The Analysis of Decomposition Methods for Support Vector Machines

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Abstract

The support vector machine (SVM) is a new and promising technique for pattern recognition. It requires the solution of a large dense quadratic programming problem. Traditional optimization methods cannot be directly applied due to memory restrictions. Up to now, very few methods can handle the memory problem and an important one is the "decomposition method." However, there is no convergence proof so far. In this paper, we connect this method to projected gradient methods and provide theoretical proofs for a version of decomposition methods. An extension to bound-constrained formulation of SVM is also provided. We then show that this convergence proof is valid for general decomposition methods if their working set selection meets a simple requirement.

I. Introduction

The support vector machine is a new and very promising classification technique for pattern recognition. Surveys of SVM are, for example, Burges [3], Cortes and Vapnik [6], Schölkopf et al. [16], and Vapnik [18]. Giving training vectors $v_i, i = 1, \ldots, m$ of length k, and a vector a defined as follows

$$a_i = \begin{cases} 1 & \text{if } v_i \text{ in class 1,} \\ -1 & \text{if } v_i \text{ in class 2,} \end{cases}$$

the support vector technique in general requires the solution of the following quadratic programming problem:

$$\begin{aligned} \min \quad & \frac{1}{2} x^T Q x - e^T x \\ & a^T x = 0, \\ & 0 \le x_i \le C, \qquad i = 1, \ldots, n, \end{aligned} \tag{1}$$

where e is the vector of all ones, C is the upper bound of all variables, Q is a positive semidefinite matrix. Possible choices of Q_{ij} are, for example, $a_i a_j ((v_i^T v_j)/k)^d$ and

$a_i a_j e^{-\|v_i - v_j\|^2/k}$. Note that v_i is considered as a *support vector* if x is the solution of (1) and $x_i > 0$.

The difficulty of solving (1) is the density of Q because Q_{ij} is in general not zero and Q becomes a fully dense matrix. Hence a prohibitive amount of memory is required to store the matrix. Thus, traditional optimization algorithms such as Newton, Quasi Newton, etc., cannot be directly applied. Several authors (e.g. [8], [9], [11], [13], [14], [15]) have proposed methods with successful implementations to conquer this difficulty. Unlike the Newton method which usually involves the whole Hessian matrix Q, these methods only calculate components of Q when they are required in the current iteration. Among them, an important one is the "decomposition method" proposed by Osuna et al. [13]. As a variation of the active set methods, they separated the index $\{1, \ldots, n\}$ of the training set to two sets B and N, where $x_i = 0$, for $i \in N$ and B is the working set if x is the current iterate of the algorithm. If we denote x_B, x_N, e_B, and e_N as vectors containing corresponding elements, the objective value is equal to $\frac{1}{2}x_B^T Q_{BB} x_B - (e_B + Q_{BN} x_N)^T x_B + \frac{1}{2}x_N^T Q_{NN} x_N - e_N^T x_N$. At each iteration, x_N is fixed and the following sub-problem with the variable x_B is solved:

$$\begin{aligned} \min \quad & \frac{1}{2}x_B^T Q_{BB} x_B - (e_B - Q_{BN} x_N)^T x_B \\ & 0 \le x_B \le C, a_B^T x_B = -a_N^T x_N, \end{aligned} \tag{2}$$

where $\left[\begin{smallmatrix} Q_{BB} & Q_{BN} \\ Q_{NB} & Q_{NN} \end{smallmatrix}\right]$ is a permutation of the matrix Q and q is the size of B. Elements of x_N are kept as zero from the beginning. The KKT condition of (2) includes $0 \le x_B \le C, a_B^T x_B = -a_N^T x_N$ of (2) and

$$\begin{aligned} Q_{BB} x_B + Q_{BN} x_N - e_B = \sigma a_B + \lambda - \mu, \\ \lambda_i (x_B)_i = 0, \mu_i (C - (x_B)_i) = 0, \lambda_i \ge 0, \mu_i \ge 0. \end{aligned} \tag{3}$$

If there is one $(x_B)_i$ satisfying $0 < (x_B)_i < C$, from (3), σ could be easily obtained. Since all components of x_N are zero, if x is an optimal solution, x_N satisfies the KKT condition of (2): $(Qx)_i - 1 - \sigma a_i \ge 0$, for $i \in N$. Therefore, the algorithm moves elements of N which violate the KKT condition of (2) to B, and move out zero elements from B to N. If the size of B is bigger than the number of support vectors, because no cycle happens, the algorithm stops in finite steps.

