## 1-5: Least-squares I

•  $A: k \times n$ . Usually

otherwise easily the minimum is zero.

Analytical solution:

$$f(x) = (Ax - b)^{T}(Ax - b)$$
$$= x^{T}A^{T}Ax - 2b^{T}Ax + b^{T}b$$

$$\nabla f(x) = 2A^T A x - 2A^T b = 0$$

## 1-5: Least-squares II

• Regularization, weights:

$$\frac{1}{2}\lambda x^{T}x + w_{1}(Ax - b)_{1}^{2} + \cdots + w_{k}(Ax - b)_{k}^{2}$$

# 2-4: Convex Combination and Convex Hull I

Convex hull is convex

$$x = \theta_1 x_1 + \dots + \theta_k x_k$$
  
$$\bar{x} = \bar{\theta}_1 \bar{x}_1 + \dots + \bar{\theta}_{\bar{k}} \bar{x}_{\bar{k}}$$

Then

$$\alpha x + (1 - \alpha)\bar{x}$$

$$= \alpha \theta_1 x_1 + \dots + \alpha \theta_k x_k + (1 - \alpha)\bar{\theta}_1 \bar{x}_1 + \dots + (1 - \alpha)\bar{\theta}_{\bar{k}} \bar{x}_{\bar{k}}$$

# 2-4: Convex Combination and Convex Hull II

Each coefficient is nonnegative and

$$\alpha\theta_1 + \dots + \alpha\theta_k + (1 - \alpha)\overline{\theta}_1 + \dots + (1 - \alpha)\overline{\theta}_{\bar{k}}$$

$$= \alpha + (1 - \alpha) = 1$$

## 2-7: Euclidean Balls and Ellipsoid I

We prove that any

$$x = x_c + Au$$
 with  $||u||_2 \le 1$ 

satisfies

$$(x-x_c)^T P^{-1}(x-x_c) \leq 1$$

Let

$$A=P^{1/2}$$

because *P* is symmetric positive definite.

Then

$$u^{T}A^{T}P^{-1}Au = u^{T}P^{1/2}P^{-1}P^{1/2}u \le 1.$$

## 2-10: Positive Semidefinite Cone I

•  $S_{+}^{n}$  is a convex cone. Let

$$X_1, X_2 \in \mathcal{S}^n_+$$

For any  $\theta_1 \geq 0, \theta_2 \geq 0$ ,

$$z^{T}(\theta_{1}X_{1} + \theta_{2}X_{2})z = \theta_{1}z^{T}X_{1}z + \theta_{2}z^{T}X_{2}z \geq 0$$

### 2-10: Positive Semidefinite Cone II

• Example:

$$\begin{bmatrix} x & y \\ y & z \end{bmatrix} \in \mathbf{S}_+^2$$

is equivalent to

$$x \ge 0, z \ge 0, xz - y^2 \ge 0$$

• If x > 0 or (z > 0) is fixed, we can see that

$$z \ge \frac{y^2}{x}$$

has a parabolic shape



### 2-12: Interaction I

• When t is fixed,

$$\{(x_1, x_2) \mid -1 \le x_1 \cos t + x_2 \cos 2t \le 1\}$$

gives a region between two parallel lines This region is convex

## 2-13: Affine Function I

• f(S) is convex:

Let

$$f(x_1) \in f(S), f(x_2) \in f(S)$$

$$\alpha f(x_1) + (1 - \alpha)f(x_2)$$

$$= \alpha (Ax_1 + b) + (1 - \alpha)(Ax_2 + b)$$

$$= A(\alpha x_1 + (1 - \alpha)x_2) + b$$

$$\in f(S)$$

### 2-13: Affine Function II

•  $f^{-1}(C)$  convex:

$$x_1,x_2\in f^{-1}(C)$$

means that

$$Ax_1 + b \in C, Ax_2 + b \in C$$

Because C is convex,

$$\alpha(Ax_1+b)+(1-\alpha)(Ax_2+b)$$
  
=A(\alpha x\_1+(1-\alpha)x\_2)+b\in C

Thus

$$\alpha x_1 + (1-\alpha)x_2 \in f^{-1}(\mathcal{C})$$

### 2-13: Affine Function III

Scaling:

$$\alpha S = \{ \alpha x \mid x \in S \}$$

Translation

$$S + a = \{x + a \mid x \in S\}$$

Projection

$$T = \{x_1 \in R^m \mid (x_1, x_2) \in S, x_2 \in R^n\}, S \subseteq R^m \times R^n$$

 Scaling, translation, and projection are all affine functions

## 2-13: Affine Function IV

For example, for projection

$$f(x) = \begin{bmatrix} I & 0 \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \end{bmatrix} + 0$$

*I*: identity matrix

Solution set of linear matrix inequality

$$C = \{S \mid S \leq 0\}$$
 is convex  
 $f(x) = x_1A_1 + \dots + x_mA_m - B = Ax + b$   
 $f^{-1}(C) = \{x \mid f(x) \leq 0\}$  is convex

## 2-13: Affine Function V

But this isn't rigourous because of some problems in arguing

$$f(x) = Ax + b$$

A more formal explanation:

$$C = \{s \in R^{p^2} \mid \mathsf{mat}(s) \in S^p \text{ and } \mathsf{mat}(s) \leq 0\}$$

is convex

$$f(x) = x_1 \text{vec}(A_1) + \cdots + x_m \text{vec}(A_m) - \text{vec}(B)$$
  
=  $\left[\text{vec}(A_1) \cdot \cdots \cdot \text{vec}(A_m)\right] x + \left(-\text{vec}(B)\right)$ 

## 2-13: Affine Function VI

$$f^{-1}(C) = \{x \mid \mathsf{mat}(f(x)) \in S^p \text{ and } \mathsf{mat}(f(x)) \leq 0\}$$
 is convex

Hyperbolic cone:

$$C = \{(z, t) \mid z^T z \le t^2, t \ge 0\}$$

is convex (by drawing a figure in 2 or 3 dimensional space)

## 2-13: Affine Function VII

We have that

$$f(x) = \begin{bmatrix} P^{1/2}x \\ c^Tx \end{bmatrix} = \begin{bmatrix} P^{1/2} \\ c^T \end{bmatrix} x$$

is affine. Then

$$f^{-1}(C) = \{x \mid f(x) \in C\}$$
  
= \{x \left| x^T Px \leq (c^T x)^2, c^T x \geq 0\}

is convex

## Perspective and linear-fractional function I

• Image convex: if S is convex, check if

$$\{P(x,t)\mid (x,t)\in S\}$$

convex or not

Note that S is in the domain of P

# Perspective and linear-fractional function

Assume

$$(x_1, t_1), (x_2, t_2) \in S$$

We hope

$$\alpha \frac{x_1}{t_1} + (1 - \alpha) \frac{x_2}{t_2} = P(A, B),$$

where

$$(A, B) \in S$$

# Perspective and linear-fractional function

We have

$$\alpha \frac{x_1}{t_1} + (1 - \alpha) \frac{x_2}{t_2} = \frac{\alpha t_2 x_1 + (1 - \alpha) t_1 x_2}{t_1 t_2}$$

$$= \frac{\alpha t_2 x_1 + (1 - \alpha) t_1 x_2}{\alpha t_1 t_2 + (1 - \alpha) t_1 t_2}$$

$$= \frac{\frac{\alpha t_2}{\alpha t_2 + (1 - \alpha) t_1} x_1 + \frac{(1 - \alpha) t_1}{\alpha t_2 + (1 - \alpha) t_1} x_2}{\frac{\alpha t_2}{\alpha t_2 + (1 - \alpha) t_1} t_1 + \frac{(1 - \alpha) t_1}{\alpha t_2 + (1 - \alpha) t_1} t_2}$$

# Perspective and linear-fractional function IV

Let

$$\theta = \frac{\alpha t_2}{\alpha t_2 + (1 - \alpha)t_1}$$

We have

$$\frac{\theta x_1 + (1 - \theta)x_2}{\theta t_1 + (1 - \theta)t_2} = \frac{A}{B}$$

Further

$$(A,B) \in S$$

because

$$(x_1, t_1), (x_2, t_2) \in S$$

# Perspective and linear-fractional function V

and

#### S is convex

- Inverse image is convex
- Given C a convex set

$$P^{-1}(C) = \{(x,t) \mid P(x,t) = x/t \in C\}$$

is convex

# Perspective and linear-fractional function VI

let

$$(x_1, t_1) : x_1/t_1 \in C$$
  
 $(x_2, t_2) : x_2/t_2 \in C$ 

Do we have

$$\theta(x_1, t_1) + (1 - \theta)(x_2, t_2) \in P^{-1}(C)$$
?

