

Data Structures and Algorithms

(資料結構與演算法)

Lecture 1: Algorithm

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Roadmap

1 the one where it all began

Lecture 1: Algorithm

- definition of algorithm
- pseudo code of algorithm
- criteria of algorithm
- correctness proof of algorithm

2 the data structures awaken

3 fantastic trees and where to find them

4 the search revolutions

5 sorting: the final frontier

definition of algorithm

Name Origin of Algorithm

Muhammad ibn Mūsā al-Ḳwārizmī on a Soviet Union stamp

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algorithm

- named after **al-Ḳwārizmī** (780–850), Persian mathematician and father of algebra
- algebra: **rules to calculate with symbols**
- algorithm: **instructions to compute with variables**

algorithm: **recipe-like instructions** for **computing**

Recipe for Cooking Dish

a recipe for hamburger on Wikibooks

figure by Gentgeen,

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Cookbook:Hamburger - Wikibooks

http://en.wikibooks.org/wiki/Cookbook:Hamburger

Cookbook:Hamburger

From Wikibooks

Cookbook | Recipe Index | Meat recipes

A **hamburger** (see [hamburger](#)), **hamburg**, or in the United Kingdom, a **hamburger**) is a variant on a sandwich involving a patty of ground meat that is almost always beef.

Ingredients

- 50g (1.1 lb) minced (ground) beef
- lettuce and onions (optional)
- cheese (optional)
- salad (lettuce, spinach, alfalfa sprouts, tomatoes, onion etc., optional)
- 1 hamburger bun for each burger

Procedure

1. Add the beef to a food processor for approximately 10 seconds.
2. Next add your lettuce and/or onion to taste. Depending on the quantity of your local beef, you may wish to add some beef stock to improve the flavor.
3. Mix at the final processing for another 10 seconds or until fully mixed.
4. If you bought the beef already ground, make sure you mix in your seasonings well. You may wish to add garlic, onion flakes, soy sauce, Worcestershire sauce and/or olive oil. You can also add 2 tsp of your favorite hot sauce for some kick. Note: If you add any liquids, the ground beef will then require not the extra price when forming patties.
5. Remove the beef from the food processor and shape by hand into burgers. You should get between 4-6 burgers from 50g (1.1 lb) of beef.
6. The burgers can be fried (about 5 min on each side for burgers which aren't too thick), grilled (same time as for frying), or baked.
7. Remove your burgers are fully cooked through before serving. If your burgers are quite thick, or if you use onions, you can cut one open to ensure the middle is not browned. If the middle are red, there is a chance that the meat is not fully cooked.
8. Serve each burger on a bun (optional, see [pretzel roll](#)), optionally with ketchup, sliced pickles, ketchup, mayonnaise, mustard, ranch dressing, cheese, lettuce, tomato and/or onion.

• For further serving suggestions, see the Wikipedia article on [hamburgers](#).

Notes, tips and variations

- You can use almost any type of minced (ground) meat to make hamburgers, including pork, chicken, turkey, lamb, bacon, venison, coriander, or even a meat substitute such as Quorn.
- If your burgers fall apart, adding an egg yolk will help keep it together. Heating lean ground beef will also help.
- You may wish to experiment with including cheese in the center of your burger before cooking.
- Spices which can work well on hamburgers include black pepper, chili powder (fresh or ground), Worcestershire Sauce and soy sauce. Experiment to find good combinations.
- Burgers can be cooked in a grill. Smoked burgers will appear red and placed on the outside, but browned on the inside.
- Smoking a burger before grilling it in an outdoor way is used in the Howard's patent.
- **Warning:** Making meat and onion together in a bowl and mixing by hand will the onions are distributed may produce better results. This will also stop your burgers from falling apart.

