

# Advanced Procedures

*Computer Organization and Assembly Languages*

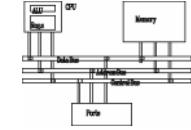
*Yung-Yu Chuang*

*2005/12/4*

*with slides by Kip Irvine*

# Overview

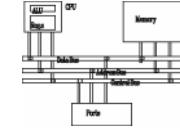
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- Stack Frames (a communication protocol between high-level-language procedures)
- Stack Parameters (passing by value, passing by reference, memory model and language specifiers)
- Local Variables (creating and initializing on the stack, scope and lifetime, LOCAL)
- Recursion
- Related directives: INVOKE, PROC, PROTO
- Creating Multimodule Programs

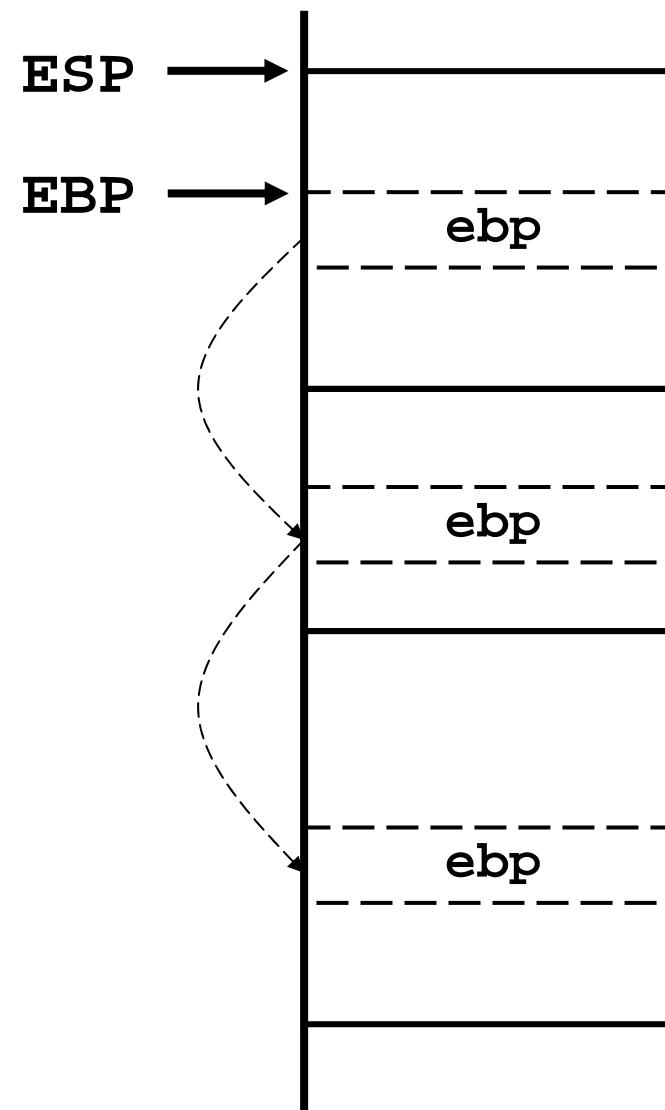
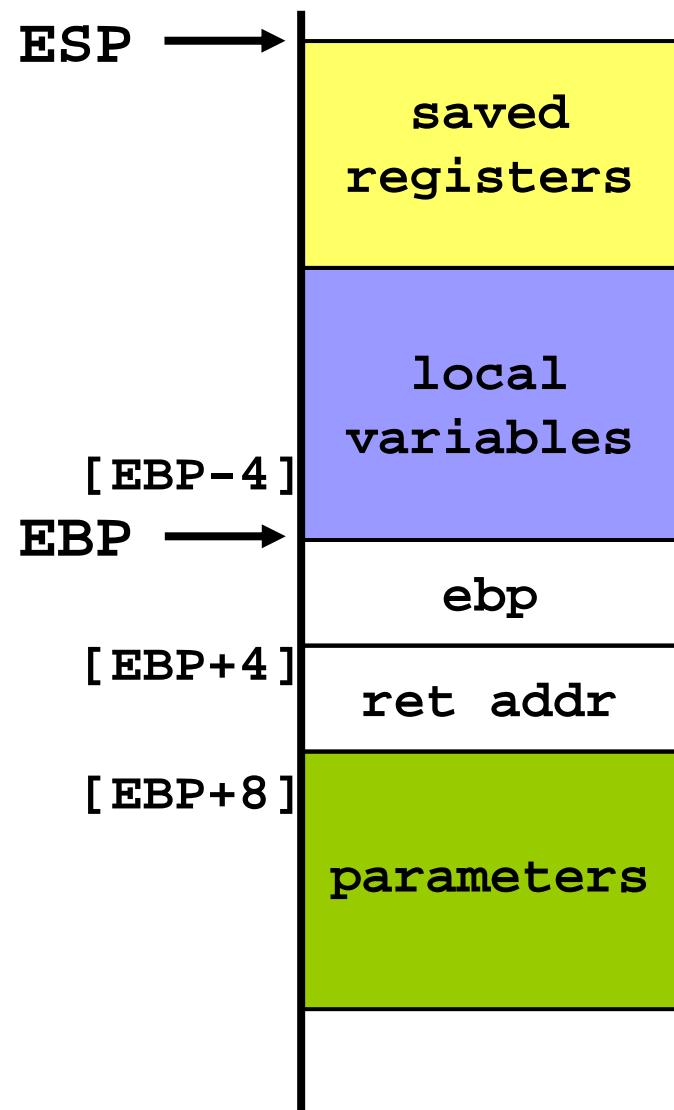
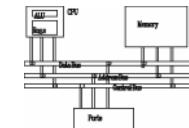
# Stack frame

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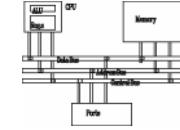
- Also known as an activation record
- Area of the stack set aside for a procedure's return address, passed parameters, saved registers, and local variables
- Created by the following steps:
  - Calling procedure pushes *arguments* on the stack and calls the procedure.
  - The subroutine is called, causing the *return address* to be pushed on the stack.
  - The called procedure pushes *EBP* on the stack, and *sets EBP to ESP*.
  - If *local variables* are needed, a constant is subtracted from *ESP* to make room on the stack.
  - The *registers needed to be saved* are pushed.