That is,

$$\frac{\theta x_1 + (1-\theta)x_2}{\theta t_1 + (1-\theta)t_2} \in C?$$

# Perspective and linear-fractional function VII

Let

$$\frac{\theta x_1 + (1 - \theta)x_2}{\theta t_1 + (1 - \theta)t_2} = \alpha \frac{x_1}{t_1} + (1 - \alpha)\frac{x_2}{t_2},$$

Earlier we had

$$\theta = \frac{\alpha t_2}{\alpha t_2 + (1 - \alpha)t_1}$$

Then

$$(\alpha(t_2-t_1)+t_1)\theta=\alpha t_2$$

# Perspective and linear-fractional function VIII

$$t_1\theta = \alpha t_2 - \alpha t_2\theta + \alpha t_1\theta$$
$$\alpha = \frac{t_1\theta}{t_1\theta + (1-\theta)t_2}$$

## 2-16: Generalized inequalities I

K contains no line:

$$\forall x \text{ with } x \in K \text{ and } -x \in K \Rightarrow x = 0$$

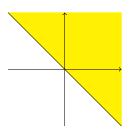
- Nonnegative polynomial on [0, 1]
- When n=2

$$x_1 \geq -tx_2, \forall t \in [0,1]$$

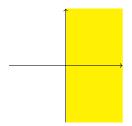
 $\bullet$  t=1



# 2-16: Generalized inequalities II

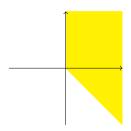






# 2-16: Generalized inequalities III

•  $\forall t \in [0,1]$ 



It really becomes a proper cone

### 2-17: I

Properties:

$$x \leq_K y, u \leq_K v$$

implies that

$$x - y \in K$$
$$u - v \in K$$

• From the definition of a convex cone,

$$(x-y)+(u-v)\in K$$

Then

$$x + u \leq_K y + v$$

## 2-18: Minimum and minimal elements I

• The minimum element

$$S \subseteq x_1 + K$$

A minimal element

$$(x_2-K)\cap S=\{x_2\}$$

## 2-19: Separating hyperplane theorem I

- We consider a simplified situation and omit part of the proof
- Assume

$$\inf_{u\in C,v\in D}\|u-v\|>0$$

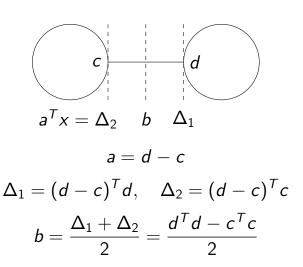
and minimum attained at c, d

We will show that

$$a \equiv d - c, \quad b \equiv \frac{\|d\|^2 - \|c\|^2}{2}$$

forms a separating hyperplane  $a^T x = b$ 

# 2-19: Separating hyperplane theorem II



# 2-19: Separating hyperplane theorem III

Assume the result is wrong so there is  $u \in D$  such that

$$a^T u - b < 0$$

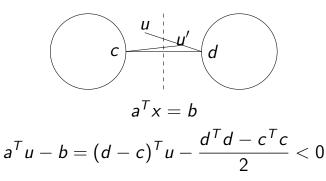
We will derive a point u' in D but closer to c than d. That is,

$$||u'-c||<||d-c||$$

Then we have a contradiction

The concept

# 2-19: Separating hyperplane theorem IV



implies that

$$(d-c)^T(u-d)+\frac{1}{2}||d-c||^2<0$$

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# 2-19: Separating hyperplane theorem V

$$\frac{d}{dt} \|d + t(u - d) - c\|^{2} \Big|_{t=0}$$

$$= 2(d + t(u - d) - c)^{T} (u - d) \Big|_{t=0}$$

$$= 2(d - c)^{T} (u - d) < 0$$

There exists a small  $t \in (0,1)$  such that

$$||d + t(u - d) - c|| < ||d - c||$$

However,

$$d + t(u - d) \in D$$
,

so there is a contradiction



## 2-20: Supporting hyperplane theorem I

### Case 1: C has an interior region

Consider 2 sets:

interior of C versus  $\{x_0\}$ ,

- where  $x_0$  is any boundary point
- If C is convex, then interior of C is also convex
- Then both sets are convex
- We can apply results in slide 2-19 so that there exists a such that

$$a^T x \leq a^T x_0, \forall x \in \text{interior of } C$$



# 2-20: Supporting hyperplane theorem II

• Then for all boundary point x we also have

$$a^T x \leq a^T x_0$$

because any boundary point is the limit of interior points

Case 2: C has no interior region

• In this situation, C is like a line in  $\mathbb{R}^3$  (so no interior). Then of course it has a supporting hyperplane

## 3-4: Examples on $R^n$ and $R^{m \times n}$ I

$$A, X \in R^{m \times n}$$
 $\operatorname{tr}(A^T X) = \sum_{j} (A^T X)_{jj}$ 
 $= \sum_{i} \sum_{j} A_{ji}^T X_{ij} = \sum_{j} \sum_{i} A_{ij} X_{ij}$ 

### 3-7: First-order Condition I

- An open set: for any x, there is a ball covering x such that this ball is in the set
- Global underestimator:

$$\frac{z-f(x)}{y-x}=f'(x)$$

$$z = f(x) + f'(x)(y - x)$$

## 3-7: First-order Condition II

•  $\Rightarrow$ : Because domain f is convex,

for all 
$$0 < t \le 1, x + t(y - x) \in \text{domain } f$$

$$f(x + t(y - x)) \le (1 - t)f(x) + tf(y)$$

$$f(y) \ge f(x) + \frac{f(x + t(y - x)) - f(x)}{t}$$

when  $t \rightarrow 0$ ,

$$f(y) \ge f(x) + \nabla f(x)^T (y - x)$$

**○** ←:

#### 3-7: First-order Condition III

For any  $0 \le \theta \le 1$ ,

$$z = \theta x + (1 - \theta)y$$

$$f(x) \ge f(z) + \nabla f(z)^{T} (x - z)$$

$$= f(z) + \nabla f(z)^{T} (1 - \theta)(x - y)$$

$$f(y) \ge f(z) + \nabla f(z)^{T} (y - z)$$

$$= f(z) + \nabla f(z)^{T} \theta (y - x)$$

$$\theta f(x) + (1 - \theta)f(y) \ge f(z)$$

#### 3-7: First-order Condition IV

• First-order condition for strictly convex function:

$$f$$
 is strictly convex if and only if  $f(y) > f(x) + \nabla f(x)^T (y - x)$ 

ullet  $\Leftarrow$ : it's easy by directly modifying  $\geq$  to >

$$f(x) > f(z) + \nabla f(z)^{T}(x - z)$$
  
=  $f(z) + \nabla f(z)^{T}(1 - \theta)(x - y)$ 

$$f(y) > f(z) + \nabla f(z)^{\mathsf{T}} (y-z) = f(z) + \nabla f(z)^{\mathsf{T}} \theta(y-x)$$

#### 3-7: First-order Condition V

•  $\Rightarrow$ : Assume the result is wrong. From the 1st-order condition of a convex function,  $\exists x, y$  such that  $x \neq y$  and

$$\nabla f(x)^{T}(y-x) = f(y) - f(x) \tag{1}$$

For this (x, y), from the strict convexity

$$f(x + t(y - x)) - f(x) < tf(y) - tf(x)$$
  
=  $\nabla f(x)^T t(y - x), \forall t \in (0, 1)$ 

#### 3-7: First-order Condition VI

Therefore,

$$f(x+t(y-x)) < f(x) + \nabla f(x)^T t(y-x), \forall t \in (0,1)$$

However, this contradicts the first-order condition:

$$f(x+t(y-x)) \ge f(x) + \nabla f(x)^T t(y-x), \forall t \in (0,1)$$

This proof was given by a student of this course before

## 3-8: Second-order condition I

- Proof of the 2nd-order condition: We consider only the simpler condition of n = 1
- $\bullet \Rightarrow$

$$f(x+t) \ge f(x) + f'(x)t$$

$$\lim_{t \to 0} 2 \frac{f(x+t) - f(x) - f'(x)t}{t^2}$$

$$= \lim_{t \to 0} \frac{2(f'(x+t) - f'(x))}{2t} = f''(x) \ge 0$$

## 3-8: Second-order condition II

● "⇐"

$$f(x + t) = f(x) + f'(x)t + \frac{1}{2}f''(\bar{x})t^2$$
  
  $\geq f(x) + f'(x)t$ 

by 1st-order condition

- The extension to general *n* is straightforward
- If  $\nabla^2 f(x) > 0$ , then f is strictly convex

## 3-8: Second-order condition III

Using 1st-order condition for strictly convex function:

$$f(y) = f(x) + \nabla f(x)^{T} (y - x) + \frac{1}{2} (y - x)^{T} \nabla^{2} f(\bar{x}) (y - x)$$
  
>  $f(x) + \nabla f(x)^{T} (y - x)$ 

Note that in our proof of 1st-order condition for strictly convex functions we used 2nd-order condition of convex function, so no contradiction here

## 3-8: Second-order condition IV

• It's possible that f is strictly convex but

$$\nabla^2 f(x) \not\succ 0$$

• Example:

$$f(x) = x^4$$

Details omitted

## 3-9: Examples I

Quadratic-over-linear

$$\frac{\partial f}{\partial x} = \frac{2x}{y}, \frac{\partial f}{\partial y} = -\frac{x^2}{y^2}$$

$$\frac{\partial^2 f}{\partial x \partial x} = \frac{2}{y}, \frac{\partial^2 f}{\partial x \partial y} = -\frac{2x}{y^2}, \frac{\partial^2 f}{\partial y \partial y} = \frac{2x^3}{y^3},$$

$$\frac{2}{y^3} \begin{bmatrix} y \\ -x \end{bmatrix} \begin{bmatrix} y & -x \end{bmatrix}$$

$$= \frac{2}{y^3} \begin{bmatrix} y^2 & -xy \\ -xy & x^2 \end{bmatrix} = 2 \begin{bmatrix} \frac{1}{y} & -\frac{x}{y^2} \\ -\frac{x}{y^2} & \frac{x^2}{y^3} \end{bmatrix}$$

### 3-10 I

$$f(x) = \log \sum_{k} \exp x_{k}$$

$$\nabla f(x) = \begin{bmatrix} \frac{e^{x_{1}}}{\sum_{k} e^{x_{k}}} \\ \vdots \\ \frac{e^{x_{n}}}{\sum_{k} e^{x_{k}}} \end{bmatrix}$$

$$\nabla_{ii}^{2} f = \frac{(\sum_{k} e^{x_{k}}) e^{x_{i}} - e^{x_{i}} e^{x_{i}}}{(\sum_{k} e^{x_{k}})^{2}}, \nabla_{ij}^{2} f = \frac{-e^{x_{i}} e^{x_{j}}}{(\sum_{k} e^{x_{k}})^{2}}, i \neq j$$