Links

Retrieved from "http://en.wikibooks.org/wiki/Cookbook:Hamburger"

Categories: Recipes | Sandwich recipes | Beef recipes | Paasted recipes

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Cookbook:Hamburger

Category:	Meat recipes
Source:	Public
Energy:	Hamburger 360 Cal / 1463 kJ Cheeseburger 790 Cal / 3303 kJ
Time:	10 minutes
Difficulty:	● ● ●

recipe

Wikipedia: a set of *instructions* that describes how to prepare or make something, especially a dish of prepared food

recipe: instructions to complete a (cooking) task

1 of 1

11/06/09 11:30 PM

Sheet Music for Playing Instrument



first page of the manuscript of

Bach's lute suite in G minor

figure licensed

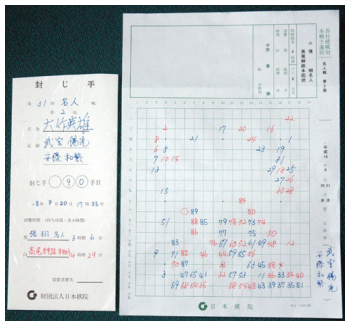
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sheet music

Wikipedia: *handwritten or printed form of musical notation ... to indicate the pitches, rhythms or chords of a song*

sheet music: instructions to play instrument (well)

Kifu for Playing Go



a Japanese kifu

figure by Velobici,

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kifu

go game record **of steps** that describe
how the game had been played

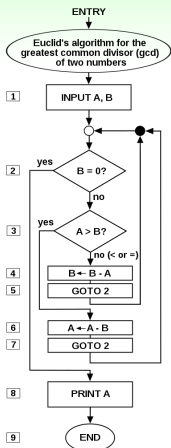
kifu: **instructions** to **mimic/learn** to play go (professionally)

Algorithm for Computing

flowchart of Euclid's algorithm for calculating the greatest common divisor (g.c.d.) of two numbers

figure by Somepics,

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algorithm

Wikipedia: *algorithm is a finite sequence of well-defined, computer-implementable instructions*, typically to solve a class of problems or to perform a computation

algorithm ~ computing recipe:

(computable) instructions to solve a computing task efficiently/correctly

Fun Time

Which of the following in the kitchen is the best metaphor for an algorithm?

- ① recipe
- ② chef
- ③ garbage
- ④ meat

Fun Time

Which of the following in the kitchen is the best metaphor for an algorithm?

- ① recipe
- ② chef
- ③ garbage
- ④ meat

Reference Answer: ①

algorithm ~ computing recipe:

(computable) instructions to solve a computing task efficiently/correctly

pseudo code of algorithm

Pseudo Code for GETMININDEX

C Version

```
/* return index to min. element
   in arr[0] ... arr[len-1] */
int getMinIndex
    (int arr[], int len){
    int i;
    int m=0;
    for(i=0;i<len;i++){
        if (arr[m] > arr[i]){
            m = i;
        }
    }
    return m;
}
```

Pseudo Code Version

GET-MIN-INDEX(A)

```
1  m = 1
2  for i = 2 to A.length
3      // update if i-th element smaller
4      if A[m] > A[i]
5          m = i
6  return m
```

pseudo code: **spoken language** of programming

Bad Pseudo Code: Too Detailed

Unnecessarily Detailed

GET-MIN-INDEX(A)

```
1  $m = 1$ 
2 for  $i = 2$  to  $A.length$ 
3     // update if  $i$ -th element smaller
4      $A_m = A[m]$ 
5      $A_i = A[i]$ 
6     if  $A_m > A_i$ 
7          $m = i$ 
8     else
9          $m = m$ 
10 return  $m$ 
```

Concise

GET-MIN-INDEX(A)

```
1  $m = 1$ 
2 for  $i = 2$  to  $A.length$ 
3     // update if  $i$ -th element smaller
4     if  $A[m] > A[i]$ 
5          $m = i$ 
6 return  $m$ 
```

goal of pseudo code: communicate efficiently

Bad Pseudo Code: Too Mysterious

Unnecessarily Mysterious

GET-MIN-INDEX(A)

```
1  $x = 1$ 
2 for  $xx = 2$  to  $A.length$ 
3
4     if  $A[x] > A[xx]$ 
5          $xx = x$ 
6 return  $xx$ 
```