The main shortcoming of the above algorithm is that we do not know the number of support vectors in priori. To handle this issue, the same authors [12] proposed to use a small number q as the size of B. Furthermore, there is no restriction that x_N must only contain zeros. To be more precise, any element in N which violates the KKT condition can enter B. Since q is small in general (less than 100), this method never faces memory problem. Good numerical results were reported so several implementations and further improvements are given in [8], [14], [15]. However, even though the strict decrease of the objective function is still hold, there is no theoretical proof to show that the sequence converges to an optimal solution. This issue has been circulated in the SVM community for a while (see, for example, the discussion in Smola and Schölkopf [17, Section 5.5.2], and Keerthi et al. [10, Section 3]), but there is no satisfactory solution so far.

The paper by Joachims [8] has drawn us a lot of attention because of its method of updating B and N. Unlike the method by Osuna et al., where a more "random" selection is used, Joachims solves the following problem in order to select the working set:

$$\begin{aligned} \min \quad & \nabla f(x_k)^T d \\ & a^T d = 0, -1 \le d \le 1 && \text{(4a)} \\ & d_i \ge 0, \text{ if } (x_k)_i = 0 && \text{(4b)} \\ & d_i \le 0, \text{ if } (x_k)_i = C && \text{(4c)} \\ & |\{d_i \mid d_i \ne 0\}| = q, && \text{(4d)} \end{aligned}$$

where x_k is the kth iterate, $f(x_k) = \frac{1}{2} x_k^T Q x_k - e^T x_k$, and $\nabla f(x_k)$ is the gradient of f at x_k. Note that $|\{d_i \mid d_i \ne 0\}|$ means the number of components of d which are not zero. The constraint (4d) implies that a descent direction involving only q variables is obtained. Then components of x with non-zero d_i are included in the working set B which is used to construct the sub-problem (2). Note that d is only used for identifying B but not as a search direction. Using the decomposition method with the new technique for selecting the working set B, Joachims reported promising numerical results. In this paper, we will demonstrate that a variation of his algorithm theoretically converges.

In Section II, we present a more general algorithm with an analysis on the selection of the working set. Then in Section III, we prove the convergence of this algorithm using the

techniques of projected gradients. Section IV extends the convergence proofs to bound-constrained formulations of SVM. In Section V, we explain that our convergence proof is valid for general decomposition methods if their working set selection meets a simple requirement. Finally we provide concluding remarks in Section VI.

A preliminary version of this paper was presented in an earlier workshop [5].

II. A General Algorithm

To describe the new algorithm, we replace the problem (1) by the following form:

$$\min f(x), \quad x \in \Omega, \tag{5}$$

where $f(x)$ is a continuously differentiable function from R^n to R and the set Ω is the feasible set of (1):

$$\Omega \equiv \{x \mid a^T x = 0, 0 \le x_i \le C, i = 1, \ldots, n\}. \tag{6}$$

Therefore, (1) is a special case of (5) when $f(x) = \frac{1}{2}x^T Q x - e^T x$. In addition, we assume C is a finite positive number so Ω is a bounded set.

In the following we describe a more general algorithm for solving (5):

Step 1: Let q be a positive integer with $2 \le q \le n$. Set $k = 1$, choose $\mu_1, \mu_2 \in (0, 1), \gamma_1, \gamma_2 \ge 0$, and find an x_1 as the initial solution.

Step 2: Solve

$$\begin{aligned} \min \quad & \nabla f(x_k)^T d \\ & x_k + d \in \Omega \\ & |\{d_i \mid d_i \neq 0\}| \le q \end{aligned} \tag{7}$$

Assume B_k contains indices for which $d_i \neq 0$.