Note that if

$$z_k = \exp x_k$$



#### 3-10 II

then

$$(zz^T)_{ij} = z_i(z^T)_j = z_iz_j$$

Cauchy-Schwarz inequality

$$(a_1b_1+\cdots+a_nb_n)^2 \leq (a_1^2+\cdots+a_n^2)(b_1^2+\cdots+b_n^2)$$

$$a_k = v_k \sqrt{z_k}, b_k = \sqrt{z_k}$$

Note that

$$z_k > 0$$



## 3-12: Jensen's inequality I

General form

$$f(\int p(z)zdz) \leq \int p(z)f(z)dz$$

Discrete situation

$$f(\sum p_i z_i) \leq \sum p_i f(z_i), \sum p_i = 1$$

## 3-12: Jensen's inequality II

Proof:

$$f(p_1z_1 + p_2z_2 + p_3z_3)$$

$$\leq (1 - p_3) \left( f\left(\frac{p_1z_1 + p_2z_2}{1 - p_3}\right) \right) + p_3f(z_3)$$

$$\leq (1 - p_3) \left(\frac{p_1}{1 - p_3}f(z_1) + \frac{p_2}{1 - p_3}f(z_2)\right) + p_3f(z_3)$$

$$= p_1f(z_1) + p_2f(z_2) + p_3f(z_3)$$

Note that

$$\frac{p_1}{1-p_3} + \frac{p_2}{1-p_3} = \frac{1-p_3}{1-p_3} = 1$$

# Positive weighted sum & composition with affine function I

Composition with affine function:

We know

$$f(x)$$
 is convex

ls

$$g(x) = f(Ax + b)$$

# Positive weighted sum & composition with affine function II

convex?

$$g((1 - \alpha)x_1 + \alpha x_2)$$
=  $f(A((1 - \alpha)x_1 + \alpha x_2) + b)$   
=  $f((1 - \alpha)(Ax_1 + b) + \alpha(Ax_2 + b))$   
 $\leq (1 - \alpha)f(Ax_1 + b) + \alpha f(Ax_2 + b)$   
=  $(1 - \alpha)g(x_1) + \alpha g(x_2)$ 

#### 3-15: Pointwise maximum I

Proof of the convexity

$$f((1 - \alpha)x_1 + \alpha x_2)$$

$$= \max(f_1((1 - \alpha)x_1 + \alpha x_2), \dots, f_m((1 - \alpha)x_1 + \alpha x_2))$$

$$\leq \max((1 - \alpha)f_1(x_1) + \alpha f_1(x_2), \dots, (1 - \alpha)f_m(x_1) + \alpha f_m(x_2))$$

$$\leq (1 - \alpha)\max(f_1(x_1), \dots, f_m(x_1)) + \alpha \max(f_1(x_2), \dots, f_m(x_2))$$

$$\leq (1 - \alpha)f(x_1) + \alpha f(x_2)$$

### 3-15: Pointwise maximum II

For

$$f(x) = x_{[1]} + \cdots + x_{[r]}$$

consider all

$$\binom{n}{r}$$

combinations

## 3-16: Pointwise supremum I

- The proof is similar to pointwise maximum
- Support function of a set C:
   When y is fixed,

$$f(x,y) = y^T x$$

is linear (convex) in x

Maximum eigenvales of symmetric matrix

$$f(X,y) = y^T X y$$

is a linear function of X when y is fixed



#### 3-19: Minimization I

Proof:

Let 
$$\epsilon > 0$$
.  $\exists y_1, y_2 \in C$  such that

$$f(x_1, y_1) \leq g(x_1) + \epsilon$$

$$f(x_2, y_2) \leq g(x_2) + \epsilon$$

$$g(\theta x_1 + (1 - \theta)x_2) = \inf_{y \in C} f(\theta x_1 + (1 - \theta)x_2, y)$$

$$\leq f(\theta x_1 + (1 - \theta)x_2, \theta y_1 + (1 - \theta)y_2)$$

$$\leq \theta f(x_1, y_1) + (1 - \theta)f(x_2, y_2)$$

$$\leq \theta g(x_1) + (1 - \theta)g(x_2) + \epsilon$$

### 3-19: Minimization II

Note that the first inequality use the peroperty that
 C is convex to have

$$\theta y_1 + (1 - \theta)y_2 \in C$$

• Because the above inequality holds for all  $\epsilon > 0$ ,

$$g(\theta x_1 + (1 - \theta)x_2) \le \theta g(x_1) + (1 - \theta)g(x_2)$$

#### 3-19: Minimization III

first example: The goal is to prove

$$A - BC^{-1}B^T \succeq 0$$

Instead of a direct proof, here we use the property in this slide. First we have that f(x, y) is convex in (x, y) because

$$\begin{bmatrix} A & B \\ B^T & C \end{bmatrix} \succeq 0$$

Consider

$$\min_{y} f(x, y)$$

#### 3-19: Minimization IV

Because

$$C \succ 0$$
,

the minimum occurs at

$$2Cy + 2B^Tx = 0$$

$$y = -C^{-1}B^T x$$

Then

$$g(x) = x^{T}Ax - 2x^{T}BC^{-1}Bx + x^{T}BC^{-1}CC^{-1}B^{T}x$$
  
=  $x^{T}(A - BC^{-1}B^{T})x$ 

#### 3-19: Minimization V

is convex. The second-order condition implies that

$$A - BC^{-1}B^T \succeq 0$$

## 3-21: the conjugate function I

- This function is useful later
- When y is fixed, maximum happens at

$$y = f'(x) \tag{2}$$

by taking the derivative on x

## 3-21: the conjugate function II

Explanation of the figure: when y is fixed

$$z = xy$$

is a straight line passing through the origin, where y is the slope of the line. Check under which x,

$$yx$$
 and  $f(x)$ 

have the largest distance

• From the figure, the largest distance happens when (2) holds



## 3-21: the conjugate function III

About the point

$$(0,-f^*(y))$$

The tangent line is

$$\frac{z-f(x_0)}{x-x_0}=f'(x_0)$$

where  $x_0$  is the point satisfying

$$y=f'(x_0)$$

When x = 0.

$$z = -x_0 f'(x_0) + f(x_0) = -x_0 y + f(x_0) = -f^*(y)$$

Chih-Jen Lin (National Taiwan Univ.)

## 3-21: the conjugate function IV

•  $f^*$  is convex: Given x,

$$y^T x - f(x)$$

is linear (convex) in y. Then we apply the property of pointwise supremum

## 3-22: examples I

negative logarithm

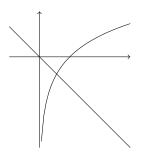
$$f(x) = -\log x$$

$$\frac{\partial}{\partial x}(xy + \log x) = y + \frac{1}{x} = 0$$

If y < 0, the picture of

$$xy + \log x$$

## 3-22: examples II



Then

$$xy + \log x = -1 - \log(-y)$$

## 3-22: examples III

strictly convex quadratic

$$Qx = y, x = Q^{-1}y$$

$$y^{T}x - \frac{1}{2}x^{T}Qx$$

$$= y^{T}Q^{-1}y - \frac{1}{2}y^{T}Q^{-1}QQ^{-1}y$$

$$= \frac{1}{2}y^{T}Q^{-1}y$$

## 3-23: quasiconvex functions I

• Figure on slide:

$$S_{\alpha} = [a, b], S_{\beta} = (-\infty, c]$$

Both are convex

 The figure is an example showing that quasi convex may not be convex

## 3-26: properties of quasiconvex functions I

Modified Jensen inequality:
 f quasiconvex if and only if

$$f(\theta x + (1-\theta)y) \le \max\{f(x), f(y)\}, \forall x, y, \theta \in [0, 1].$$

 $\bullet \Rightarrow \mathsf{Let}$ 

$$\Delta = \max\{f(x), f(y)\}$$

 $S_{\Delta}$  is convex

$$x \in S_{\Delta}, y \in S_{\Delta}$$
  
 $\theta x + (1 - \theta)y \in S_{\Delta}$ 



## 3-26: properties of quasiconvex functions

$$f(\theta x + (1 - \theta)y) \leq \Delta$$

and the result is obtained

•  $\Leftarrow$  If results are wrong, there exists  $\alpha$  such that  $S_{\alpha}$  is not convex.

$$\exists x,y,\theta$$
 with  $x,y\in\mathcal{S}_{\!lpha},\theta\in[0,1]$  such that

$$\theta x + (1 - \theta)y \notin S_{\alpha}$$

## 3-26: properties of quasiconvex functions

Then

$$f(\theta x + (1 - \theta)y) > \alpha \ge \max\{f(x), f(y)\}$$

This violates the assumption

• First-order condition (this is exercise 3.43):

## 3-26: properties of quasiconvex functions IV

"\Rightarrow"
$$f((1-t)x + ty) \le \max(f(x), f(y)) = f(x)$$

$$\frac{f(x + t(y - x)) - f(x)}{t} \le 0$$

$$\lim_{t \to 0} \frac{f(x + t(y - x)) - f(x)}{t} = \nabla f(x)^{T} (y - x) \le 0$$

## 3-26: properties of quasiconvex functions

•  $\Leftarrow$ : If results are wrong, there exists  $\alpha$  such that  $S_{\alpha}$  is not convex.

