

# Regular Operations I

- Regular operations can be used to study whether languages are regular or not
- That is, we aim to check for a given language, whether there are finite automata to recognize it or not
- Three definitions
  - $A, B$  are given languages
    - union

$$A \cup B$$

# Regular Operations II

- concatenation

$$A \circ B = \{xy \mid x \in A, y \in B\}$$

- star:

$$A^* = \{x_1 \cdots x_k \mid k \geq 0, x_i \in A\}$$

# Regular Operations III

- If  $k = 0$ , what do we mean

$$x_1 \cdots x_k?$$

We define

$\epsilon$  : empty string

in this situation

- Thus

$$\epsilon \in A^*$$

# Regular Operations IV

- Example

$$\Sigma = \{a, \dots, z\}$$

$$A = \{good, bad\}$$

$$B = \{boy, girl\}$$

$$A \circ B = \{goodboy, \dots\}$$

$$A^* : \{\epsilon, good, bad, goodgood, \dots\}$$

- We say an operation  $R$  is **closed** if the following property holds

$$\text{if } x \in A, y \in A, \text{ then } xRy \in A$$

# Regular Operations V

Example:  $N = \{1, 2, \dots\}$  is closed under multiplication

- Th 1.25: regular languages are closed under the union operation

$A_1, A_2$  are regular languages  
 $\Rightarrow A_1 \cup A_2$  is regular

- Proof

# Regular Operations VI

Assume we are given two automata

$$M_1 = (Q_1, \Sigma, \delta_1, q_1, F_1)$$

$$M_2 = (Q_2, \Sigma, \delta_2, q_2, F_2)$$

# Regular Operations VII

- Construct a new machine

$$M = (Q, \Sigma, \delta, q, F)$$

$$Q = \{(r_1, r_2) \mid r_1 \in Q_1, r_2 \in Q_2\}$$

$$\delta((r_1, r_2), a) = (\delta_1(r_1, a), \delta_2(r_2, a))$$

$$q_0 = (q_1, q_2)$$

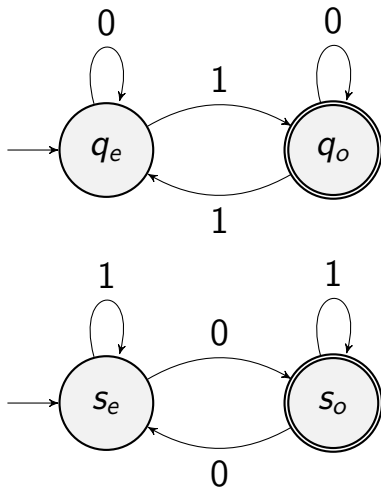
$$F = \{(r_1, r_2) \mid r_1 \in F_1 \text{ or } r_2 \in F_2\}$$

- Example: combining

$$\{w \mid w \text{ has an odd \# 1's}\} \cup$$

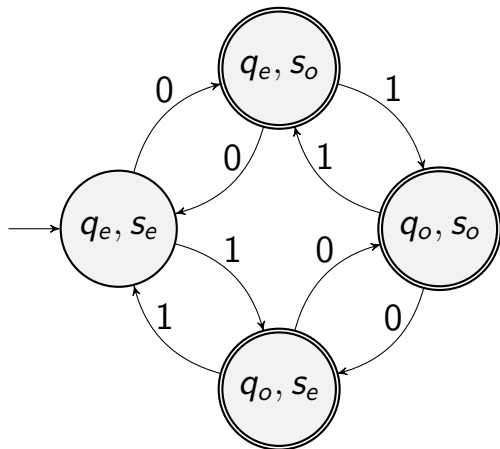
$$\{w \mid w \text{ has an odd \# 0's}\}$$

# Regular Operations VIII





# Regular Operations IX



- Is this proof rigourously enough ?

# Regular Operations X

A formal proof should be done by induction. But we don't provide it here

- Th 1.26: closed under concatenation

If  $A, B$  are regular, then  $A \circ B$  is regular

But the proof is not easy

It's unclear where to break the input

- To easily do the proof, we introduce a new technique called **nondeterminism**