Step 3: If $d = 0$, stop and output x_k as an optimal solution. Otherwise, let

$$\bar{\nabla} f(x_k)_i \equiv \begin{cases} \nabla f(x_k) & \text{if } i \in B_k, \\ 0 & \text{otherwise.} \end{cases} \tag{8}$$

Define

$$\Omega_k \equiv \{x \mid x \in \Omega, x_i = (x_k)_i, i \notin B_k\},$$
$$R_k^n \equiv \{x \mid x \in R^n, x_i = (x_k)_i, i \notin B_k\}.$$

The mapping $P_k : R_k^n \to \Omega_k$ is defined by

$$P_k(x) \equiv \text{argmin}\{\|z - x\| : z \in \Omega_k\}.$$

For $\alpha \geq 0$, define $x_k(\alpha) \equiv P_k(x_k - \alpha\tilde{\nabla} f(x_k))$, and select $\alpha_k > 0$ so that $x_k^C \equiv P_k(x_k - \alpha_k\tilde{\nabla} f(x_k))$,

$$f(x_k^C) \leq f(x_k) + \mu_1\tilde{\nabla} f(x_k)^T(x_k^C - x_k), \tag{9}$$

and

$$\alpha_k \geq \gamma_1 \text{ or } \alpha_k \geq \gamma_2\bar{\alpha}_k > 0,$$

where $\bar{\alpha}_k$ satisfies

$$f(x_k(\bar{\alpha}_k)) > f(x_k) + \mu_2\nabla f(x_k)^T(x_k(\bar{\alpha}_k) - x_k). \tag{10}$$

Step 4: Find a new x_{k+1} such that

$$(x_{k+1})_i = (x_k)_i \text{ if } i \notin B_k \text{ and } f(x_{k+1}) \leq f(x_k^C).$$

Go to Step 2.

Joachims requires that $|\{d_i \mid d_i \neq 0\}| = q$. This may not be always possible, so we modify the equality to $|\{d_i \mid d_i \neq 0\}| \leq q$. Another difference is that we use $x_k + d \in \Omega$ instead of (4a)-(4c). Note that (4a)-(4c) do not ensure the feasibility $x_k + d \in \Omega$ so the solution of (4) may not be a feasible direction. However, we use this property for the convergence proof as will be explained in Section III.

The Step 3 of the algorithm searches for a point x_k^C by following a "partial" projected gradient direction. We called $\tilde{\nabla} f(x_k)$ a "partial" gradient direction since it contains only components of $\nabla f(x_k)$ which are in the working set of x_k. For a standard projected gradient method (e.g. Calamai and Moré [4], and Bertsekas [2]), $x_k^C \equiv P(x_k - \alpha\nabla f(x_k))$, where P is the projection into Ω, and x_k^C is referred as the *Cauchy point*. The sufficient decrease of function values at Cauchy points is guaranteed by conditions (9) and (10) and is used for the convergence proof. Then by requiring $f(x_{k+1}) \leq f(x_k^C)$, we will prove that any limit point of the sequence $\{x_k\}$ is an optimal solution. Note that when $f(x) = \frac{1}{2}x^TQx - e^Tx$, x_k^C in our algorithm is a feasible solution to the subproblem (2).

Since most existing decomposition methods obtain the optimal solution of (2) for x_{k+1}, in Step 4 the requirement that $f(x_{k+1}) \leq f(x_k^C)$ is satisfied.

Because the use of $\tilde{\nabla} f(x_k)$ instead of $\nabla f(x_k)$ makes a difference in Step 3, there is a new obstacle proving the convergence. In the rest of this section, we will construct the relation between (7) and the following problem:

$$\begin{aligned} \min \quad & \nabla f(x_k)^T d \\ & x_k + d \in \Omega \end{aligned} \tag{11}$$

Definition II.1: A point x^* is a stationary point of (5) if $\nabla f(x^*)^T (x - x^*) \geq 0$, for all $x \in \Omega$.

Immediately, we have the following result:

Lemma II.2: If $OP(x_k)$ is the optimal objective value of (11), $OP(x_k) = 0$ if and only if x_k is a stationary point.

Proof: If x_k is not a stationary point, there is an $x \in \Omega$ such that $\nabla f(x_k)^T (x - x_k) < 0$. Thus $d = x - x_k$ is a feasible solution of (11). The objective value $\nabla f(x_k)^T d < 0$ causes a contradiction. The proof of the other direction is similar. ■