 $\exists x, y, \theta$  with  $x, y \in S_{\alpha}, \theta \in [0, 1]$  such that

$$\theta x + (1 - \theta)y \notin S_{\alpha}$$

Then

$$f(\theta x + (1 - \theta)y) > \alpha \ge \max\{f(x), f(y)\}$$
 (3)

## 3-26: properties of quasiconvex functions VI

Because f is differentiable, it is continuous. Without loss of generality, we have

$$f(z) \ge f(x), f(y), \forall z ext{ between } x ext{ and } y$$
  $z = x + heta(y - x), heta \in (0, 1)$   $\nabla f(z)^T (- heta(y - x)) \le 0$   $\nabla f(z)^T (y - x - heta(y - x)) \le 0$ 

Then

$$\nabla f(z)^T(y-x)=0, \forall \theta \in (0,1)$$

## 3-26: properties of quasiconvex functions VII

$$f(x + \theta(y - x))$$

$$= f(x) + \nabla f(t)^{T} \theta(y - x)$$

$$= f(x), \forall \theta \in [0, 1)$$

This contradicts (3).

# 3-27: Log-concave and log-convex functions I

Powers:

$$\log(x^a) = a \log x$$
$$\log x \text{ is concave}$$

Probability densities:

$$\log f(x) = -\frac{1}{2}(x - \bar{x})^T \Sigma^{-1}(x - \bar{x}) + \text{ constant}$$

 $\Sigma^{-1}$  is positive definite. Thus log f(x) is concave

## 3-27: Log-concave and log-convex functions II

Cumulative Gaussian distribution

$$\log \Phi(x) = \log \int_{-\infty}^{x} e^{-u^{2}/2} du$$

$$\frac{d}{dx} \log \Phi(x) = \frac{e^{-x^{2}/2}}{\int_{-\infty}^{x} e^{-u^{2}/2} du}$$

$$\frac{d^{2}}{d^{2}x} \log \Phi(x)$$

$$= \frac{(\int_{-\infty}^{x} e^{-u^{2}/2} du) e^{-x^{2}/2} (-x) - e^{-x^{2}/2} e^{-x^{2}/2}}{(\int_{-\infty}^{x} e^{-u^{2}/2} du)^{2}}$$

# 3-27: Log-concave and log-convex functions III

Need to prove that

$$\left(\int_{-\infty}^{x}e^{-u^2/2}du\right)x+e^{-x^2/2}>0$$

Because

$$x \ge u$$
 for all  $u \in (-\infty, x]$ ,

# 3-27: Log-concave and log-convex functions IV

we have

$$\left(\int_{-\infty}^{x} e^{-u^{2}/2} du\right) x + e^{-x^{2}/2}$$

$$= \int_{-\infty}^{x} x e^{-u^{2}/2} du + e^{-x^{2}/2}$$

$$\geq \int_{-\infty}^{x} u e^{-u^{2}/2} du + e^{-x^{2}/2}$$

$$= -e^{-u^{2}/2}|_{-\infty}^{x} + e^{-x^{2}/2}$$

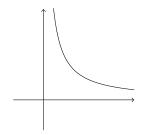
$$= -e^{-x^{2}/2} + e^{-x^{2}/2} = 0$$

# 3-27: Log-concave and log-convex functions V

This proof was given by a student (and polished by another student) of this course before

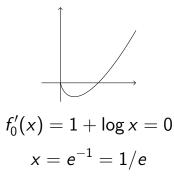
## 4-3: Optimal and locally optimal points I

• 
$$f_0(x) = 1/x$$



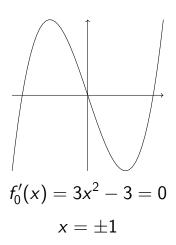
•  $f_0(x) = x \log x$ 

## 4-3: Optimal and locally optimal points II



• 
$$f_0(x) = x^3 - 3x$$

## 4-3: Optimal and locally optimal points III



## 4-9: Optimality criterion for differentiable for I

←: easyFrom first-order condition

$$f_0(y) \ge f_0(x) + \nabla f_0(x)^T (y - x)$$

Together with

$$\nabla f_0(x)^T(y-x) \geq 0$$

we have

$$f_0(y) \geq f_0(x)$$
, for all feasible  $y$ 

# 4-9: Optimality criterion for differentiable $f_0$ II

• ⇒ Assume the result is wrong. Then

$$\nabla f_0(x)^T(y-x)<0$$

Let

$$egin{aligned} z(t) &= ty + (1-t)x \ &rac{d}{dt}f_0(z(t)) &= 
abla f_0(z(t))^T(y-x) \ &rac{d}{dt}f_0(z(t))igg|_{t=0} &= 
abla f_0(x)^T(y-x) < 0 \end{aligned}$$

## 4-9: Optimality criterion for differentiable $f_0$ III

There exists t such that

$$f_0(z(t)) < f_0(x)$$

Note that

is feasible because

$$f_i(z(t)) \leq tf_i(x) + (1-t)f_i(y) \leq 0$$

and

$$A(tx+(1-t)y) = tAx+(1-t)Ay = tb+(1-b)b = b$$

Chih-Jen Lin (National Taiwan Univ.)

# 4-9: Optimality criterion for differentiable $f_0$ IV

#### 4-10 I

Unconstrained problem:

Let

$$y = x - t \nabla f_0(x)$$

It is feasible (unconstrained problem). Optimality condition implies

$$\nabla f_0(x)^T (y-x) = -t \|\nabla f_0(x)\|^2 \ge 0$$

Thus

$$\nabla f_0(x) = 0$$

Equality constrained problem

### 4-10 II

 $\Leftarrow$  Easy. For any feasible y,

$$Ay = b$$

$$\nabla f_0(x)(y-x) = -\nu^T A(y-x) = -\nu^T (b-b) = 0 \ge 0$$

So x is optimal

⇒: more complicated. We only do a rough explanation

### 4-10 III

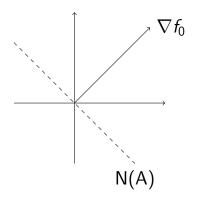
From optimality condition

$$\nabla f_0(x)^T \nu = \nabla f_0(x)^T ((x+\nu)-x) \geq 0, \forall \nu \in N(A)$$

N(A) is a subspace in 2-D. Thus

$$\nu \in N(A) \Rightarrow -\nu \in N(A)$$

## 4-10 IV



### 4-10 V

We have

$$abla f_0(x)^T 
u = 0, \forall 
u \in N(A)$$

$$abla f_0(x) \perp N(A), 
abla f_0(x) \in R(A^T)$$

$$\Rightarrow \exists 
u \text{ such that } 
abla f_0(x) + A^T 
u = 0$$

Minimization over nonnegative orthant
 Easy

### 4-10 VI

For any  $y \succeq 0$ ,

$$\nabla_i f_0(x)(y_i - x_i) = \begin{cases} \nabla_i f_0(x) y_i \ge 0 & \text{if } x_i = 0 \\ 0 & \text{if } x_i > 0. \end{cases}$$

Therefore,

$$\nabla f_0(x)^T(y-x) \geq 0$$

and

x is optimal



### 4-10 VII

 $\Rightarrow$  If  $x_i = 0$ , we claim

$$\nabla_i f_0(x) \geq 0$$

Otherwise,

$$\nabla_i f_0(x) < 0$$

Let

$$y = x$$
 except  $y_i \to \infty$ 

$$\nabla f_0(x)^T(y-x) = \nabla_i f_0(x)(y_i-x_i) \to -\infty$$

This violates the optimality condition



### 4-10 VIII

If  $x_i > 0$ , we claim

$$\nabla_i f_0(x) = 0$$

Otherwise, assume

$$\nabla_i f_0(x) > 0$$

Consider

$$y = x \text{ except } y_i = x_i/2 > 0$$



#### 4-10 IX

It is feasible. Then

$$\nabla f_0(x)^T(y-x) = \nabla_i f_0(x)(y_i-x_i) = -\nabla_i f_0(x)x_i/2 < 0$$

violates the optimality condition. The situation for

$$\nabla_i f_0(x) < 0$$

is similar



## 4-23: examples I

$$ar{c} \equiv E(C)$$

$$\Sigma \equiv E_C((C - ar{c})(C - ar{c}))$$

$$Var(C^T x) = E_C((C^T x - ar{c}^T x)(C^T x - ar{c}^T x))$$

$$= E_C(x^T (C - ar{c})(C - ar{c})^T x)$$

$$= x^T \Sigma x$$

## 4-25: second-order cone programming I

Cone was defined on slide 2-8

$$\{(x,t) \mid ||x|| \le t\}$$

## 4-35: generalized inequality constraint I

•  $f_i \in R^n \to R^{k_i}$   $K_i$ -convex:

$$f_i(\theta x + (1-\theta)y) \leq_{K_i} \theta f_i(x) + (1-\theta)f_i(y)$$

See page 3-31

#### 4-37: LP and SOCP as SDP I

LP and equivalent SDP

$$Ax = \begin{bmatrix} a_{11} & \cdots & a_{1n} \\ & \vdots & \\ a_{m1} & \cdots & a_{mn} \end{bmatrix} \begin{bmatrix} x_1 \\ \vdots \\ x_n \end{bmatrix}$$

$$x_{1}\begin{bmatrix} a_{11} & & & \\ & \ddots & & \\ & & a_{m1} \end{bmatrix} + \cdots + x_{n}\begin{bmatrix} a_{1n} & & \\ & \ddots & \\ & & a_{mn} \end{bmatrix}$$

$$-\begin{bmatrix} b_{1} & & \\ & \ddots & \\ & & b \end{bmatrix} \leq 0$$

### 4-37: LP and SOCP as SDP II

• For SOCP and SDP we will use results in 4-39:

$$\begin{bmatrix} tI_{p\times p} & A_{p\times q} \\ A^T & tI_{q\times q} \end{bmatrix} \succeq 0 \Leftrightarrow A^TA \preceq t^2I_{q\times q}, t \geq 0$$

Now

$$p = m, q = 1$$
 $A = A_i x + b_i, t = c_i^T x + d_i$ 
 $||A_i x + b_i||^2 \le (c_i^T x + d_i)^2,$ 
 $c_i^T x + d_i \ge 0$ 

