Small-o I

- Two different concepts:

  O: no more than something

  o: less than something
- Definition

$$f(n) = o(g(n))$$

  if

$$\lim_{n \to \infty} \frac{f(n)}{g(n)} = 0.$$

Small-o II

- The definition of this limit:

$$\forall c > 0, \exists n_0, \forall n \geq n_0, \frac{f(n)}{g(n)} \leq c.$$

- O versus o:

$$\exists c > 0, \exists n_0, \forall n \geq n_0, f(n) \leq cg(n)$$
$$\forall c > 0, \exists n_0, \forall n \geq n_0, f(n) \leq cg(n)$$

The $\forall c$ causes o to be something smaller

Small-o III

- $\sqrt{n} = o(n)$

$$\lim_{n\to\infty} \frac{\sqrt{n}}{n} = \lim_{n\to\infty} \frac{1}{\sqrt{n}} = 0$$

- $f(n) \neq o(f(n))$

$$\lim_{n\to\infty} \frac{f(n)}{f(n)} = 1 \neq 0$$

Example: $A = \{0^k 1^k \mid k \geq 0\}$ I

- Let's count the number of steps in the algorithm discussed before
- Check if the input is

$$0...01...1$$

  This takes $O(n)$
- Move back: $O(n)$
- Cross off each 0 and 1: $O(n)$

  How many such crosses: $n/2$

$$n/2 \times O(n) = O(n^2)$$

Example: $A = \{0^k 1^k \mid k \geq 0\}$ II

- Accept or not?

  $O(n)$ to go through from beginning to end
- Total:

$$O(n) + O(n^2) + O(n) = O(n^2)$$

Time complexity class I

- Definition:

$$\begin{aligned}&\text{TIME}(t(n))\\ \equiv&\{L \mid L \text{ a language decided by an } O(t(n)) \text{ TM}\}\end{aligned}$$

- We have

$$\{0^k 1^k \mid k \geq 0\} \in \text{TIME}(n^2)$$

Can we make it faster?

New Algorithm for $A = \{0^k1^k \mid k \geq 0\}$ I

- The procedure: cross off every other 0 and 1

  $\underline{0}0\underline{0}0\underline{0}\underline{1}1\underline{1}1\underline{1}$

  $\underline{0}0\underline{1}1$

  $\underline{0}\underline{1}$

  ϵ

  key: length of the string left must be always even
- A failed algorithm

  $\underline{0}0\underline{0}0\underline{1}1$

  001
- Algorithm

New Algorithm for $A = \{0^k 1^k \mid k \geq 0\}$ II

1. check 0...0 1...1
2. repeat if not empty
   total # 0 & 1: odd $\Rightarrow$ reject
   cross off every other 0 and 1
3. no 0 & 1 remain, accept

- If 13 "0" $\Rightarrow$ 6 "0" $\Rightarrow$ 3 "0" $\Rightarrow$ 1 "0"
  $1 + \log_2 n$ iterations
- Each iteration: $O(n)$ operations
  Note that length of tape contents is still n as we only "mark" elements
- Total cost: $O(n \log n)$

New Algorithm for $A = \{0^k1^k \mid k \geq 0\}$ III

- Therefore

$$\{0^k1^k \mid k \geq 0\} \in \mathrm{TIME}(n \log n)$$

- Can we do better? no
- Any language decided in $o(n \log n)$ on a single-tape TM $\Rightarrow$ regular (not proved here)
- But we know that

$$\{0^k1^k \mid k \geq 0\}$$

is not regular

Using two-tape TM for $\{0^k 1^k \mid k \geq 0\}$ I

- We can have an $O(n)$ procedure
  1. check 0...0 1...1
  2. copy 0 to the second tape
     find the first 1
  3. sequentially cut 1 and 0
     if no “0” reject
  4. if “1” left, reject
     otherwise, accept
- Each step $O(n)$