We will investigate (11) in detail. Except $a^T d = 0$, constraints of (11) on d can be written as

$$l_i \leq d_i \leq u_i, \ i = 1, \ldots, n,$$

where $u_i \equiv C - (x_k)_i \geq 0, l_i \equiv -(x_k)_i \leq 0$, and $u_i \neq l_i$, for all i. The KKT condition of (11) is

$$\begin{aligned} \nabla f(x_k) - \sigma a &= \lambda - \mu, \\ \lambda_i (d_i - l_i) = 0, \lambda_i \geq 0, i &= 1, \ldots, n, \\ \mu_i (u_i - d_i) = 0, \mu_i \geq 0, i &= 1, \ldots, n, \end{aligned} \tag{12}$$

where σ is a real number, and λ_i and μ_i are multipliers associated with inequalities. If d is an optimal solution of (11), the property $l_i \neq u_i, \forall i$ and (12) imply that

$$\begin{aligned} d_i = l_i \quad &\Rightarrow \quad \nabla f(x_k)_i - \sigma a_i = \lambda_i \geq 0, \mu_i = 0, \\ l_i < d_i < u_i \quad &\Rightarrow \quad \nabla f(x_k)_i - \sigma a_i = 0 = \lambda_i = \mu_i, \\ d_i = u_i \quad &\Rightarrow \quad \nabla f(x_k)_i - \sigma a_i = -\mu_i \leq 0, \lambda_i = 0. \end{aligned}$$

As $OP(x_k)$ is the optimal objective value of (11),

$$\begin{aligned} OP(x_k) &= \sum_{i=1}^{n} \nabla f(x_k)_i d_i \\ &= \sum_{i:d_i=l_i} l_i(\nabla f(x_k)_i - \sigma a_i) + \sum_{i:d_i=u_i} u_i(\nabla f(x_k)_i - \sigma a_i). \end{aligned} \tag{13}$$

Since each term of (13) is non-positive, if j is the index such that $d_i(\nabla f(x_k)_i - \sigma a_i)$ has the smallest value,

$$d_j(\nabla f(x_k)_j - \sigma a_j) \leq \frac{OP(x_k)}{n}. \tag{14}$$

Hence if $OP(x_k) < 0$, $d_j \neq 0$. As $a_j d_j \neq 0$, with the constraint $a^T d = 0$, there are some elements i satisfying

$$(a_i d_i)(a_j d_j) < 0. \tag{15}$$

Furthermore, since all $a_i, i = 1, \ldots, n$ are 1 or -1, there is at least one index r satisfying (15) such that

$$|d_r| \geq \frac{|d_j|}{n}. \tag{16}$$

Now we are ready to switch back to (7) by considering the following problem:

$$\begin{aligned} \min \quad & \nabla f(x_k)^T d \\ & x_k + d \in \Omega \\ & d_i = 0, \forall i \neq j \text{ and } i \neq r. \end{aligned} \tag{17}$$

If $|d_r| \leq |d_j|$, define $\bar{d}_j \equiv (|d_r|/|d_j|)d_j$. Then $a_j \bar{d}_j + a_r d_r = 0$, and from (13), (14), and (16),

$$\begin{aligned} & d_r \nabla f(x_k)_r + \bar{d}_j \nabla f(x_k)_j \\ = \; & d_r \nabla f(x_k)_r + \frac{|d_r|}{|d_j|} d_j \nabla f(x_k)_j \\ = \; & d_r(\nabla f(x_k)_r - \sigma a_r) + \frac{|d_r|}{|d_j|} d_j(\nabla f(x_k)_j - \sigma a_j) \\ \leq \; & \frac{1}{n^2} OP(x_k). \end{aligned}$$

A similar setting can be made for the case when $|d_r| > |d_j|$. Therefore, by assigning other elements of d to be zero, we find a feasible solution of (17) with objective value less than $\frac{1}{n^2}OP(x_k)$. As this solution is also a feasible point of (7), if $\tilde{OP}(x_k)$ is the optimal objective value of (7), we have the following lemma:

Lemma II.3: $n^2\tilde{OP}(x_k) \leq OP(x_k)$.

From Lemma II.3, the fact that $OP(x_k) \leq \tilde{OP}(x_k)$, and Lemma II.2, we have

Lemma II.4: $\tilde{OP}(x_k) = 0$ if and only if x_k is a stationary point.

III. The Convergence Proof

Based on the analysis on special structure of SVM formulations in Section II, in this section we will provide the convergence proof.