Thus

$$||A_ix + b_i|| \le c_i^T x + d_i$$

### 4-39: matrix norm minimization I

 Following 4-38, we have the following equivalent problem

We then use

$$||A||_2 \le t \Leftrightarrow A^T A \le t^2 I, t \ge 0$$
$$\Leftrightarrow \begin{bmatrix} tI & A \\ A^T & tI \end{bmatrix} \succeq 0$$

### 4-39: matrix norm minimization II

to have the SDP

min 
$$t$$
 subject to 
$$\begin{bmatrix} tI & A(x) \\ A(x)^T & tI \end{bmatrix} \succeq 0$$

Next we prove

$$\begin{bmatrix} tI_{\rho\times\rho} & A_{\rho\times q} \\ A^T & tI_{q\times q} \end{bmatrix} \succeq 0 \Leftrightarrow A^TA \preceq t^2I_{q\times q}, t \geq 0$$

#### 4-39: matrix norm minimization III

⇒ we immediately have

$$t \ge 0$$

If t > 0,

$$\begin{bmatrix} -v^{T}A^{T} & tv^{T} \end{bmatrix} \begin{bmatrix} tI_{p\times p} & A_{p\times q} \\ A^{T} & tI_{q\times q} \end{bmatrix} \begin{bmatrix} -Av \\ tv \end{bmatrix}$$

$$= \begin{bmatrix} -v^{T}A^{T} & tv^{T} \end{bmatrix} \begin{bmatrix} -tAv + tAv \\ -A^{T}Av + t^{2}v \end{bmatrix}$$

$$= t(t^{2}v^{T}v - v^{T}A^{T}Av) \ge 0$$

$$v^{T}(t^{2}I - A^{T}A)v \ge 0, \forall y \in \mathbb{R}$$

### 4-39: matrix norm minimization IV

and hence

$$t^2I - A^TA \succeq 0$$

If t = 0

$$\begin{bmatrix} -v^{T}A^{T} & v^{T} \end{bmatrix} \begin{bmatrix} 0 & A \\ A^{T} & 0 \end{bmatrix} \begin{bmatrix} -Av \\ v \end{bmatrix}$$

$$= \begin{bmatrix} -v^{T}A^{T} & v^{T} \end{bmatrix} \begin{bmatrix} Av \\ -A^{T}Av \end{bmatrix}$$

$$= -2v^{T}A^{T}Av \ge 0, \forall v$$

**Therefore** 

$$A^TA \prec 0$$

### 4-39: matrix norm minimization V

← Consider

$$\begin{bmatrix} u^T & v^T \end{bmatrix} \begin{bmatrix} tI & A \\ A^T & tI \end{bmatrix} \begin{bmatrix} u \\ v \end{bmatrix}$$
$$= \begin{bmatrix} u^T & v^T \end{bmatrix} \begin{bmatrix} tu + Av \\ A^T u + tv \end{bmatrix}$$
$$= tu^T u + 2v^T A^T u + tv^T v$$

We hope to have

$$tu^T u + 2v^T A^T u + tv^T v \ge 0, \forall (u, v)$$



### 4-39: matrix norm minimization VI

If t > 0

$$\min_{u} tu^{T}u + 2v^{T}A^{T}u + tv^{T}v$$

has optimum at

$$u = \frac{-Av}{t}$$

We have

$$tu^{T}u + 2v^{T}A^{T}u + tv^{T}v$$

$$= tv^{T}v - \frac{v^{T}A^{T}Av}{t}$$

$$= \frac{1}{t}v^{T}(t^{2}I - A^{T}A)v \ge 0.$$

#### 4-39: matrix norm minimization VII

Hence

$$\begin{bmatrix} tI & A \\ A^T & tI \end{bmatrix} \succeq 0$$

If t = 0

$$A^{T}A \leq 0$$
$$v^{T}A^{T}Av \leq 0, v^{T}A^{T}Av = ||Av||^{2} = 0$$

Thus

$$Av = 0, \forall v$$

$$\begin{bmatrix} u^T & v^T \end{bmatrix} \begin{bmatrix} 0 & A \\ A^T & 0 \end{bmatrix} \begin{bmatrix} u \\ v \end{bmatrix}$$

$$= \begin{bmatrix} u^T & v^T \end{bmatrix} \begin{bmatrix} 0 \\ A^T u \end{bmatrix} = 0 \ge 0$$

#### 4-39: matrix norm minimization VIII

Thus

$$\begin{bmatrix} 0 & A \\ A^T & 0 \end{bmatrix} \succeq 0$$

### 4-40: Vector optimization I

Though

$$f_0(x)$$
 is a vector

note that

$$f_i(x)$$
 is still  $R^n \to R^1$ 

K-convex

See 3-31 though we didn't discuss it earlier

# 4-41: optimal and pareto optimal points I

Optimal

$$O \subseteq \{x\} + K$$

Pareto optimal

$$(x-K)\cap O=\{x\}$$

## 5-3: Lagrange dual function I

 Note that g is concave no matter if the original problem is convex or not

$$f_0(x) + \sum \lambda_i f_i(x) + \sum \nu_i h_i(x)$$

is convex (linear) in  $\lambda, \nu$  for each x

### 5-3: Lagrange dual function II

Use pointwise supremum on 3-16

$$\sup_{x\in D}(-f_0(x)-\sum \lambda_i f_i(x)-\sum \nu_i h_i(x))$$

is convex. Hence

$$\inf(f_0(x) + \sum \lambda_i f_i(x) + \sum \nu_i h_i(x))$$

is concave. Note that

$$-\sup(-\cdots) = -\operatorname{convex}$$
  
=  $\inf(\cdots) = \operatorname{concave}$ 



# 5-8: Lagrange dual and conjugate function

$$f_0^*(-A^T\lambda - c^T\nu)$$

$$= \sup_{x} ((-A^T\lambda - c^T\nu)^Tx - f_0(x))$$

$$= -\inf_{x} (f_0(x) + (A^T\lambda + c^T\nu)^Tx)$$

### 5-10: weak and strong duality I

 We don't discuss the SDP problem on this slide because we omitted 5-7

# 5-11: Slater's constraint qualification I

- We omit the proof because of no time
- "linear inequality do not need to hold with strict inequality": for linear inequalities we DO NOT need constraint qualification
- We will see some explanation later

### 5-12: inequality from LP I

 If we have only linear constraints, then constraint qualification holds

### 5-15: geometric interpretation l

• Explanation of  $g(\lambda)$ : when  $\lambda$  is fixed

$$\lambda u + t = \Delta$$

is a line. We lower  $\Delta$  until it touches the boundary of G

The  $\Delta$  value then becomes  $g(\lambda)$ 

When

$$u=0 \Rightarrow t=\Delta$$

so we see the point marked as  $g(\lambda)$  on t-axis

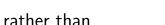


### 5-15: geometric interpretation II

• We have  $\lambda \geq 0$ , so

$$\lambda u + t = \Delta$$

must be like





### 5-15: geometric interpretation III

Explanation of p\*:
 In G, only points satisfying

$$u \leq 0$$

are feasible

#### 5-16 I

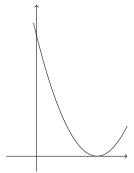
- We do not discuss a formal proof of Slater condition ⇒ strong duality Instead, we explain this result by figures
- Reason of using A: G may not be convex
- Example:

#### 5-16 II

This is a convex optimization problem

$$G = \{(x+2, x^2) \mid x \in R\}$$

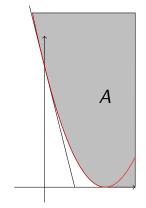
is only a quadratic curve



#### 5-16 III

The curve is not convex

• However, *A* is convex



#### 5-16 IV

Primal problem:

$$x = -2$$

optimal objective value = 4

• Dual problem:

$$g(\lambda) = \min_{x} x^2 + \lambda(x+2)$$
 
$$x = -\lambda/2$$
 
$$\max_{\lambda \geq 0} -\frac{\lambda^2}{4} + 2\lambda$$
 optimal  $\lambda = 4$ 

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#### 5-16 V

optimal objective value 
$$= -\frac{16}{4} + 8 = 4$$

Proving that A is convex

$$(u_1, t_1) \in A, (u_2, t_2) \in A$$

 $\exists x_1, x_2 \text{ such that }$ 

$$f_1(x_1) \leq u_1, f_0(x_1) \leq t_1$$

$$f_1(x_2) \leq u_2, f_0(x_2) \leq t_2$$

Consider

$$x = \theta x_1 + (1 - \theta)x_2$$



#### 5-16 VI

We have

$$f_1(x) \le \theta u_1 + (1 - \theta)u_2$$
  
 $f_0(x) \le \theta t_1 + (1 - \theta)t_2$ 

So

$$\begin{bmatrix} u \\ t \end{bmatrix} = \theta \begin{bmatrix} u_1 \\ t_1 \end{bmatrix} + (1 - \theta) \begin{bmatrix} u_2 \\ t_2 \end{bmatrix} \in A$$

Note that we have

Slater condition  $\Rightarrow$  strong duality However, it's possible that Slater condition doesn't hold but strong duality holds

#### 5-16 VII

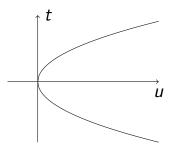
Example from exercise 5.22:

Slater condition doesn't hold because no x satisfies

$$x^{2} < 0$$

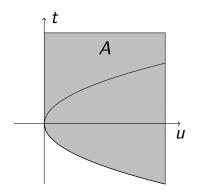
$$G = \{(x^2, x) \mid x \in R\}$$

### 5-16 VIII



There is only one feasible point (0,0)

### 5-16 IX



#### 5-16 X

$$g(\lambda) = \min_{x} x + x^{2}\lambda$$

$$x = \begin{cases} -1/(2\lambda) & \text{if } \lambda > 0\\ -\infty & \text{if } \lambda = 0 \end{cases}$$

Dual problem

$$\max_{\lambda\geq 0} -1/(4\lambda)$$
  $\lambda o\infty, ext{ objective value } o0$   $d^*=0, p^*=0$ 

Strong duality holds



## 5-17: complementary slackness I

In deriving the inequality we use

$$h_i(x^*) = 0$$
 and  $f_i(x^*) \le 0$ 

Complementary slackness
 compare the earlier results in 4-10

 For the problem on p5-16, neither slater condition nor KKT condition holds

$$1 \neq \lambda 0$$

Therefore, for convex problems,

$$KKT \Rightarrow optimality$$

but not vice versa.