# Stack frame



# Explicit access to stack parameters

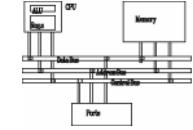
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- A procedure can explicitly access stack parameters using constant offsets from **EBP**.
  - Example: `[ebp + 8]`
- **EBP** is often called the base pointer or frame pointer because it holds the base address of the stack frame.
- **EBP** does not change value during the procedure.
- **EBP** must be restored to its original value when a procedure returns.

# Parameters

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- Two types: register parameters and stack parameters.
- Stack parameters are more convenient than register parameters.

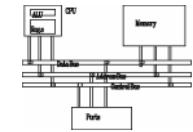
```
pushad  
mov esi,OFFSET array  
mov ecx,LENGTHOF array  
mov ebx,TYPE array  
call DumpMem  
popad
```

register parameters

```
push TYPE array  
push LENGTHOF array  
push OFFSET array  
call DumpMem
```

stack parameters

# Parameters



call by value

```
int sum=AddTwo(a, b);
```

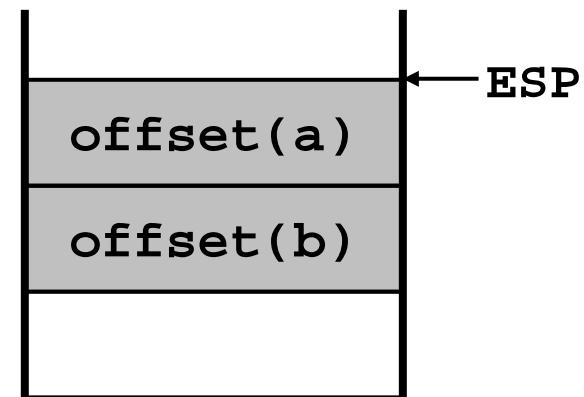
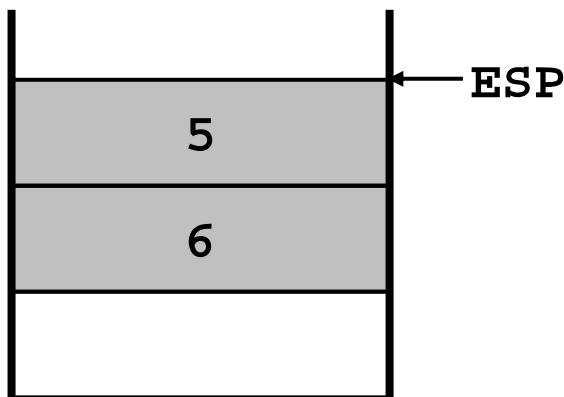
call by reference

```
int sum=AddTwo(&a, &b);
```

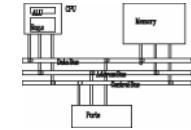
```
.date  
a    DWORD   5  
b    DWORD   6
```

```
push b  
push a  
call AddTwo
```

```
push OFFSET b  
push OFFSET a  
call AddTwo
```

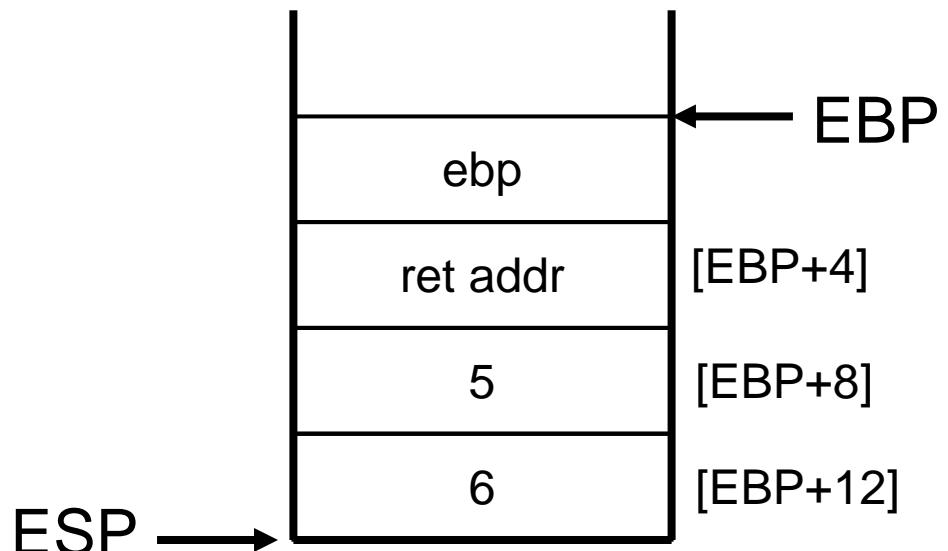


# Stack frame example

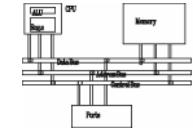


```
.data  
sum DWORD ?  
.code  
    push 6          ; second argument  
    push 5          ; first argument  
    call AddTwo    ; EAX = sum  
    mov  sum,eax   ; save the sum
```

```
AddTwo PROC  
    push ebp  
    mov  ebp,esp  
    .  
    .
```