As the only difference between P_k and P is the restriction on the domain variable x such that $x_i = (x_k)_i, i \notin B_k$, following from Calamai and Moré [4], we have

Lemma III.1: Let P_k be the projection into Ω_k:

(a) If $z \in \Omega_k$, then

$$(P_k(x) - x)^T(z - P_k(x)) \geq 0, \forall x \in R_k^n. \tag{18}$$

(b) P_k is a monotone operator, that is

$$(P_k(y) - P_k(x))^T(y - x) \geq 0, \forall x, y \in R_k^n. \tag{19}$$

If $P_k(y) \neq P_k(x)$ then strict inequality holds.

(c) Given $x \in R_k^n$ and $d \in R^n$, $d_i = 0, i \notin B_k$, the function ψ defined by

$$\psi(\alpha) \equiv \frac{\|P_k(x + \alpha d) - x\|}{\alpha},\ \alpha > 0, \tag{20}$$

is non-increasing.

We define $x_k(\alpha) \equiv P_k(x_k - \alpha\tilde{\nabla} f(x_k))$, a function of α, which yields the following inequalities: From (18),

$$\tilde{\nabla} f(x_k)^T(x_k - x_k(\alpha)) \geq \frac{\|x_k(\alpha) - x_k\|^2}{\alpha}, \alpha > 0. \tag{21}$$

Therefore,

$$\|(x_k - \alpha\tilde{\nabla} f(x_k)) - x_k\|\|x_k(\alpha) - x_k\| \geq$$
$$((x_k - \alpha\tilde{\nabla} f(x_k)) - x_k)^T(x_k(\alpha) - x_k) \geq \|x_k(\alpha) - x_k\|^2,$$

and consequently

$$\|\alpha\tilde{\nabla} f(x_k)\| \geq \|x_k(\alpha) - x_k\|. \tag{22}$$

From (19), if $\alpha \geq \beta$, then

$$\tilde{\nabla} f(x_k)^T(x_k - x_k(\alpha)) \geq \tilde{\nabla} f(x_k)^T(x_k - x_k(\beta)). \tag{23}$$

The following lemma ensures that there is an x_k^C satisfying the requirements in Step 3 of the algorithm.

Lemma III.2: If x_k is not a stationary point of (5), then

(a) For all $\alpha > 0$, $\tilde{\nabla} f(x_k)^T(x_k(\alpha) - x_k) < 0$.

(b) There is a μ in $(0, 1)$, and an $\hat{\alpha} > 0$ such that for all $0 < \alpha < \hat{\alpha}$,

$$f(x_k(\alpha)) \leq f(x_k) + \mu\tilde{\nabla} f(x_k)^T(x_k(\alpha) - x_k), \tag{24}$$

and

(c) There are positive constants γ_1 and γ_2, and constants μ_1 and μ_2 in $(0, 1)$, and $\alpha_k > 0$ such that (9) and (10) are satisfied.

Proof: If (a) is false, we can find an $\bar{\alpha}$ such that

$$\tilde{\nabla} f(x_k)^T(x_k(\bar{\alpha}) - x_k) \geq 0.$$

Then from (23), for all $0 \leq \alpha \leq \bar{\alpha}$, $\tilde{\nabla} f(x_k)^T(x_k(\alpha) - x_k) \geq 0$. However, from (21),

$$\tilde{\nabla} f(x_k)^T(x_k(\alpha) - x_k) \leq -\frac{\|x_k(\alpha) - x_k\|^2}{\alpha} \leq 0.$$

Therefore, $x_k(\alpha) = x_k$, for all $0 \leq \alpha \leq \bar{\alpha}$. Because x_k is not a stationary point, from Lemma II.4, (7) has a solution $d^* \neq 0$ and there exists an $\hat{\alpha} < \bar{\alpha}$ such that $\|\hat{\alpha}\tilde{\nabla} f(x_k)\| \leq \|d^*\|$. Because $\tilde{\nabla} f(x_k)^T d^* < 0$, we can find an $r \neq 0$ such that $(-\hat{\alpha}\tilde{\nabla} f(x_k) - \frac{d^*}{r})^T d^* = 0$. Then

$$\|(x_k + \frac{d^*}{r}) - (x_k - \hat{\alpha}\tilde{\nabla} f(x_k))\| < \|\hat{\alpha}\tilde{\nabla} f(x_k)\|.$$

This contradicts the property that $\text{argmin}\{\|z - (x_k - \hat{\alpha}\tilde{\nabla} f(x_k))\| : z \in \Omega_k\} = x_k$.

