



When loaded into the hardware simulator, the built-in `Screen.hdl` chip opens up a screen window; the simulator then refreshes this window from the screen memory map several times each second.

The Screen chip has a basic RAM chip functionality:

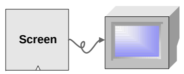
- read logic: `out = Screen[address]`
- write logic: `if load then Screen[address] = in`

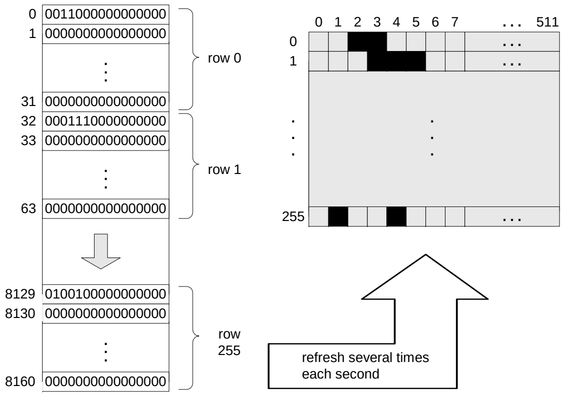
Side effect:

Continuously refreshes a 256 by 512 black-and-white screen device

Screen memory map

In the Hack platform, the screen is implemented as an 8K 16-bit RAM chip.

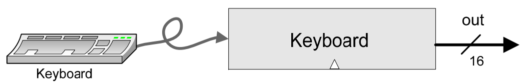




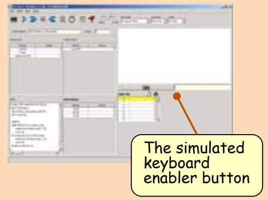
How to set the (`row,col`) pixel of the screen to black or to white:

- Low-level (machine language): Set the `col%16` bit of the word found at `Screen[row*32+col/16]` to `1` or to `0` (`col/16` is integer division)
- High-level: Use the OS command `drawPixel(row,col)` (effects the same operation, discussed later in the course, when we'll write the OS).

Keyboard



Simulated keyboard:



The keyboard is implemented as a built-in **`Keyboard.hdl`** chip. When this java chip is loaded into the simulator, it connects to the regular keyboard and pipes the scan-code of the currently pressed key to the keyboard memory map.

Keyboard chip: a single 16-bit register

Input: scan-code (16-bit value) of the currently pressed key, or 0 if no key is pressed

Output: same

Special keys:

Key pressed	Keyboard output	Key pressed	Keyboard output
newline	128	end	135
backspace	129	page up	136
left arrow	130	page down	137
up arrow	131	insert	138
right arrow	132	delete	139
down arrow	133	esc	140
home	134	f1-f12	141-152

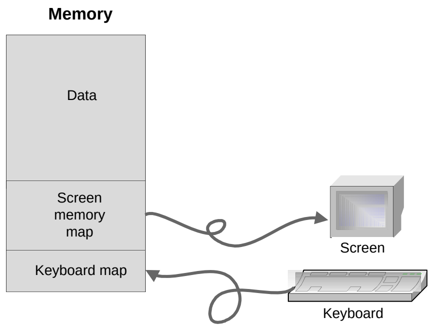
How to read the keyboard:

- Low-level (hardware): probe the contents of the `Keyboard` chip
- High-level: use the OS command `keyPressed()` (effects the same operation, discussed later in the course, when we'll write the OS).

Lecture / construction plan

- Instruction memory
- Memory:
  - Data memory
  - Screen
  - Keyboard
- CPU
- Computer

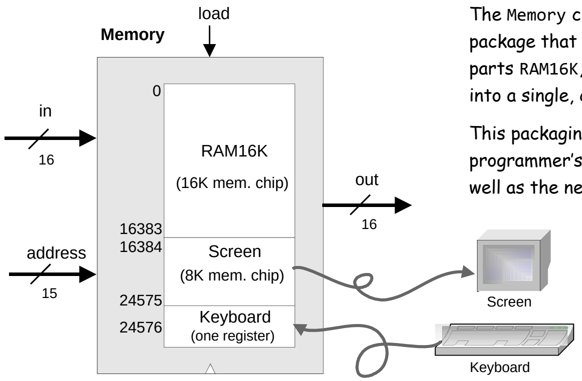
Memory: conceptual / programmer's view



Using the memory:

- To record or recall values (e.g. variables, objects, arrays), use the first 16K words of the memory
- To write to the screen (or read the screen), use the next 8K words of the memory
- To read which key is currently pressed, use the next word of the memory.

Memory: physical implementation



The Memory chip is essentially a package that integrates the three chip-parts RAM16K, Screen, and Keyboard into a single, contiguous address space.

This packaging effects the programmer's view of the memory, as well as the necessary I/O side-effects.