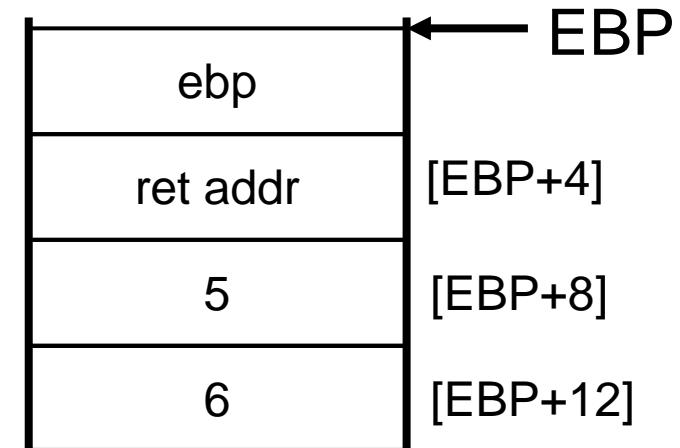


# Stack frame example



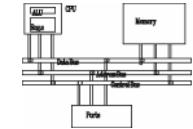
```
AddTwo PROC  
    push ebp  
    mov ebp,esp          ; base of stack frame  
    mov eax,[ebp + 12]   ; second argument (6)  
    add eax,[ebp + 8]    ; first argument (5)  
    pop ebp  
    ret 8                ; clean up the stack  
AddTwo ENDP            ; EAX contains the sum
```

Who should be responsible to remove arguments? It depends on the language model.



# RET Instruction

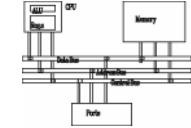
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- *Return from subroutine*
- Pops stack into the instruction pointer (EIP or IP). Control transfers to the target address.
- Syntax:
  - RET
  - RET *n*
- Optional operand *n* causes *n* bytes to be added to the stack pointer after EIP (or IP) is assigned a value.

# Passing arguments by reference

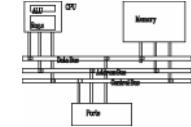
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- The **ArrayFill** procedure fills an array with 16-bit random integers
- The calling program passes the address of the array, along with a count of the number of array elements:

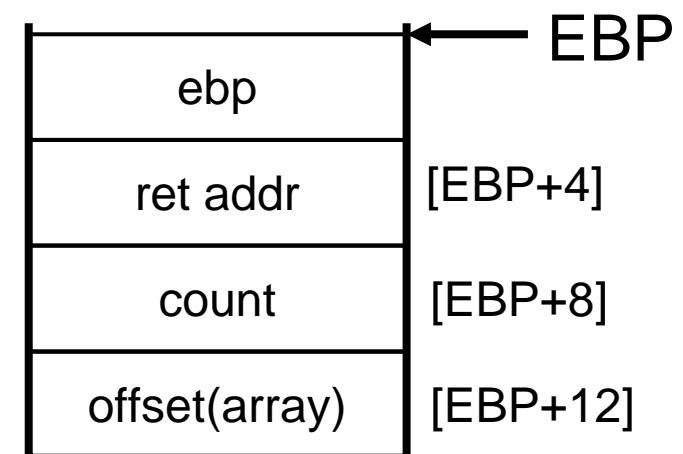
```
.data  
count = 100  
array WORD count DUP(?)  
.code  
    push OFFSET array  
    push COUNT  
    call ArrayFill
```

# Passing arguments by reference

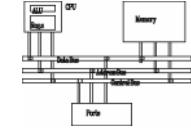


**ArrayFill** can reference an array without knowing the array's name:

```
ArrayFill PROC  
    push ebp  
    mov  ebp,esp  
    pushad  
    mov  esi,[ebp+12]  
    mov  ecx,[ebp+8]  
    .  
    .
```



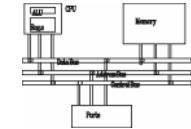
# Passing 8-bit and 16-bit arguments



- When passing stack arguments, it is best to push 32-bit operands to keep ESP aligned on a doubleword boundary.

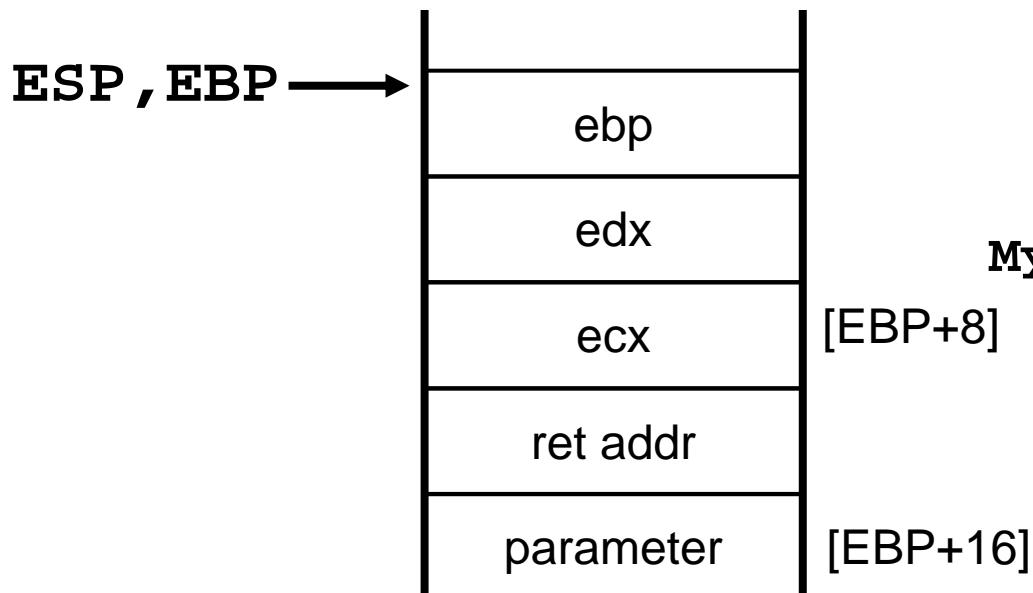
```
Uppercase PROC                push    'x'    ; error
                                Call    Uppercase
    push    ebp
    mov     ebp, esp
    mov     al, [ebp+8]
    cmp     al, 'a'
    jb      L1
    cmp     al, 'z'
    ja      L1
    sub     al, 32
L1:   pop    ebp
    ret    4
Uppercase ENDP
                                .data
                                charVal BYTE  'x'
                                .code
                                movzx eax, charVal
                                push    eax
                                Call    Uppercase
```

# Saving and restoring registers



- When using stack parameters, avoid **USES**.