For (b), because x_k is not a stationary point and f is continuously differentiable, from (a) of this Lemma and (21), we have

$$\begin{aligned} |\frac{-f(x_k)+f(x_k(\alpha))}{\tilde{\nabla}f(x_k)^T(x_k(\alpha)-x_k)}-1| &\le \frac{o(\|x_k(\alpha)-x_k\|)}{-\tilde{\nabla}f(x_k)^T(x_k(\alpha)-x_k)} \\ &\le \frac{\alpha o(\|x_k(\alpha)-x_k\|)}{\|x_k(\alpha)-x_k\|^2}. \end{aligned} \tag{25}$$

From (20), $\|x_k(\alpha)-x_k\|/\alpha$ is non-increasing. On the other hand, from (22), $\|x_k(\alpha)-x_k\| \le \alpha\|\tilde{\nabla}f(x_k)\|$, so

$$\frac{\alpha}{\|x_k(\alpha)-x_k\|} = O(1). \tag{26}$$

Thus, (25) implies

$$\lim_{\alpha\to 0^+} \frac{-f(x_k)+f(x_k(\alpha))}{\tilde{\nabla}f(x_k)^T(x_k(\alpha)-x_k)} = 1.$$

Hence the result (24) immediately follows.

We prove (c) by giving an example from Bertsekas [1]: Define $\alpha_k = \gamma_2^{m_k}\gamma_1$, where m_k is the smallest integer such that (24) is satisfied. From (b), we know this α_k exists. Then by using this $x_k(\alpha_k)$ as x_k^C, both (9) and (10) are satisfied. ■

The following theorem from Calamai and Moré [4] is very important to our convergence proof. Because we are now using $\tilde{\nabla}f(x_k)$ instead of $\nabla f(x_k)$ as the search direction of the Cauchy step, to demonstrate its validity, here we present the whole proof.

Theorem III.3: Let $f : R^n \to R$ be continuously differentiable on Ω.

$$\lim_{k\to\infty} \frac{\|x_k^C - x_k\|}{\alpha_k} = 0.$$

Proof: Assume there is an infinite subset K_0 such that

$$\frac{\|x_k^C - x_k\|}{\alpha_k} \ge \epsilon, k \in K_0 \text{ and } \lim_{k\to\infty, k\in K_0} x_k = x^*. \tag{27}$$

We will prove that this assumption leads to a contradiction. From (27), for all $k \in K_0$,

$$\frac{\|x_k^C - x_k\|^2}{\alpha_k} \ge \epsilon \max\{\epsilon\alpha_k, \|x_k^C - x_k\|\}. \tag{28}$$

Since Ω is bounded and $x_k \in \Omega, \forall k$, $\{x_k\}$ is bounded, $\{f(x_k)\}_{k\in K_0}$ converges, and condition (9) implies $\{\tilde{\nabla}f(x_k)^T(x_k - x_k^C)\}_{k\in K_0} \to 0$. Hence (21) and (28) show that

$$\lim_{k\in K_0, k\to\infty} \alpha_k = 0 \text{ and } \lim_{k\in K_0, k\to\infty} \|x_k^C - x_k\| = 0.$$

In particular, we have shown that eventually $\alpha_k < \gamma_1$ and hence $\alpha_k \geq \gamma_2 \bar{\alpha}_k$ where $\bar{\alpha}_k$ satisfies (10). Now assign $\bar{x}_k^C = x_k(\bar{\alpha}_k)$ and $\beta_k = \min(\alpha_k, \bar{\alpha}_k)$. With this, (20) implies

$$\frac{\|x_k(\beta_k) - x_k\|^2}{\beta_k} \geq \beta_k (\frac{\|x_k^C - x_k\|}{\alpha_k})(\frac{\|\bar{x}_k^C - x_k\|}{\bar{\alpha}_k}).$$

Since $\|x_k^C - x_k\| \geq \epsilon \alpha_k$ and $\alpha_k \geq \gamma_2 \bar{\alpha}_k$ for $k \in K_0$, we obtain

$$\frac{\|x_k(\beta_k) - x_k\|^2}{\beta_k} \geq \epsilon \min\{1, \gamma_2\} \|\bar{x}_k^C - x_k\|. \tag{29}$$

Together with (21) and (23), (29) implies that for $k \in K_0$,

$$\begin{aligned} & \min\{\tilde{\nabla} f(x_k)^T (x_k - x_k^C), \tilde{\nabla} f(x_k)^T (x_k - \bar{x}_k^C)\} \\ \geq \; & \epsilon \min\{1, \gamma_2\} \|\bar{x}_k^C - x_k\|. \end{aligned} \tag{30}$$

