Decidability and CFL I

- CFG

\[ A_{CFG} = \{ \langle G, w \rangle \mid G : CFG, \text{ generates } w \} \]

- We prove that \( A_{CFG} \) is decidable
- But an issue is the \( \infty \) possible derivations of a CFG
- For example,
  \[ A \to B, B \to A \]

- Chomsky normal form

\[ A \to BC \]
\[ A \to a \]
Decidability and CFL II

- Any $w$, $|w| = n$, derivation in exactly $2n - 1$ steps
- If $q$: # rules, check $q^{2n-1}$ branches
- Proof
  1. Convert $G$ to Chomsky
  2. Check all $q^{2n-1}$ branches
- Results apply to PDA as well
\[ E_{CFG} = \{ \langle G \rangle \mid G : CFG, L(G) = \emptyset \} \]

- idea: bottom up setting to see if any string can be generated from the start variable. From

\[ A \rightarrow a \]

We search if \( \exists \)

\[ B \rightarrow A \]

- Proof:
1. Mark all terminals
2. Repeat until no new variables marked
   if
   
   \[ A \to U_1 \cdots U_k \]

   and
   
   all \( U_1, \ldots, U_k \) marked

   \[ \Rightarrow \text{mark } A \]
3. If start state marked, accept
   Otherwise, reject
$EQ_{CFG}$ I

$EQ_{CFG} = \{ \langle G, H \rangle \mid G, H : CFG, L(G) = L(H) \}$

- Remember that $EQ_{DFA}$ is decidable
- However, we cannot apply the same proof as CFL is not closed for $\cap$ and complementation
- It’s proved in Chapter 5 that this language is not decidable
- We do not discuss details
CFL decidable I

- This question different from $A_{CFG}$ decidable or not
- How about converting PDA to a TM?
- For nondeterministic PDA we can do NTM
- But nondeterministic PDA may have $\infty$-long branches
  TM runs forever
- So converting PDA to TM does not really work
- A proof that works:
  Find grammar $G$ for this CFL
  Run TM for $\langle G, w \rangle$ by using $A_{CFG}$
Classes of languages I

- Fig 4.10

Diagram showing the hierarchy of classes of languages:
- Regular
- Context-free
- Decidable
- Turing-recognizable