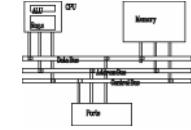
```
MySub2 PROC USES ecx, edx
    push ebp
    mov ebp, esp
    mov eax, [ebp+8]
    pop ebp
    ret 4
MySub2 ENDP
```



```
MySub2 PROC
    push ecx
    push edx
    push ebp
    mov ebp, esp
    mov eax, [ebp+8]
    pop ebp
    pop edx
    pop ecx
    ret 4
MySub2 ENDP
```

# Local variables

---



- The variables defined in the data segment can be taken as *static global variables*.

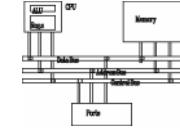
A bracket diagram where a vertical line connects the first bullet point to a horizontal line that branches into two arrows. The top arrow points to the text 'visibility=the whole program'. The bottom arrow points to the text 'lifetime=program duration'.

*visibility=the whole program*

*lifetime=program duration*

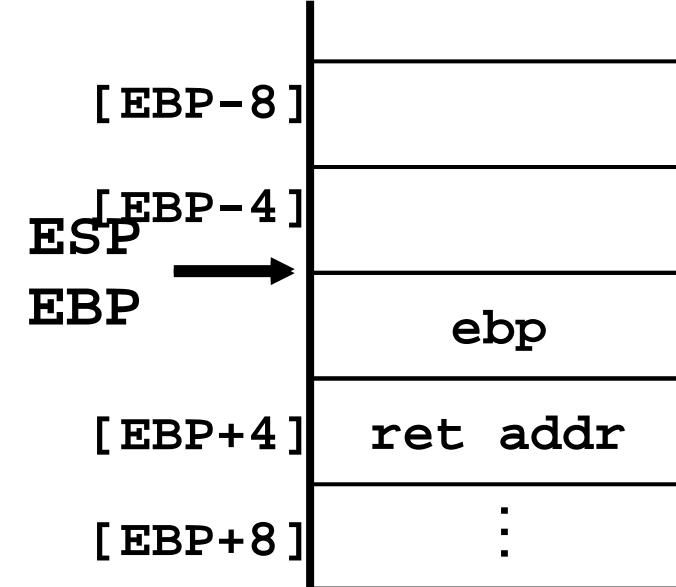
- A local variable is created, used, and destroyed within a single procedure (block)
- Advantages of local variables:
  - Restricted access: easy to debug, less error prone
  - Efficient memory usage
  - Same names can be used in two different procedures
  - Essential for recursion

# Creating local variables

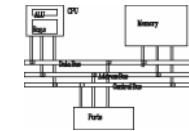


- Local variables are created on the runtime stack, usually above EBP.
- To explicitly create local variables, subtract their total size from ESP.

```
MySub PROC
    push ebp
    mov  ebp,esp
    sub  esp,8
    mov  [ebp-4],123456h
    mov  [ebp-8],0
    .
    .
```



# Local variables



- They can't be initialized at assembly time but can be assigned to default values at runtime.

```
MySub PROC
    push ebp
    mov  ebp, esp
    sub  esp, 8
void MySub()
{
    int x=10;
    int y=20;
    ...
}
    mov  DWORD PTR [ebp-4], 10
    mov  DWORD PTR [ebp-8], 20
    ...
    mov  esp, ebp
    pop  ebp
    ret
MySub ENDP
```

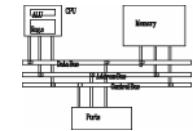
The diagram shows the state of the stack during the execution of the `MySub` function. The stack grows downwards from the current top of the stack (ESP). The stack contains the following values:

- Top of stack (ESP): 20
- Second value: 10
- Third value: EBP (the base pointer)
- Fourth value: Return address (the address of the instruction following the `ret` instruction)
- Bottom of stack (EBP): stack

Arrows point from the labels "ESP" and "EBP" to their respective stack frames. The label "return address" points to the fourth value on the stack.

# Local variables

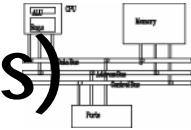
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```
x_local EQU DWORD PTR [ebp-4]  
y_local EQU DWORD PTR [ebp-8]
```

```
MySub PROC  
    push ebp  
    mov  ebp, esp  
    sub  esp, 8  
    mov  x_local, 10  
    mov  y_local, 20  
    ...  
    mov  esp, ebp  
    pop  ebp  
    ret  
MySub ENDP
```

# LEA instruction (load effective address)



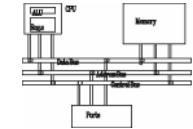
- The **LEA** instruction returns offsets of both direct and indirect operands.
  - **OFFSET** operator can only return constant offsets.
- **LEA** is required when obtaining the offset of a stack parameter or local variable. For example:

```
CopyString PROC,  
    count:DWORD  
    LOCAL temp[20]:BYTE  
  
    mov edi,OFFSET count; invalid operand  
    mov esi,OFFSET temp ; invalid operand  
    lea edi,count        ; ok  
    lea esi,temp         ; ok
```



# ENTER and LEAVE

---



- **ENTER** instruction creates stack frame for a called procedure
  - pushes EBP on the stack **push ebp**
  - set EBP to the base of stack frame **mov ebp, esp**
  - reserves space for local variables **sub esp, n**
- **ENTER nbytes, nestinglevel**
  - **nbytes** (for local variables) is rounded up to a multiple of 4 to keep ESP on a doubleword boundary
  - **nestinglevel**: 0 for now

**MySub PROC**

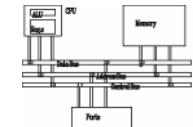
**enter 8,0**

**MySub PROC**

**push ebp  
mov ebp,esp  
sub esp,8**

# ENTER and LEAVE

---

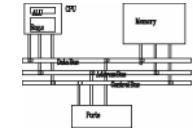


- **LEAVE** reverses the action of a previous **ENTER** instruction.