Clear

GET-MIN-INDEX(A)

```
1  $m = 1$  // store current min. index
2 for  $i = 2$  to  $A.length$ 
3     // update if  $i$ -th element smaller
4     if  $A[m] > A[i]$ 
5          $m = i$ 
6 return  $m$ 
```

goal of pseudo code: communicate correctly

Bad Pseudo Code: Too Abstract

Unnecessarily Abstract

GET-MIN-INDEX(A)

- 1 $m = 1$ // store current min. index
- 2 **run a loop through A**
that updates m in every iteration
- 3 **return m**

Concrete

GET-MIN-INDEX(A)

- 1 $m = 1$ // store current min. index
- 2 **for** $i = 2$ **to** $A.length$
- 3 // update if i -th element smaller
- 4 **if** $A[m] > A[i]$
- 5 $m = i$
- 6 **return** m

goal of pseudo code: **communicate effectively**

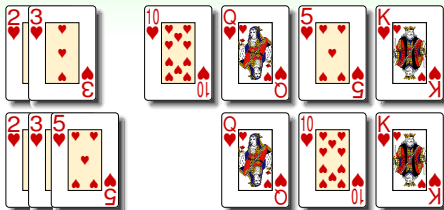
From GET-MIN-INDEX to SELECTION-SORT

GET-MIN-INDEX(A, ℓ, r)

```

1   $m = \ell$  // store current min. index
2  for  $i = \ell + 1$  to  $r$ 
3      // update if  $i$ -th element smaller
4      if  $A[m] > A[i]$ 
5           $m = i$ 
6  return  $m$ 

```



Good Pseudo Code

- modularize, just like coding
- depends on speaker/listener
- usually no formal definition

SELECTION-SORT(A)

```

1  for  $i = 1$  to  $A.length$ 
2       $m = \text{GET-MIN-INDEX}(A, i, A.length)$ 
3      SWAP( $A[i], A[m]$ )
4  return  $A$  // which has been sorted in place

```

follow any textbook if you really need a definition

Fun Time

Which of the following can be used to describe good pseudo code?

- ① clear
- ② concise
- ③ concrete
- ④ all of the above

Fun Time

Which of the following can be used to describe good pseudo code?

- ① clear
- ② concise
- ③ concrete
- ④ all of the above

Reference Answer: ④

Have fun communicating with other programmers using good pseudo code! :-)

criteria of algorithm

Criteria of Recipe



figure by Larry, licensed under CC BY-NC-ND 2.0 via Flickr

Cocktail Recipe: Screwdriver (from Wikipedia)

inputs: 5 cl vodka, 10 cl orange juice

- 1 mix inputs in a highball glass with ice
- 2 garnish with orange slice and serve

output: a glass of delicious cocktail

- input:
ingredients
- definiteness:
clear instructions
- effectiveness:
feasible instructions
- finiteness:
completable instructions
- output:
delicious drink

algorithm ~ recipe: same five criteria for algorithm

(Knuth, The Art of Computer Programming)

Input of Algorithm

... quantities which are given to it initially before the algorithm begins.
These inputs are taken from **specified sets of objects**. (Knuth, TAOCP)

GET-MIN-INDEX(A)

```
1   $m = 1$  // store current min. index
2  for  $i = 2$  to  $A.length$ 
3      // update if  $i$ -th element smaller
4      if  $A[m] > A[i]$ 
5           $m = i$ 
6  return  $m$ 
```

one algorithm, many uses (on different **legal inputs**)

Definiteness of Algorithm

*Each step of an algorithm must be precisely defined; the actions to be carried out must be **rigorously & unambiguously specified**.* (Knuth, TAOCP)