 Next we explain why for linear constraints we don't need constraint qualification

• Consider the situation of inequality constraints only:

min 
$$f_0(x)$$
  
subject to  $f_i(x) \leq 0, i = 1, ..., m$ 

- Consider an optimial solution x. We would like to prove that x satisfies KKT condition
- Because x is optimal, from the optimality condition on slide 4-9, for any feasible direction  $\delta x$ ,

$$\nabla f_0(x)^T \delta x \geq 0.$$

• A feasible  $\delta x$  means

$$f_i(x + \delta x) \leq 0, \forall i$$

Because

$$f_i(x + \delta x) \approx f_i(x) + \nabla f_i(x)^T \delta x$$

from

$$f_i(x) \leq 0, \forall i$$

we have

$$\nabla f_i(x)^T \delta x \leq 0$$
 if  $f_i(x) = 0$ .

We claim that

$$\nabla f_0(x) = \sum_{\lambda_i \ge 0, f_i(x) = 0} -\lambda_i \nabla f_i(x) \tag{4}$$

Assume the result is wrong. First let's consider

$$\nabla f_0(x)$$
 =linear combination of  $\{\nabla f_i(x) \mid f_i(x) = 0\}$   
  $+ \Delta$ ,

where

$$\Delta 
eq 0$$
 and  $\Delta^T 
abla f_i(x) = 0$ ,  $\forall i: f_i(x) = 0$ 

• Then there exists  $\alpha < 0$  such that

$$\delta x \equiv \alpha \Delta$$

satisfies

$$\nabla f_i(x)^T \delta x = 0 \text{ if } f_i(x) = 0$$

$$f_i(x+\delta x)\leq 0$$
 if  $f_i(x)<0$ 

and

$$\nabla f_0(x)^T \delta x = \alpha \Delta^T \Delta < 0$$

This contradicts the optimality condition

- By a similar setting we can further prove (4).
- Assume

$$\nabla f_0(x) = \sum_{i:f_i(x)=0} -\lambda_i \nabla f_i(x)$$

and there exists i' such that

$$\lambda_{i'} < 0, \nabla f_{i'}(x) \neq 0$$
, and  $f_{i'}(x) = 0$ 

Let

$$ar{\lambda} = \arg\min_{\lambda} \| \nabla f_{i'}(x) - \sum_{i:i \neq i', f_i(x) = 0} \nabla f_i(x) \lambda_i \|$$

$$\Delta = \|\nabla f_{i'}(x) - \sum_{i:i\neq i',f_i(x)=0} \nabla f_i(x)\bar{\lambda}_i\|$$

Then

$$\nabla f_i(x)^T \Delta = 0, \forall i \neq i', f_i(x) = 0$$
 (5)

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We have

$$\nabla f_{i'}(x)^T \Delta \geq 0$$

lf

$$\nabla f_{i'}(x)^T \Delta = 0$$

then

$$-\lambda_{i'}\nabla f_{i'}(x)$$

can be rearranged to use linear combination of

$$\{\nabla f_i(x) \mid i \neq i', f_i(x) = 0\}$$



Otherwise

$$\nabla f_{i'}(x)^T \Delta > 0$$

Let

$$\delta x = \alpha \Delta, \alpha < 0.$$

• From (5),

$$\nabla f_i(x)^T \delta x = 0, \forall i \neq i', f_i(x) = 0$$

$$\nabla f_{i'}(x)^T \delta x = \alpha \nabla f_{i'}(x)^T \Delta < 0.$$

Hence  $\delta x$  is a feasible direction.

However,

$$\nabla f_0(x)^T \delta x = -\alpha \lambda_{i'} \nabla f_{i'}(x)^T \Delta < 0$$

contradicts the optimality condition

- ullet This proof is not rigourous because of pprox
- For linear the proof becomes rigourous

#### 5-25 I

• Explanation of  $f_0^*(\nu)$ 

$$\inf_{y} (f_0(y) - \nu^T y)$$
=  $-\sup_{y} (\nu^T y - f_0(y)) = -f_0^*(\nu)$ 

where  $f_0^*(\nu)$  is the conjugate function

#### 5-26 I

The original problem

$$g(\lambda, \nu) = \inf_{x} ||Ax - b|| = \text{constant}$$

• Dual norm:

$$\|\nu\|_* \equiv \sup\{\nu^T y \mid \|y\| \le 1\}$$

If 
$$\|\nu\|_* > 1$$
,

$$\nu^T y^* > 1, ||y^*|| \le 1$$



#### 5-26 II

$$\inf \|y\| + \nu^T y$$

$$\leq \| - y^*\| - \nu^T y^* < 0$$

$$\| - ty^*\| - \nu^T (ty^*) \to -\infty \text{ as } t \to \infty$$

Hence

$$\inf_{y} \|y\| + \nu^T y = -\infty$$

If  $\|\nu\|_* \leq 1$ , we claim that

$$\inf_{y} \|y\| + \nu^T y = 0$$

$$y = 0 \Rightarrow ||y|| + \nu^T y = 0$$

#### 5-26 III

If  $\exists y$  such that

$$||y|| + \nu^T y < 0$$

then

$$\|-y\|<-\nu^T y$$

We can scale y so that

$$\sup\{\nu^T y\mid \|y\|\leq 1\}>1$$

but this causes a contradiction

### 5-27: implicit constraint l

The dual function

$$c^{T}x + \nu^{T}(Ax - b)$$

$$= -b^{T}\nu + x^{T}(A^{T}\nu + c)$$

$$\inf_{-1 \le x_{i} \le 1} x_{i}(A^{T}\nu + c)_{i} = -|(A^{T}\nu + c)_{i}|$$

### 5-30: semidefinite program l

- From 5-29 we need that Z is non-negative in the dual cone of  $S_{\perp}^{k}$
- Dual cone of  $S_+^k$  is  $S_+^k$  (we didn't discuss dual cone so we assume this result)
- Why

$$tr(Z(\cdots))$$
?

We are supposed to do component-wise produt between

$$Z$$
 and  $x_1F_1 + \cdots + x_nF_n - G$ 



### 5-30: semidefinite program II

Trace is the component-wise product

$$tr(AB)$$

$$= \sum_{i} (AB)_{ii}$$

$$= \sum_{i} \sum_{j} A_{ij} B_{ji} = \sum_{i} \sum_{j} A_{ij} B_{ij}$$

Note that we take the property that B is symmetric

#### 7-4

Uniform noise

$$p(z) = \begin{cases} \frac{1}{2a} & \text{if } |z| \le a \\ 0 & \text{otherwise} \end{cases}$$

# 8-10: Dual of maximum margin problem I

Largangian:

$$\frac{\|a\|}{2} - \sum_{i} \lambda_{i} (a^{T} x_{i} + b - 1) + \sum_{i} \mu_{i} (a^{T} y_{i} + b + 1)$$

$$= \frac{\|a\|}{2} + a^{T} (-\sum_{i} \lambda_{i} x_{i} + \sum_{i} \mu_{i} y_{i})$$

$$+ b(-\sum_{i} \lambda_{i} + \sum_{i} \mu_{i}) + \sum_{i} \lambda_{i} + \sum_{i} \mu_{i}$$

Because of

$$b(-\sum_{i}\lambda_{i}+\sum_{i}\mu_{i})$$

# 8-10: Dual of maximum margin problem II

we have

$$\inf_{a,b} L$$

$$= \begin{cases}
\inf_{a,b} \frac{\|a\|}{2} - \sum_{i} \lambda_{i} a^{T} x_{i} + \sum_{i} \mu_{i} a^{T} y_{i} & \text{if } \sum_{i} \lambda_{i} = \sum_{i} \mu_{i} \\
-\infty & \text{if } \sum_{i} \lambda_{i} \neq \sum_{i} \mu_{i}
\end{cases}$$

For

$$\inf_{a} \frac{\|a\|}{2} - \sum_{i} \lambda_{i} a^{T} x_{i} + \sum_{i} \mu_{i} a^{T} y_{i}$$

# 8-10: Dual of maximum margin problem

we can denote it as

$$\inf_{a} \frac{\|a\|}{2} + v^{T} a$$

where v is a vector. We cannot do derivative because ||a|| is not differentiable. Formal solution:

# 8-10: Dual of maximum margin problem IV

• Case 1: If  $||v|| \le 1/2$ :

$$a^T v \ge -\|a\| \|v\| \ge -\frac{\|a\|}{2}$$

SO

$$\inf_{a} \frac{\|a\|}{2} + v^{T} a \ge 0.$$

However,

$$a = 0 \rightarrow \frac{\|a\|}{2} + v^T a = 0$$

# 8-10: Dual of maximum margin problem V

**Therefore** 

$$\inf_{a} \frac{\|a\|}{2} + v^T a = 0.$$

• If ||v|| > 1/2, let

$$a = \frac{-tv}{\|v\|}$$

$$\frac{\|a\|}{2} + v^{T} a$$

$$= \frac{t}{2} - t \|v\|$$

$$= t(\frac{1}{2} - \|v\|) \to \infty \text{ if } t \to -\infty$$

# 8-10: Dual of maximum margin problem VI

Thus

$$\inf_{a} \frac{\|a\|}{2} + v^{T} a = -\infty$$

#### 8-14

$$\theta = \begin{bmatrix} \operatorname{vec}(P) \\ q \\ r \end{bmatrix}, F(z) = \begin{bmatrix} \vdots \\ z_i z_j \\ \vdots \\ z_i \\ \vdots \\ 1 \end{bmatrix}$$