```
MySub PROC
    enter 8, 0
    .
    .
    .
    .
    leave
    ret
MySub ENDP
```

```
MySub PROC
    push ebp
    mov  ebp, esp
    sub  esp, 8
    .
    .
    .
    mov  esp, ebp
    pop  ebp
    ret
MySub ENDP
```

# **LOCAL** directive



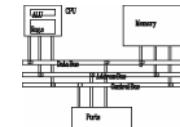
- The **LOCAL** directive declares a list of local variables
  - immediately follows the **PROC** directive
  - each variable is assigned a type
- Syntax:

**LOCAL varlist**

Example:

```
MySub PROC  
    LOCAL var1:BYTE, var2:WORD, var3:SDWORD
```

# MASM-generated code



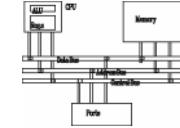
```
BubbleSort PROC  
    LOCAL temp:DWORD, SwapFlag:BYTE  
    . . .  
    ret  
BubbleSort ENDP
```

MASM generates the following code:

```
BubbleSort PROC  
    push ebp  
    mov ebp,esp  
    add esp,0FFFFFFF8h ; add -8 to ESP  
    . . .  
    mov esp,ebp  
    pop ebp  
    ret  
BubbleSort ENDP
```

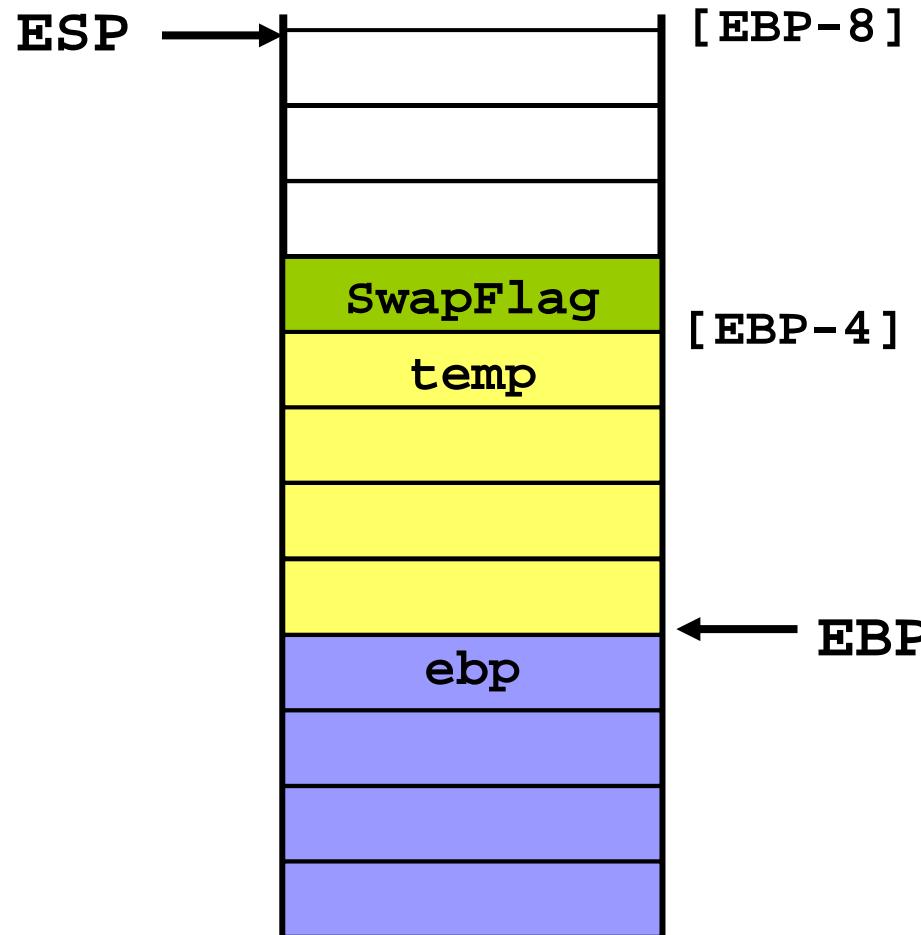
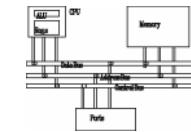
# Non-Doubleword Local Variables

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- Local variables can be different sizes
- How created in the stack by **LOCAL** directive:
  - 8-bit: assigned to next available byte
  - 16-bit: assigned to next even (word) boundary
  - 32-bit: assigned to next doubleword boundary

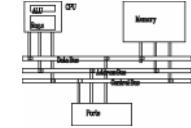
# MASM-generated code



```
mov    eax, temp           mov    eax, [ebp-4]  
mov    bl, SwapFlag        mov    bl, [ebp-5]
```

# Reserving stack space

---



- .STACK 4096
- Sub1 calls Sub2, Sub2 calls Sub3, how many bytes will you need in the stack?

Sub1 PROC

```
LOCAL array1[50]:DWORD ; 200 bytes
```

Sub2 PROC

```
LOCAL array2[80]:WORD ; 160 bytes
```

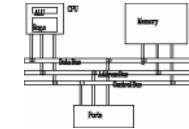
Sub3 PROC