Clear

GET-MIN-INDEX(A)

```
1  $m = 1$  // store current min. index
2 for  $i = 2$  to  $A.length$ 
3     // update if  $i$ -th element smaller
4     if  $A[m] > A[i]$ 
5          $m = i$ 
6 return  $m$ 
```

Ambiguous

GET-ZERO-INDEX(A)

```
1
2 for  $i = 1$  to  $A.length$ 
3
4     if  $A[m]$  is almost zero
5         return  $m$ 
6 // what to return here?
```

definiteness: **clarity** of algorithm

Effectiveness of Algorithm

... all of the operations to be performed in the algorithm must be sufficiently basic that they can *in principle be done exactly and in a finite length of time* by a man using paper and pencil. (Knuth, TAOCP)

Effective

GET-MIN-INDEX(A)

```
1  $m = 1$  // store current min. index
2 for  $i = 2$  to  $A.length$ 
3     // update if  $i$ -th element smaller
4     if  $A[m] > A[i]$ 
5          $m = i$ 
6 return  $m$ 
```

Ineffective

GET-SOFT-MIN(A)

```
1  $s = 0$  // sum of exponentiated values
2 for  $i = 1$  to  $A.length$ 
3      $s = s + \exp(-A[i] \cdot 1126)$ 
4
5
6 return  $-\log(s)/1126$ 
```

floating point errors may make some steps ineffective on some computers

Finiteness of Algorithm

*An algorithm must always terminate after a finite number of steps ...
a very finite number, a reasonable number.*

(Knuth, TAOCP)

GET-MIN-INDEX(A)

```
1   $m = 1$  // store current min. index
2  for  $i = 2$  to  $A.length$ 
3      // update if  $i$ -th element smaller
4      if  $A[m] > A[i]$ 
5           $m = i$ 
6  return  $m$ 
```

finiteness (& efficiency): often requiring analysis for sophisticated algorithms (to be taught later)

Output of Algorithm

... quantities which have a *specified relation* to the inputs (Knuth, TAOCP)

```
GET-MIN-INDEX(A)
```

```
1 m = 1 // store current min. index  
2 for i = 2 to A.length  
3     // update if i-th element smaller  
4     if  $A[m] > A[i]$   
5         m = i  
6 return m
```

output (*correctness*): needs *proving*
with respect to requirements

Fun Time

What best describes the input/output relationship of the selection sort algorithm below?

```
SELECTION-SORT(A)
```

```
1 for i = 1 to A.length  
2     m = GET-MIN-INDEX(A, i, A.length)  
3     SWAP(A[i], A[m])  
4 return A // which has been sorted in place
```

- 1 input: an ascending array;
output: the same array
sorted in descending order
- 2 input: an arbitrary array;
output: the same array
sorted in descending order
- 3 input: an arbitrary array;
output: the same array
sorted in ascending order
- 4 none of the other choices

Fun Time

What best describes the input/output relationship of the selection sort algorithm below?

```
SELECTION-SORT(A)
```

```
1 for i = 1 to A.length
2     m = GET-MIN-INDEX(A, i, A.length)
3     SWAP(A[i], A[m])
4 return A // which has been sorted in place
```

- 1 input: an ascending array;
output: the same array sorted in descending order
- 2 input: an arbitrary array;
output: the same array sorted in descending order
- 3 input: an arbitrary array;
output: the same array sorted in ascending order
- 4 none of the other choices

Reference Answer: 3

The selection sort algorithm re-arranges an arbitrary array into ascending order.

correctness proof of algorithm

Claim



figure by Nick Youngson, licensed CC BY-SA 3.0 via Picpedia.Org

GET-MIN-INDEX(A)

```
1  $m = 1$  // store current min. index
2 for  $i = 2$  to  $A.length$ 
3     // update if  $i$ -th element smaller
4     if  $A[m] > A[i]$ 
5          $m = i$ 
6 return  $m$ 
```