# 10-3: initial point and sublevel set I

• The condition that S is closed if

$$f(x) \to \infty$$
 as  $x \to$  boundary of domain  $f$ 

Proof: if not, consider

$$\{x_i\}\subset S$$

such that

$$x_i \rightarrow \text{boundary}$$

Then

$$f(x_i) \to \infty > f(x_i^0)$$

# 10-3: initial point and sublevel set II

and

S is not closed

Example

$$f(x) = \log(\sum_{i=1}^{m} \exp(a_i^T x + b_i))$$

domain  $= R^n$ 

# 10-3: initial point and sublevel set III

Example

$$f(x) = -\sum_{i} \log(b_i - a_i^T x)$$

domain 
$$\neq R^n$$

We use the condition that

$$f(x) \to \infty$$
 as  $x \to$  boundary of domain  $f$ 

# 10-4: strong convexity and implications I

• S is bounded. Otherwise, there exists a set

$$\{y_i \mid y_i = x + \Delta_i\} \subset S$$

satisfying

$$\lim_{i\to\infty} |\Delta_i| = \infty$$

Then

$$f(y_i) \ge f(x) + \nabla f(x)^T \Delta_i + \frac{m}{2} ||\Delta_i||^2 \to \infty$$

This contradicts

$$f(y) \le f(x^0)$$

# 10-4: strong convexity and implications II

Proof of

$$p^* > -\infty$$

and

$$f(x) - p^* \le \frac{1}{2m} \|\nabla f(x)\|^2$$

From

$$f(y) \ge f(x) + \nabla f(x)^{T} (y - x) + \frac{m}{2} ||x - y||^{2}$$

Minimize the right-hand side with respect to y

$$\nabla f(x) + m(y - x) = 0$$



# 10-4: strong convexity and implications III

$$\begin{split} \tilde{y} &= x - \frac{\nabla f(x)}{m} \\ f(y) &\geq f(x) + \nabla f(x)^{T} (\tilde{y} - x) + \frac{m}{2} \|\tilde{y} - x\|^{2} \\ &= f(x) - \frac{1}{2m} \|\nabla f(x)\|^{2}, \forall y \end{split}$$

Then

$$p^* \ge f(x) - \frac{1}{2m} \|\nabla f(x)\|^2$$

and

$$f(x) - p^* \le \frac{1}{2m} \|\nabla f(x)\|^2$$



#### 10-5: descent methods I

If

$$f(x+t\Delta x) < f(x)$$

then

$$\nabla f(x)^T \Delta x < 0$$

Proof: From the first-order condition of a convex function

$$f(x + t\Delta x) \ge f(x) + t\nabla f(x)^T \Delta x$$

Then

$$t\nabla f(x)^T \Delta x \leq f(x+t\Delta x)-f(x)<0$$



# 10-6: line search types l

Why

$$\alpha \in (0, \frac{1}{2})$$
?

The use of 1/2 is for convergence though we won't discuss details

• Finite termination of backtracking line search. We argue that  $\exists t^* > 0$  such that

$$f(x + t\Delta x) < f(x) + \alpha t \nabla f(x)^T \Delta x, \forall t \in (0, t^*)$$

Otherwise,

$$\exists \{t_k\} \rightarrow 0$$

# 10-6: line search types II

such that

$$f(x + t_k \Delta x) \ge f(x) + \alpha t_k \nabla f(x)^T \Delta x, \forall k$$

$$\lim_{t_k \to 0} \frac{f(x + t_k \Delta x) - f(x)}{t_k}$$

$$= \nabla f(x)^T \Delta x \ge \alpha \nabla f(x)^T \Delta x$$

However,

$$\nabla f(x)^T \Delta x < 0 \text{ and } \alpha > 0$$

cause a contradiction



# 10-6: line search types III

• Geometric interpretation: the tangent line passes through (0, f(x)), so the equation is

$$\frac{y - f(x)}{t - 0} = \nabla f(x)^{\mathsf{T}} \Delta x$$

Because

$$\nabla f(x)^T \Delta x < 0,$$

we see that the line of

$$f(x) + \alpha t \nabla f(x)^T \Delta x$$

# 10-6: line search types IV

is above that of

$$f(x) + t \nabla f(x)^T \Delta x$$

#### 10-7 I

- Linear convergence. We consider exact line search; proof for backtracking line search is more complicated
- S closed and bounded

$$\nabla^2 f(x) \leq MI, \forall x \in S$$

$$f(y) \le f(x) + \nabla f(x)^T (y - x) + \frac{M}{2} ||y - x||^2$$

Solve

$$\min_{t} f(x) - t\nabla f(x)^{T} \nabla f(x) + \frac{t^{2}M}{2} \nabla f(x)^{T} \nabla f(x)$$

#### 10-7 II

$$t=rac{1}{M}$$

$$f(x_{\text{next}}) \le f(x - \frac{1}{M} \nabla f(x)) \le f(x) - \frac{1}{2M} \nabla f(x)^T \nabla f(x)$$

The first inequality is from the fact that we use exact line search

$$f(x_{\text{next}}) - p^* \le f(x) - p^* - \frac{1}{2M} \nabla f(x)^T \nabla f(x)$$

From slide 10-4,

$$-\|\nabla f(x)\|^2 \le -2m(f(x)-p^*)$$

#### 10-7 III

Hence

$$f(x_{\mathsf{next}}) - p^* \le (1 - \frac{m}{M})(f(x) - p^*)$$

#### 10-8 I

#### Assume

$$x_1^k = \gamma (\frac{\gamma - 1}{\gamma + 1})^k, x_2^k = (-\frac{\gamma - 1}{\gamma + 1})^k,$$

$$\nabla f(x_1, x_2) = \begin{bmatrix} x_1 \\ \gamma x_2 \end{bmatrix}$$

$$\min_t \frac{1}{2} ((x_1 - tx_1)^2 + \gamma (x_2 - t\gamma x_2)^2)$$

$$\min_t \frac{1}{2} (x_1^2 (1 - t)^2 + \gamma x_2^2 (1 - t\gamma)^2)$$

#### 10-8 II

$$-x_1^2(1-t) + \gamma x_2^2(1-t\gamma)(-\gamma) = 0$$

$$-x_1^2 + tx_1^2 - \gamma^2 x_2^2 + \gamma^3 tx_2^2 = 0$$

$$t(x_1^2 + \gamma^3 x_2^2) = x_1^2 + \gamma^2 x_2^2$$

$$t = \frac{x_1^2 + \gamma^2 x_2^2}{x_1^2 + \gamma^3 x_2^2} = \frac{\gamma^2 (\frac{\gamma - 1}{\gamma + 1})^{2k} + \gamma^2 (\frac{\gamma - 1}{\gamma + 1})^{2k}}{\gamma^2 (\frac{\gamma - 1}{\gamma + 1})^{2k} + \gamma^3 (\frac{\gamma - 1}{\gamma + 1})^{2k}}$$

$$= \frac{2\gamma^2}{\gamma^2 + \gamma^3} = \frac{2}{1 + \gamma}$$

$$x^{k+1} = x^k - t\nabla f(x^k) = \begin{bmatrix} x_1^k (1 - t) \\ x_2^k (1 - \gamma t) \end{bmatrix}$$

#### 10-8 III

$$x_1^{k+1} = \gamma \left(\frac{\gamma - 1}{\gamma + 1}\right)^k \left(\frac{\gamma - 1}{1 + \gamma}\right) = \gamma \left(\frac{\gamma - 1}{\gamma + 1}\right)^{k+1}$$

$$x_2^{k+1} = \left(-\frac{\gamma - 1}{\gamma + 1}\right)^k \left(1 - \frac{2\gamma}{1 + \gamma}\right)$$

$$= \left(-\frac{\gamma - 1}{\gamma + 1}\right)^k \left(\frac{1 - \gamma}{1 + \gamma}\right) = \left(-\frac{\gamma - 1}{\gamma + 1}\right)^{k+1}$$

 Why gradient is orghogonal to the tangent line of the contour curve?