```
LOCAL array3[300]:WORD ; 300 bytes
```

660+8(ret addr)+saved registers...

# WriteStackFrame Procedure

---

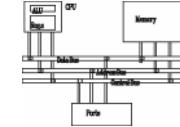


- Displays contents of current stack frame

```
WriteStackFrame PROTO,  
    numParam:DWORD, ; # of passed parameters  
    numLocalVal: DWORD, ; # of DwordLocal  
                      ; variables  
    numSavedReg: DWORD ; # of saved registers
```

# WriteStackFrame Example

---

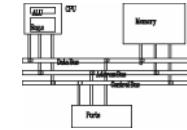


```
main PROC
    mov eax, 0EAEEAEAEAh
    mov ebx, 0EBEBEBEBBh
    INVOKE aProc, 1111h, 2222h
    exit
main ENDP

aProc PROC USES eax ebx,
    x: DWORD, y: DWORD
    LOCAL a:DWORD, b:DWORD
    PARAMS = 2
    LOCALS = 2
    SAVED_REGS = 2
    mov a,0AAAAh
    mov b,0BBBBh
    INVOKE WriteStackFrame, PARAMS, LOCALS, SAVED_REGS
```

# WriteStackFrame Example

---



## Stack Frame

00002222 ebp+12 (parameters)

00001111 ebp+8 (parameters)

00401083 ebp+4 (return address)

0012FFF0 ebp+0 (saved ebp) ← ebp

0000AAAA ebp-4 (local variable)

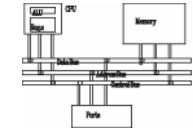
0000BBBB ebp-8 (local variable)

EAEAEAEA ebp-12 (saved register)

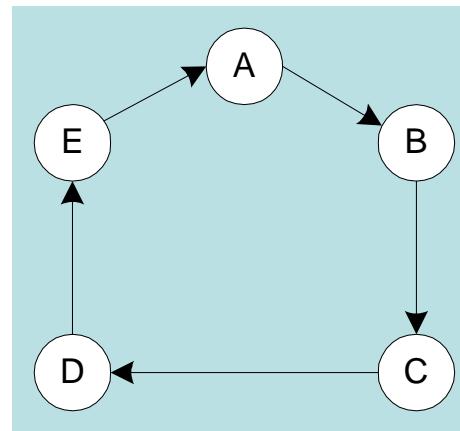
EBEBEBEB ebp-16 (saved register) ← esp

# Recursion

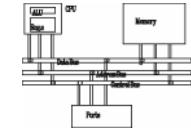
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- The process created when . . .
  - A procedure calls itself
  - Procedure A calls procedure B, which in turn calls procedure A
- Using a graph in which each node is a procedure and each edge is a procedure call, recursion forms a cycle:



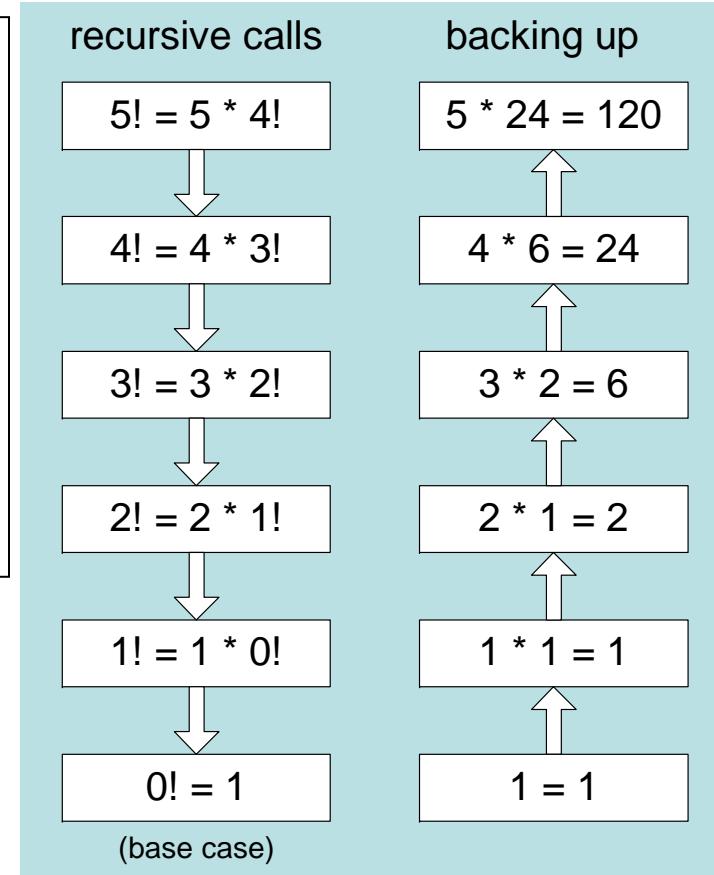
# Calculating a factorial



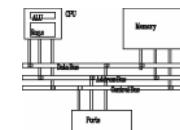
This function calculates the factorial of integer  $n$ .  
A new value of  $n$  is saved in each stack frame:

```
int factorial(int n)
{
    if (n == 0)
        return 1;
    else
        return n*factorial(n-1);
}
```

**factorial(5);**



# Calculating a factorial



```
Factorial PROC
```

```
    push ebp
    mov  ebp,esp
    mov  eax,[ebp+8]      ; get n
    cmp  eax,0            ; n > 0?
    ja   L1               ; yes: continue
    mov  eax,1            ; no: return 1
    jmp  L2
L1:dec eax
    push eax              ; Factorial(n-1)
    call Factorial
```

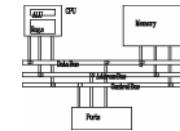
```
ReturnFact:
```