Access logic:

- Access to any address from 0 to 16,383 results in accessing the RAM16K chip-part
- Access to any address from 16,384 to 24,575 results in accessing the Screen chip-part
- Access to address 24,576 results in accessing the keyboard chip-part
- Access to any other address is invalid.

Lecture / construction plan

- Instruction memory
- Memory:
  - Data memory
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  - Keyboard
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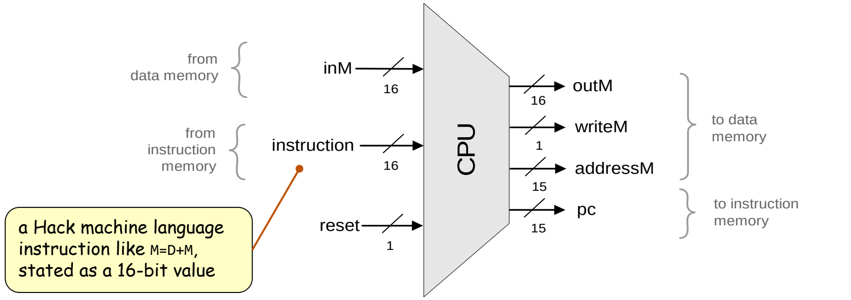
- Computer

CPU



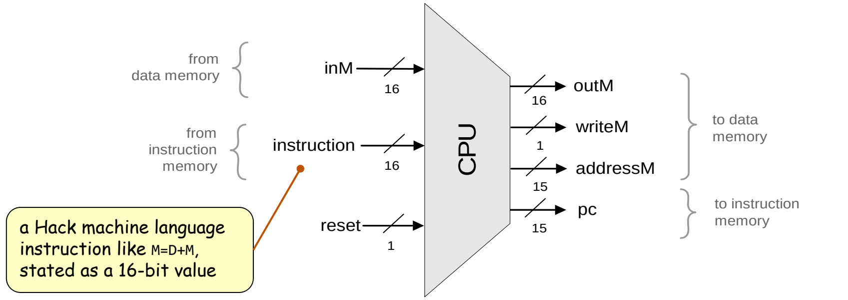
CPU internal components (invisible in this chip diagram): ALU and 3 registers: A, D, PC

CPU execute logic:

The CPU executes the `instruction` according to the Hack language specification:

- The `D` and `A` values, if they appear in the instruction, are read from (or written to) the respective CPU-resident registers
- The `M` value, if there is one in the instruction's RHS, is read from `inM`
- If the instruction's LHS includes `M`, then the ALU output is placed in `outM`, the value of the CPU-resident `A` register is placed in `addressM`, and `writeM` is asserted.

CPU



CPU internal components (invisible in this chip diagram): ALU and 3 registers: A, D, PC

CPU fetch logic:

Recall that:

1. the `instruction` may include a jump directive (expressed as non-zero jump bits)
2. the `ALU` emits two control bits, indicating if the `ALU` output is zero or less than zero

If `reset==0`: the CPU uses this information (the jump bits and the ALU control bits) as follows:
If there should be a jump, the `PC` is set to the value of `A`; else, `PC` is set to `PC+1`

If `reset==1`: the `PC` is set to `0`. (restarting the computer)

The *C*-instruction revisited

dest = comp; jump — *comp*, *dest*, *jump*

binary: `1 1 1 a c1 c2 c3 c4 c5 c6 d1 d2 d3 j1 j2 j3`

(when a=0) *comp*	c1	c2	c3	c4	c5	c6	(when a=1) *comp*
0	1	0	1	0	1	0	
1	1	1	1	1	1	1	
-1	1	1	1	0	1	0	
D	0	0	1	1	0	0	
A	1	1	0	0	0	0	M
!D	0	0	1	1	0	1	
!A	1	1	0	0	0	1	!M
-D	0	0	1	1	1	1	
-A	1	1	0	0	1	1	-M
D+1	0	1	1	1	1	1	
A+1	1	1	0	1	1	1	M+1
D-1	0	0	1	1	1	0	
A-1	1	1	0	0	1	0	M-1
D+A	0	0	0	0	1	0	D+M
D-A	0	1	0	0	1	1	D-M
A-D	0	0	0	1	1	1	M-D
D&A	0	0	0	0	0	0	D&M
D|A	0	1	0	1	0	1	D|M

d1	d2	d3	*Mnemonic*	*Destination (where to store the computed value)*
0	0	0	null	The value is not stored anywhere
0	0	1	M	Memory[A] (memory register addressed by A)
0	1	0	D	D register
0	1	1	MD	Memory[A] and D register
1	0	0	A	A register
1	0	1	AM	A register and Memory[A]
1	1	0	AD	A register and D register
1	1	1	AMD	A register, Memory[A], and D register

j1 ($out < 0$)	j2 ($out = 0$)	j3 ($out > 0$)	Mnemonic	Effect
0	0	0	null	No jump
0	0	1	JGT	If $out > 0$ jump
0	1	0	JEQ	If $out = 0$ jump
0	1	1	JGE	If $out \geq 0$ jump
1	0	0	JLT	If $out < 0$ jump
1	0	1	JNE	If $out \neq 0$ jump
1	1	0	JLE	If $out \leq 0$ jump
1	1	1	JMP	Jump

