

The use of pumping lemma I

- Recall that from previous examples we need to choose a string s
- Some are confused about why we need to choose s but not p
- Here we will discuss in detail about how we really use the pumping lemma

The use of pumping lemma II

- Let's start from the following statement

$$\forall n, 1 + \dots + n = \frac{n(n+1)}{2}$$

What is the opposite statement?

$$\exists n \text{ such that } 1 + \dots + n \neq \frac{n(n+1)}{2}$$

- Formally, in the pumping lemma
regular \Rightarrow some properties

The use of pumping lemma III

“Some properties” are in fact

$$\exists p, \{ \forall s \in A, |s| \geq p [\exists x, y, z \text{ such that } s = xyz \text{ and } (xy^i z \in A, \forall i \geq 0, \text{ and } |y| > 0 \text{ and } |xy| \leq p)] \}$$

- To do the proof by contradiction, we need the opposite statement of the right-hand side

$$\forall p, \{ \exists s \in A, |s| \geq p, [\forall x, y, z \text{ (opposite of } (s = xyz \text{ and } (xy^i z \in A, \forall i \geq 0, \text{ and } |y| > 0 \text{ and } |xy| \leq p)))] \}$$

The use of pumping lemma IV

- What is the opposite of

$$s = xyz \text{ and } xy^i z \in A, \forall i \geq 0, \text{ and} \\ |y| > 0 \text{ and } |xy| \leq p?$$

- Various equivalent forms are possible. The one we consider is

$$(s = xyz \text{ and } |y| > 0 \text{ and } |xy| \leq p) \\ \rightarrow \exists i \geq 0, xy^i z \notin A$$

The use of pumping lemma V

- Therefore, the opposite statement of the right-hand side that we really use is

$$\begin{aligned} \forall p, \{ \exists s \in A, |s| \geq p, [\forall x, y, z \\ ((s = xyz \text{ and } |y| > 0 \text{ and } |xy| \leq p) \\ \rightarrow \exists i \geq 0, xy^i z \notin A)] \} \end{aligned} \quad (1)$$

- Note that the opposite of

$A \& B$

The use of pumping lemma VI

is

$$A \rightarrow \neg B$$

See the truth table

A	B	$A \& B$	$\neg(A \& B)$	$\neg B$	$A \rightarrow \neg B$
0	0	0	1	1	1
0	1	0	1	0	1
1	0	0	1	1	1
1	1	1	0	0	0

- To prove (1), the “exists” part is important
- You can see that we need to choose s and find an i

The use of pumping lemma VII

- In a sense we guess s or i to see if the statement can be proved.

If not, we may try other choices

- About

$$\forall x, y, z, \dots$$

in (1) you can see in examples that we go through all possible cases of x, y, z

Example 1.75 I

- $F = \{ww \mid w \in \{0, 1\}^*\}$ not regular
- We choose

$$s = 0^p 1 0^p 1 \in F$$

- If

$$s = xyz, |xy| \leq p, |y| > 0,$$

then

$$y = 0 \dots 0$$

and thus

$$xy^2z = 0 \dots 0 1 0^p 1 \neq ww$$

Example 1.75 II

- What if we say there exists $i = 3$ such that

$$xy^i z = xy^3 z = 0 \dots 010^p 1 \neq ww?$$

The proof is still correct.

- Note that we only need to **find an i** such that $xy^i z$ is not in the language

Example 1.76 I

- $D = \{1^{n^2} \mid n \geq 0\}$ not regular
- For this language we have

$$n = 0, 1^0 = \epsilon$$

$$n = 1, 1^1 = 1$$

$$n = 2, 1^4 = 1111$$

$$n = 3, 1^9 = 111111111$$

- We choose

$$s = 1^{p^2} \in D$$

Example 1.76 II

- If

$$s = xyz, |xy| \leq p, |y| > 0$$

then

$$p^2 < |xy^2z| \leq p^2 + p < (p + 1)^2$$

and therefore

$$xy^2z \notin D$$

Example 1.76 III

- What if we consider $i = 0$? It seems that

$$(p - 1)^2 < p^2 - p \leq |xy^0z| < p^2$$

can also give us xy^0z not in the language. However,

$$(p - 1)^2 < p^2 - p$$

may not hold at $p = 1$.

- Note that p is any positive integer now
- However, we still have the proof because as we said, one i is enough