```
    mov  ebx,[ebp+8]      ; get n
    mul  ebx              ; edx:eax=eax*ebx
L2:pop  ebp              ; return EAX
    ret  4                ; clean up stack
```

```
Factorial ENDP
```

# Calculating a factorial

```
push 12  
call Factorial
```

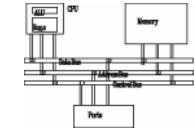


```
Factorial PROC  
    push ebp  
    mov  ebp,esp  
    mov  eax,[ebp+8]  
    cmp  eax,0  
    ja   L1  
    mov  eax,1  
    jmp  L2  
L1:dec  eax  
    push eax  
    call Factorial
```

```
ReturnFact:  
    mov  ebx,[ebp+8]  
    mul  ebx  
  
L2:pop  ebp  
    ret  4  
Factorial ENDP
```

# .MODEL directive

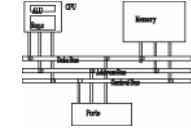
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- **.MODEL** directive specifies a program's memory model and model options (language-specifier).
- Syntax:  
`.MODEL memorymodel [,modeloptions]`
- ***memorymodel*** can be one of the following:
  - tiny, small, medium, compact, large, huge, or flat
- ***modeloptions*** includes the language specifier:
  - procedure naming scheme
  - parameter passing conventions
- **.MODEL flat, STDCALL**

# Memory models

---



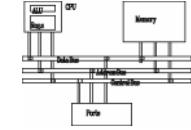
- A program's memory model determines the number and sizes of code and data segments.
- Real-address mode supports tiny, small, medium, compact, large, and huge models.
- Protected mode supports only the flat model.

Small model: code < 64 KB, data (including stack) < 64 KB.  
All offsets are 16 bits.

Flat model: single segment for code and data, up to 4 GB.  
All offsets are 32 bits.

# Language specifiers

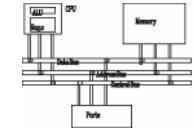
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- STDCALL (used when calling Windows functions)
  - procedure arguments pushed on stack in reverse order (right to left)
  - called procedure cleans up the stack
  - `_name@nn` (for example, `_AddTwo@8`)
- C
  - procedure arguments pushed on stack in reverse order (right to left)
  - calling program cleans up the stack (variable number of parameters such as `printf`)
  - `_name` (for example, `_AddTwo`)
- PASCAL
  - arguments pushed in forward order (left to right)
  - called procedure cleans up the stack
- BASIC, FORTRAN, SYSCALL

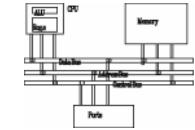
# INVOKE directive

---



- The **INVOKE** directive is a powerful replacement for Intel's **CALL** instruction that lets you pass multiple arguments
- Syntax:  
`INVOKE procedureName [, argumentList]`
- **ArgumentList** is an optional comma-delimited list of procedure arguments
- Arguments can be:
  - immediate values and integer expressions
  - variable names
  - address and ADDR expressions
  - register names

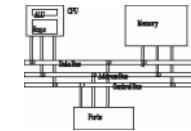
# INVOKE examples



```
.data  
byteVal BYTE 10  
wordVal WORD 1000h  
.code  
; direct operands:  
INVOKE Sub1,byteVal,wordVal  
  
; address of variable:  
INVOKE Sub2,ADDR byteVal  
  
; register name, integer expression:  
INVOKE Sub3,eax,(10 * 20)  
  
; address expression (indirect operand):  
INVOKE Sub4,[ebx]
```

# INVOKE example

---

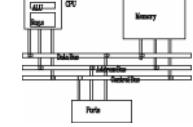


```
.data  
val1 DWORD 12345h  
val2 DWORD 23456h  
.code  
    INVOKE AddTwo, val1, val2
```

```
push val1  
push val2  
call AddTwo
```

# ADDR operator

---

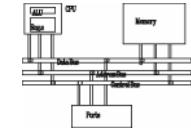


- Returns a near or far pointer to a variable, depending on which memory model your program uses:
  - Small model: returns 16-bit offset
  - Large model: returns 32-bit segment/offset
  - Flat model: returns 32-bit offset
- Simple example:

```
.data  
myWord WORD ?  
.code  
INVOKE mySub, ADDR myWord
```

# ADDR example

---

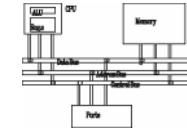


```
.data  
Array DWORD 20 DUP( ? )  
.code  
...  
INVOKE Swap, ADDR Array, ADDR [Array+4]
```

```
push OFFSET Array+4  
push OFFSET Array  
Call Swap
```

# PROC directive

---



- The **PROC** directive declares a procedure with an optional list of named parameters.
- Syntax:

```
label PROC [attributes] [USES] paramList
```

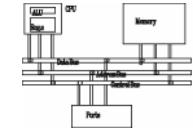
- **paramList** is a list of parameters separated by commas. Each parameter has the following syntax:

```
paramName:type
```

*type* must either be one of the standard ASM types (BYTE, SBYTE, WORD, etc.), or it can be a pointer to one of these types.

- Example: **foo PROC C USES eax, param1:DWORD**

# PROC example

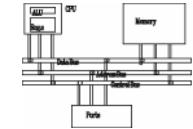


- The AddTwo procedure receives two integers and returns their sum in EAX.
- C++ programs typically return 32-bit integers from functions in EAX.