We now use (30) to obtain the desired contradiction. Since $\{\tilde{\nabla} f(x_k)^T (x_k - x_k^C)\}$ converges to zero, (30) implies that $\{\|\bar{x}_k^C - x_k\|\}$ converges to zero. Thus the continuity of ∇f, the fact $\nabla f(x_k)^T (x_k - \bar{x}_k^C) = \tilde{\nabla} f(x_k)^T (x_k - \bar{x}_k^C)$, and the boundedness of $\{x_k\}$ show that

$$|f(\bar{x}_k^C) - f(x_k) - \tilde{\nabla} f(x_k)^T (\bar{x}_k^C - x_k)| = o(\|\bar{x}_k^C - x_k\|).$$

Therefore, if

$$\rho_k(\alpha) = \frac{f(x_k) - f(x_k(\alpha))}{\tilde{\nabla} f(x_k)^T (x_k - x_k(\alpha))},$$

then

$$|\rho_k(\bar{\alpha}) - 1| \leq \frac{o(\|\bar{x}_k^C - x_k\|)}{\tilde{\nabla} f(x_k)^T (x_k - \bar{x}_k^C)}.$$

Hence (30) establishes that $\rho_k(\bar{\alpha}) > \mu_2$ for all $k \in K_0$ sufficiently large. This is the desired contradiction because (10) guarantees that $\rho(\bar{\alpha}_k) < \mu_2$. ■

We are now in a position to prove the main theorem of this paper.

Theorem III.4: Let $f : R^n \to R$ be continuously differentiable on Ω. Any limit point of $\{x_k\}$ is a stationary point of (5).

Proof: For any $z \in \Omega_k$, from (18),

$$\begin{aligned} & \alpha_k \tilde{\nabla} f(x_k)^T (x_k^C - z) \leq -(x_k^C - x_k)^T (x_k^C - z) \\ \leq \; & -(x_k^C - x_k)^T (x_k - z) \leq \|x_k^C - x_k\| \|x_k - z\|. \end{aligned}$$

Hence,

$$\begin{aligned} & \tilde{\nabla} f(x_k)^T (x_k - z) \\ \leq \quad & \tilde{\nabla} f(x_k)^T (x_k - x_k^C) + \frac{\|x_k^C - x_k\|}{\alpha_k} \|x_k - z\|. \end{aligned}$$

If d is a direction at x_k with $x_k + d \in \Omega_k$, then setting $z = x_k + d$ yields

$$-\tilde{\nabla} f(x_k)^T d \leq \tilde{\nabla} f(x_k)^T (x_k - x_k^C) + \frac{\|x_k^C - x_k\|}{\alpha_k} \|d\|.$$

Since this inequality is true for all such d, with Lemma II.3, for any $x \in \Omega$,

$$\begin{aligned} & \nabla f(x_k)^T (x_k - x) \leq -OP(x_k) \leq -n^2 \tilde{OP}(x_k) \\ = \quad & n^2 \max_{x_k + d \in \Omega_k} -\tilde{\nabla} f(x_k)^T d \\ \leq \quad & n^2 (\tilde{\nabla} f(x_k)^T (x_k - x_k^C) + \frac{\|x_k^C - x_k\|}{\alpha_k} \max_{x_k + d \in \Omega} \|d\|). \end{aligned}$$

Theorem III.3 implies that $\|x_k^C - x_k\|/\alpha_k$ converges to zero. Since $\{f(x_k)\}$ converges, the inequality (9) implies that $\tilde{\nabla} f(x_k)^T (x_k - x_k^C)$ converges to zero. Since $f(x)$ is continuously differentiable on Ω and $\|d\|$ is bounded, if x^* is a limit point of a subsequence $\{x_k\}_{k \in K_0}$,

$$\lim_{k \in K_0, k \to \infty} \nabla f(x_k)^T (x_k - x) = \nabla f(x^*)^T (x^* - x) \leq 0.$$

Since this inequality is true for all $x \in \Omega$, from Definition II.1, x^* is a stationary point of (5). ■

Since the quadratic program of SVM is convex, Theorem III.4 assures that any limit point of $\{x_k\}$ is a global minimum.

