

Definition of GNFA I

- Here we give the formal definition of generalized NFA
- Between any two states: a regular expression
- $(Q, \Sigma, \delta, q_{start}, q_{accept})$
- Single accept state. No longer a set F
- The δ function:

$$(Q - \{q_{accept}\}) \times (Q - \{q_{start}\}) \rightarrow R$$

R : all regular expressions over Σ

- DFA \rightarrow GNFA

Definition of GNFA II

Two new states: q_{start}, q_{accept}

$q_{start} \rightarrow q_0$ with ϵ

any $q \in F \rightarrow q_{accept}$ with ϵ

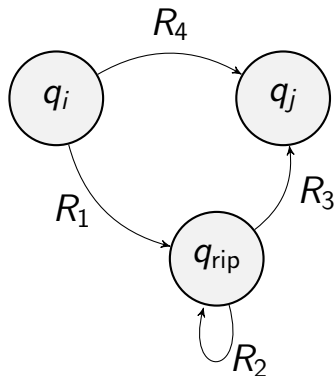
- In the definition, between any two states there is an expression

But what if in the graph two states are not connected ?

$\emptyset \in R$ so if no connection, we simply consider \emptyset as the expression between two states

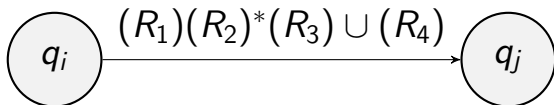
GNFA \rightarrow regular expression I

- q_{rip} is the state being removed



- The new regular expression between q_i and q_j is

GNFA \rightarrow regular expression II



- In our example
3-state DFA \rightarrow 5-state GNFA \rightarrow 4-state $\dots \rightarrow$
2-state GNFA \rightarrow regular expression
- In the procedure any any (i, j) related to q_{rip} considered
- Algorithm: convert(G)
 - 1 k : # of G
 - 2 If $k = 2$

GNFA \rightarrow regular expression III

return R between q_s and q_a

- 3 If $k > 2$, choose any $q_{rip} \in Q \setminus \{q_s, q_a\}$ for removal

$$Q' = Q - \{q_{rip}\}$$

$$\forall q_i \in Q' - \{q_{accept}\}, q_j \in Q' - \{q_{start}\}$$

$$\delta'(q_i, q_j) = R_1 R_2^* R_3 \cup R_4,$$

where

$$R_1 = \delta(q_i, q_{rip}), R_2 = \delta(q_{rip}, q_{rip}),$$

$$R_3 = \delta(q_{rip}, q_j), R_4 = \delta(q_i, q_j)$$

GNFA \rightarrow regular expression IV

- ④ Run $\text{convert}(G')$, where

$$G' = (Q', \Sigma, \delta', q_s, q_a)$$

- You can see we have a recursive setting
The process stops when $k = 2$
- Why in the textbook we modify DFA to GNFA?
Is it ok to use NFA?
Seems ok??