```
AddTwo PROC,  
    val1:DWORD,  
    val2:DWORD  
  
    mov eax, val1  
    add eax, val2  
    ret  
AddTwo ENDP
```

```
AddTwo PROC,  
    push ebp  
    mov ebp, esp  
    mov eax, dword ptr [ebp+8]  
    add eax, dword ptr [ebp+0Ch]  
    leave  
    ret 8  
AddTwo ENDP
```

# PROC example



```
Read_File PROC USES eax, ebx,  
    pBuffer:PTR BYTE  
    LOCAL fileHandle:DWORD
```

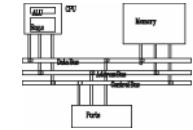
```
    mov  esi, pBuffer  
    mov  fileHandle, eax  
    .  
    .  
    ret
```

```
Read_File ENDP
```

```
Read_File PROC  
    push ebp  
    mov  ebp, esp  
    add  esp, 0FFFFFFFCh  
    push eax  
    push ebx  
    mov  esi, dword ptr [ebp+8]  
    mov  dword ptr [ebp-4], eax  
    .  
    .  
    pop  ebx  
    pop  eax  
    ret  
Read_File ENDP
```

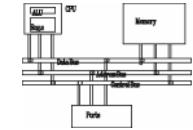
# PROTO directive

---



- Creates a procedure prototype
- Syntax:
  - *label* **PROTO** *paramList*
- Every procedure called by the **INVOKE** directive must have a prototype
- A complete procedure definition can also serve as its own prototype

# PROTO directive

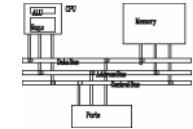


- Standard configuration: **PROTO** appears at top of the program listing, **INVOKE** appears in the code segment, and the procedure implementation occurs later in the program:

```
MySub PROTO      ; procedure prototype  
  
.code  
INVOKE MySub      ; procedure call  
  
  
MySub PROC      ; procedure implementation  
    .  
    .  
MySub ENDP
```

# PROTO example

---



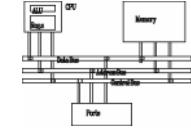
- Prototype for the ArraySum procedure, showing its parameter list:

```
ArraySum PROTO,  
    ptrArray:PTR DWORD, ; points to the array  
    szArray:DWORD        ; array size
```

```
ArraySum PROC USES esi, ecx,  
    ptrArray:PTR DWORD, ; points to the array  
    szArray:DWORD        ; array size
```

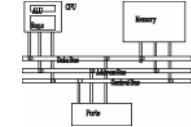
# Parameter classifications

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- An input parameter is data passed by a calling program to a procedure.
  - The called procedure is not expected to modify the corresponding parameter variable, and even if it does, the modification is confined to the procedure itself.
- An output parameter is created by passing a pointer to a variable when a procedure is called.
  - The procedure does not use any existing data from the variable, but it fills in a new value before it returns.
- An input-output parameter represents a value passed as input to a procedure, which the procedure may modify.
  - The same parameter is then able to return the changed data to the calling program.

# Example: exchanging two integers

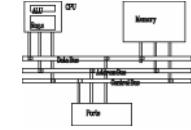


The Swap procedure exchanges the values of two 32-bit integers. pValX and pValY do not change values, but the integers they point to are modified.

```
Swap PROC USES eax esi edi,  
    pValX:PTR DWORD, ; pointer to first integer  
    pValY:PTR DWORD  ; pointer to second integer  
  
    mov esi,pValX      ; get pointers  
    mov edi,pValY  
    mov eax,[esi]       ; get first integer  
    xchg eax,[edi]     ; exchange with second  
    mov [esi],eax       ; replace first integer  
    ret ; MASM changes it to ret 8 due to PROC  
Swap ENDP
```

# Multimodule programs

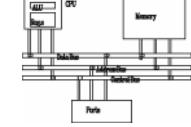
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- A multimodule program is a program whose source code has been divided up into separate ASM files.
- Each ASM file (module) is assembled into a separate OBJ file.
- All OBJ files belonging to the same program are linked using the link utility into a single EXE file.
  - This process is called static linking

# Advantages

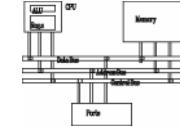
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- Large programs are easier to write, maintain, and debug when divided into separate source code modules.
- When changing a line of code, only its enclosing module needs to be assembled again. Linking assembled modules requires little time.
- A module can be a container for logically related code and data
  - encapsulation: procedures and variables are automatically hidden in a module unless you declare them public

# Creating a multimodule program

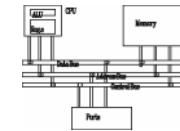
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- Here are some basic steps to follow when creating a multimodule program:
  - Create the main module
  - Create a separate source code module for each procedure or set of related procedures
  - Create an include file that contains procedure prototypes for external procedures (ones that are called between modules)
  - Use the INCLUDE directive to make your procedure prototypes available to each module

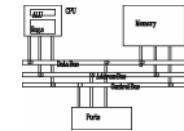
# Multimodule programs

---



- **MySub PROC PRIVATE**
- Sub PROC PUBLIC**
- **EXTERN sub1@0:PROC**
- **PUBLIC count, SYM1**  
**SYM1=10**
- .data**  
**Count DWORD 0**
- **EXTERN name:type**

# INCLUDE file



The sum.inc file contains prototypes for external functions that are not in the Irvine32 library:

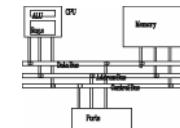
```
INCLUDE Irvine32.inc

PromptForIntegers PROTO,
    ptrPrompt:PTR BYTE,           ; prompt string
    ptrArray:PTR DWORD,          ; points to the array
    arraySize:DWORD              ; size of the array

ArraySum PROTO,
    ptrArray:PTR DWORD,          ; points to the array
    count:DWORD                  ; size of the array

DisplaySum PROTO,
    ptrPrompt:PTR BYTE,          ; prompt string
    theSum:DWORD                 ; sum of the array
```

# Main.asm



```
TITLE Integer Summation Program

INCLUDE sum.inc

.code
main PROC
    call Clrscr

    INVOKE PromptForIntegers,
        ADDR prompt1,
        ADDR array,
        Count

    ...
    call Crlf
    INVOKE ExitProcess,0
main ENDP
END main
```