IV. Bound-constrained Formulations of SVM

Recently several authors [11], [7] have proposed different methods for formulating support vector machines. The resulting quadratic program is a bound-constrained problem:

$$\begin{aligned} \min \quad & \frac{1}{2} x^T Q x - e^T x \\ & 0 \leq x_i \leq C, \qquad i = 1, \ldots, n. \end{aligned}$$

Note that the matrix Q here is not the same as the matrix Q in (1). In this section we would like to demonstrate that if the proposed algorithm in Section II is applied to this bound-constrained problem, the convergence proof can still follow.

If $\tilde{OP}(x_k)$ and $OP(x_k)$ are defined by the same way, now it is much easier to establish a relation between them. Instead of Lemma II.3, we have:

Lemma IV.1: $\lceil \frac{n}{q} \rceil \tilde{OP}(x_k) \leq OP(x_k)$.

Proof: First we assume d^* is an optimal solution of (11). We then divide $\{1, \ldots, n\}$ to $\lceil \frac{n}{q} \rceil$ subsets: $I_1, \ldots, I_{\lceil n/q \rceil}$. Let $I_1 = B_k$ and we distribute the rest elements to $I_2, \ldots, I_{\lceil n/q \rceil}$ such that each subset has no more than q elements. If we define vectors $d^i, i = 1, \ldots, \lceil \frac{n}{q} \rceil$ as

$$d^i_j = \begin{cases} d^*_j & j \in I_i \\ 0 & \text{otherwise,} \end{cases}$$

then all $d^i, i = 1, \ldots, \lceil \frac{n}{q} \rceil$, are feasible solutions of (7). Thus,

$$\nabla f(x_k)^T d^i \geq \tilde{OP}(x_k), \forall i = 1, \ldots, \lceil \frac{n}{q} \rceil.$$

Therefore,

$$\begin{aligned} & OP(x_k) = \nabla f(x_k)^T d^* \\ = & \sum_{i=1}^{\lceil n/q \rceil} \nabla f(x_k)^T d^i \geq \lceil \frac{n}{q} \rceil \tilde{OP}(x_k). \end{aligned}$$

■

It can be clearly seen that Lemmas III.1 and III.2, and Theorem III.3 are still true without any modifications. Then in the proof of Theorem III.4, we use Lemma IV.1 to obtain the convergence result.

V. General Working Set Selections

From the proofs in Section III and the discussion in Section IV, we realize that if the optimal objective value $\tilde{OP}(x_k)$ of the working set selection sub-problem is relatively small enough, the algorithm presented in Section II always converges.

Corollary V.1: Suppose Step 2 of the algorithm in Section II is replaced by another strategy and B_k contains indices of the selected working set. If the optimal objective value $\tilde{OP}(x_k)$ of the following problem

$$\begin{aligned} \min \quad & \nabla f(x_k)^T d \\ & x_k + d \in \Omega \qquad (31) \\ & d_i = 0, \text{ if } i \notin B_k, \qquad (32) \end{aligned}$$

satisfies

$$M \cdot \tilde{OP}(x_k) \leq OP(x_k) \leq 0, \forall k,$$

where M is a positive constant, Theorem III.4 still holds.

This corollary gives us the flexibility of choosing different types of working sets. For example, a strategy can be as follows:

Solve

$$\begin{aligned} \min \quad & \nabla f(x_k)^T d \\ & x_k + d \in \Omega \\ & |\{d_i \mid d_i \neq 0\}| \leq q/2 \end{aligned} \tag{33}$$

and define

$$B_k \equiv \{i \mid d_i \neq 0\} \cup \{q/2 \text{ indices with } d_i = 0 \text{ selected from another strategy}\}.$$

If $\hat{OP}(x_k)$ is the optimal objective value of (33), from Lemma II.3, we have

$$n^2 \tilde{OP}(x_k) \leq n^2 \hat{OP}(x_k) \leq OP(x_k).$$

Hence the algorithm converges. Such a strategy might be useful in practice as (7) is a steepest-descent type selection which may not be the best choice for quick convergence.

VI. Conclusion

In this paper we discuss the convergence of decomposition methods for support vector machines. The proof is based on the connection between a more general algorithm and the techniques of projected gradients. We also demonstrate that the proof can be easily applied to general bound-constrained formulations of SVM. Examples of flexible working set selections are also given.

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