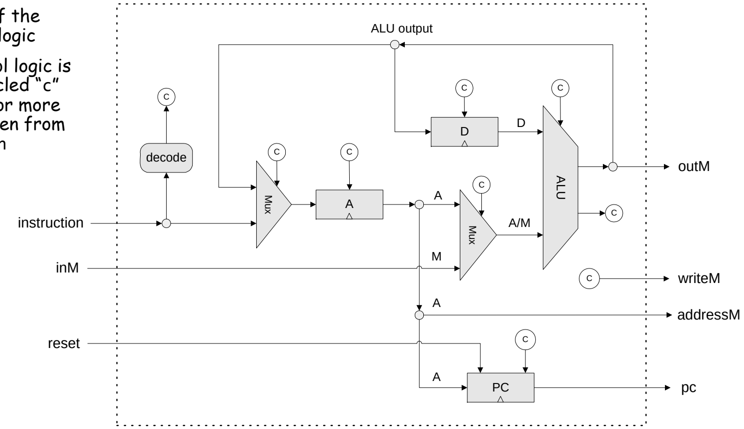
CPU implementation

dest = comp; jump — *comp*, *dest*, *jump*

binary: `1 1 1 a c1 c2 c3 c4 c5 c6 d1 d2 d3 j1 j2 j3`

Chip diagram:

- Includes most of the CPU's execution logic
- The CPU's control logic is hinted: each circled "c" represents one or more control bits, taken from the instruction
- The "decode" bar does not represent a chip, but rather indicates that the instruction bits are decoded somehow.



Cycle:

- Execute
- Fetch

Execute logic:

- Decode
- Execute

Fetch logic:

If there should be a jump, set `PC` to `A`
else set `PC` to `PC+1`

Resetting the computer:

Set `reset` to `1`, then set it to `0`.

Lecture / construction plan

- Instruction memory
- Memory:
  - Data memory
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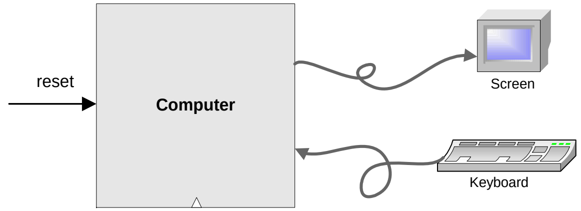
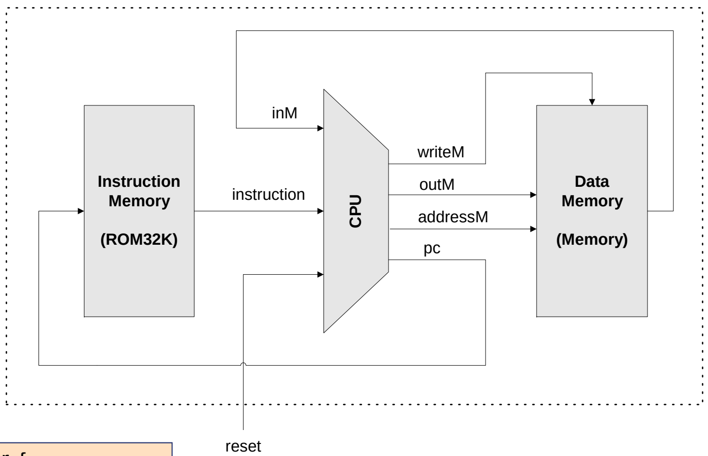
Computer-on-a-chip interface




```
Chip Name: Computer  // Topmost chip in the Hack platform
Input:     reset
Function:  When reset is 0, the program stored in the
           computer's ROM executes. When reset is 1, the
           execution of the program restarts. Thus, to start a
           program's execution, reset must be pushed "up" (1)
           and "down" (0).

           From this point onward the user is at the mercy of
           the software. In particular, depending on the
           program's code, the screen may show some output and
           the user may be able to interact with the computer
           via the keyboard.
```


Computer-on-a-chip implementation




```
CHIP Computer {
    IN reset;
    PARTS:
    // implementation missing
}
```

Implementation:

Simple, the chip-parts do all the hard work.

Perspective: from here to a "real" computer

- Caching
- More I/O units
- Special-purpose processors (I/O, graphics, communications, …)
- Multi-core / parallelism
- Efficiency
- Energy consumption considerations
- And more ...