#### 10-8 IV

Assume f(g(t)) is the countour with

$$g(0) = x$$

Then

$$egin{aligned} 0 &= f(g(t)) - f(g(0)) \ 0 &= \lim_{t o 0} rac{f(g(t)) - f(g(0))}{t} \ &= \lim_{t o 0} 
abla f(g(t))^T 
abla g(t) \ &= 
abla f(x)^T 
abla g(0) \end{aligned}$$

#### 10-8 V

where

$$x + t\nabla g(0)$$

is the tangent line

#### 10-10 I

• linear convergence: from slide 10-7

$$f(x^k) - p^* \le c^k (f(x^0) - p^*)$$

$$\log(c^{k}(f(x^{0}) - p^{*})) = k \log c + \log(f(x^{0}) - p^{*})$$

is a straight line. Note that now k is the x-axis

# 10-11: steepest descent method l

• (unnormalized) steepest descent direction:

$$\Delta x_{\rm sd} = \|\nabla f(x)\|_* \Delta x_{nsd}$$

Here  $\|\cdot\|_*$  is the dual norm

 We didn't discuss much about dual norm, but we can still explain some examples on 10-12

#### 10-12 l

• Euclidean:  $\Delta x_{nsd}$  is by solving

subject to 
$$\|v\| = 1$$
 
$$\nabla f^T v = \|\nabla f\| \|v\| \cos \theta = -\|\nabla f\| \text{ when } \cos \theta = \pi$$
 
$$\Delta x_{\mathsf{nsd}} = \frac{-\nabla f(x)}{\|\nabla f(x)\|}$$
 
$$\|\nabla f(x)\|_* = \|\nabla f(x)\|$$
 
$$\|\nabla f(x)\|_* \Delta x_{\mathsf{nsd}} = \|\nabla f(x)\|_* \frac{-\nabla f(x)}{\|\nabla f(x)\|_*} = -\nabla f(x)$$

min  $\nabla f^T v$ 

#### 10-12 II

• Quadratic norm:  $\Delta x_{nsd}$  is by solving

Now

$$||v||_P = \sqrt{v^T P v},$$

where P is symmetric positive definite

### 10-12 III

Let

$$w = P^{1/2}v$$

The optimization problem becomes

$$\min_{w} \quad \nabla f^T P^{-1/2} w$$
 subject to  $\|w\| = 1$ 

optimal 
$$w = \frac{-P^{-1/2}\nabla f}{\|P^{-1/2}\nabla f\|}$$
$$= \frac{-P^{-1/2}\nabla f}{\sqrt{\nabla f^T P^{-1}\nabla f}}$$

## 10-12 IV

optimal 
$$v = \frac{-P^{-1}\nabla f}{\sqrt{\nabla f^T P^{-1}\nabla f}}$$

Dual norm

$$||z||_* = ||P^{-1/2}z||$$

Therefore

$$\Delta x_{sd} = \sqrt{\nabla f^T P^{-1} \nabla f} \frac{-P^{-1} \nabla f}{\sqrt{\nabla f^T P^{-1} \nabla f}}$$
$$= -P^{-1} \nabla f$$

## 10-12 V

Explanation of the figure:

$$-\nabla f(x)^T \Delta x_{\mathsf{nsd}} = \|-\nabla f(x)\| \|\Delta x_{\mathsf{nsd}}\| \cos \theta$$

 $\|-\nabla f(x)\|$  is a constant. From a point  $\Delta x_{\rm nsd}$  on the boundary, the projected point on  $-\nabla f(x)$  indicates

$$\|\Delta x_{\mathsf{nsd}}\|\cos\theta$$

In the figure, we see that the chosen  $\Delta x_{\rm nsd}$  has the largest  $\|\Delta x_{\rm nsd}\|\cos\theta$ 

• We omit the discussion of  $l_1$ -norm



## 10-13 I

- The two figures are by using two P matrices
- The left one has faster convergence
- Gradient descent after change of variables

$$\bar{x} = P^{1/2}x, x = P^{-1/2}\bar{x}$$

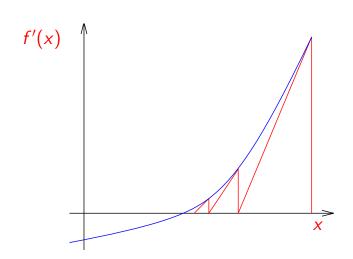
$$\min_{x} f(x) \Rightarrow \min_{\bar{x}} f(P^{-1/2}\bar{x})$$

$$\bar{x} \leftarrow \bar{x} - \alpha P^{-1/2}\nabla_{x} f(P^{-1/2}\bar{x})$$

$$P^{1/2}x \leftarrow P^{1/2}x - \alpha P^{-1/2}\nabla_{x} f(x)$$

$$x \leftarrow x - \alpha P^{-1}\nabla_{x} f(x)$$

## 10-14 I



## 10-14 II

Solve

$$f'(x) = 0$$

Fining the tangent line at  $x_k$ :

$$\frac{y-f'(x_k)}{x-x_k}=f''(x_k)$$

 $x_k$ : the current iterate

Let 
$$y = 0$$

$$x_{k+1} = x_k - f'(x_k)/f''(x_k)$$

## 10-16 I

$$\hat{f}(y) = f(x) + \nabla f(x)^{T} (y - x) + \frac{1}{2} (y - x)^{T} \nabla^{2} f(x) (y - x)$$

$$\nabla \hat{f}(y) = 0 = \nabla f(x) + \nabla^{2} f(x) (y - x)$$

$$y - x = -\nabla^{2} f(x)^{-1} \nabla f(x)$$

$$\inf_{y} \hat{f}(y) = f(x) - \frac{1}{2} \nabla f(x)^{T} \nabla^{2} f(x)^{-1} \nabla f(x)$$

$$f(x) - \inf_{y} \hat{f}(y) = \frac{1}{2} \lambda(x)^{2}$$

## 10-16 II

Norm of the Newton step in the quadratic Hessian norm

$$\Delta x_{\rm nt} = -\nabla^2 f(x)^{-1} \nabla f(x)$$

$$\Delta x_{\mathsf{nt}}^T \nabla^2 f(x) \Delta x_{\mathsf{nt}} = \nabla f(x)^T \nabla^2 f(x)^{-1} \nabla f(x) = \lambda(x)^2$$

Directional derivative in the Newton direction

$$\lim_{t \to 0} \frac{f(x + t\Delta_{nt}) - f(x)}{t}$$

$$= \nabla f(x)^T \Delta x_{nt}$$

$$= -\nabla f(x)^T \nabla^2 f(x)^{-1} \nabla f(x) = -\lambda(x)^2$$

## 10-16 III

Affine invariant

$$\bar{f}(y) \equiv f(Ty) = f(x)$$

Assume T is an invertable square matrix. Then

$$\bar{\lambda}(y) = \lambda(Ty)$$

Proof:

$$abla ar{f}(y) = T^T \nabla f(Ty)$$

$$abla^2 ar{f}(y) = T^T \nabla^2 f(Ty) T$$

## 10-16 IV

$$\bar{\lambda}(y)^{2} = \nabla \bar{f}(y)^{T} \nabla^{2} \bar{f}(y)^{-1} \nabla \bar{f}(y) 
= \nabla f(Ty)^{T} T T^{-1} \nabla^{2} f(Ty)^{-1} T^{-T} T^{T} \nabla f(Ty) 
= \nabla f(Ty)^{T} \nabla^{2} f(Ty)^{-1} \nabla f(Ty) 
= \lambda (Ty)^{2}$$

## 10-17 I

#### Affine invariant

$$\Delta y_{nt} = \nabla^2 \bar{f}(y)^{-1} \nabla \bar{f}(y)$$

$$= T^{-1} \nabla^2 f(Ty)^{-1} \nabla f(Ty)$$

$$= T^{-1} \Delta x_{nt}$$

Note that

$$y_k = T^{-1} x_k$$

SO

$$y_{k+1} = T^{-1} x_{k+1}$$



## 10-17 II

#### But how about line search

$$\nabla \bar{f}(y)^{T} \Delta y_{nt}$$

$$= \nabla f(Ty)^{T} T T^{-1} \Delta x_{nt}$$

$$= \nabla f(x)^{T} \Delta x_{nt}$$

## 10-19 I

$$\eta \in (0, \frac{m^2}{L})$$
$$\|\nabla f(x_k)\| \le \eta \le \frac{m^2}{L}$$
$$\frac{L}{2m^2} \|\nabla f(x_k)\| \le \frac{1}{2}$$

## 10-20 I

$$f(x_{l}) - f(x^{*})$$

$$\leq \frac{1}{2m} \|\nabla f(x_{l})\|^{2} \quad \text{(from p10-4)}$$

$$\leq \frac{1}{2m} \frac{4m^{4}}{L^{2}} (\frac{1}{2})^{2^{l-k} \cdot 2}$$

$$= \frac{2m^{3}}{L^{2}} (\frac{1}{2})^{2^{l-k+1}} \leq \epsilon$$

Let

$$\epsilon_0 = \frac{2m^3}{L^2}$$



## 10-20 II

$$egin{aligned} \log_2 \epsilon_0 - 2^{l-k+1} &\leq \log_2 \epsilon \ 2^{l-k+1} &\geq \log_2(\epsilon_0/\epsilon) \ l &\geq k-1 + \log_2 \log_2(\epsilon_0/\epsilon) \ k &\leq rac{f(x_0) - p^*}{r} \end{aligned}$$

In at most

$$\frac{f(x_0) - p^*}{r} + \log_2 \log_2(\epsilon_0/\epsilon)$$

iterations, we have

$$f(x_l) - f(x^*) \le \epsilon$$

## 10-29: implementation I

$$\lambda(x) = (\nabla f(x) \nabla^2 f(x)^{-1} \nabla f(x))^{1/2}$$
  
=  $(g^T L^{-T} L^{-1} g)^{1/2} = ||L^{-1} g||_2$ 

# 10-30: example of dense Newton systems with structure I

$$abla f(x) = egin{bmatrix} \psi_1'(x_1) \ dots \ \psi_n'(x_n) \end{bmatrix} + A^T 
abla \psi_0(Ax + b)$$

$$abla^2 f(x) = \begin{bmatrix} \psi_1''(x_1) & & & \\ & \ddots & & \\ & & \psi_n''(x_n) \end{bmatrix} + A^T 
abla^2 \psi_0^2 (Ax + b) A$$

method 2:

$$\Delta x = D^{-1}(-g - A^T L_o w)$$

## 10-30: example of dense Newton systems with structure II

$$L_0^T A D^{-1} (-g - A^T L_0 w) = w$$
$$(I + L_0^T A D^{-1} A^T L_0) w = -L_0^T A D^{-1} g$$

Cost

$$L_0: p \times p$$
  
 $A^T L_0: n \times p, \text{cost} : O(np)$   
 $(L_0^T A) D^{-1} (A^T L_0) : O(p^2 n)$ 

# 10-30: example of dense Newton systems with structure III

Note that Cholesky factorization of  $H_0$  costs

$$\frac{1}{3}p^3 \le p^2n$$

as

$$p \ll n$$