Correctness of GET-MIN-INDEX

Upon exiting GET-MIN-INDEX(A),

$$A[m] = \min_{1 \leq j \leq n} A[j]$$

with $n = A.length$

claim: math. statement that **declares correctness**

Invariant



invariants when constructing fractals
figures by Johannes Rössel,

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GET-MIN-INDEX(A)

```

1   $m = 1$  // store current min. index
2  for  $i = 2$  to  $A.length$ 
3      // update if  $i$ -th element smaller
4      if  $A[m] > A[i]$ 
5           $m = i$ 
6  return  $m$ 

```

Correctness of GET-MIN-INDEX

Upon exiting GET-MIN-INDEX(A),

$$A[m] = \min_{1 \leq j \leq n} A[j]$$

with $n = A.length$



Invariant within GET-MIN-INDEX

Upon finishing the loop with $i = k$,
denote m by m_k ,

$$A[m_k] \leq A[j] \text{ for } j = 1, 2, \dots, k$$

(loop) invariant: property that algorithm maintains

Proof of Loop Invariant

Mathematical Induction

Base

when $i = 2$, invariant true because

...

- assume **invariant true** for $i = t - 1$
- when $i = t$,
 - if $A[m_{t-1}] > A[t] \Rightarrow m_t = t$

$$A[m_t] = A[t] \leq A[t] \\ < A[m_{t-1}] \leq A[j] \text{ for } j < t$$

- if $A[m_{t-1}] \leq A[t] \Rightarrow m_t = m_{t-1}$

$$A[m_t] = A[m_{t-1}] \leq A[t] \\ = A[m_{t-1}] \leq A[j] \text{ for } j < t$$

—by mathematical induction,
invariant true for $i = 2, 3, \dots, k$

GET-MIN-INDEX(A)

```

1  m = 1 // store current min. index
2  for i = 2 to A.length
3      // update if i-th element smaller
4      if A[m] > A[i]
5          m = i
6  return m
```

Correctness of GET-MIN-INDEX



Invariant within GET-MIN-INDEX

Upon **finishing the loop** with $i = k$,
denote m by m_k ,

$$A[m_k] \leq A[j] \text{ for } j = 1, 2, \dots, k$$

⇒

proof of (loop) invariants ⇒ correctness claim of algorithm

Fun Time

Which of the following is a loop invariant to selection sort?

```
SELECTION-SORT(A)
```

```
1  for i = 1 to A.length  
2      m = GET-MIN-INDEX(A, i, A.length)  
3      SWAP(A[i], A[m])  
4  return A // which has been sorted in place
```

- 1 Upon finishing the loop with $i = k$, $A[1] \geq A[2] \geq \dots \geq A[k]$.
- 2 Upon finishing the loop with $i = k$, $A[1] \leq A[2] \leq \dots \leq A[k]$.
- 3 Upon finishing the loop with $i = k$, $A[k + 1] \geq \dots \geq A[A.length]$.
- 4 Upon finishing the loop with $i = k$, $A[k + 1] \leq \dots \leq A[A.length]$.

Fun Time

Which of the following is a loop invariant to selection sort?

```
SELECTION-SORT(A)
```

```
1  for  $i = 1$  to  $A.length$   
2       $m = \text{GET-MIN-INDEX}(A, i, A.length)$   
3       $\text{SWAP}(A[i], A[m])$   
4  return  $A$  // which has been sorted in place
```

- 1 Upon finishing the loop with $i = k$, $A[1] \geq A[2] \geq \dots \geq A[k]$.
- 2 Upon finishing the loop with $i = k$, $A[1] \leq A[2] \leq \dots \leq A[k]$.
- 3 Upon finishing the loop with $i = k$, $A[k + 1] \geq \dots \geq A[A.length]$.
- 4 Upon finishing the loop with $i = k$, $A[k + 1] \leq \dots \leq A[A.length]$.

Reference Answer: ②

The selection sort algorithm essentially picks the smallest element, the 2nd-smallest, and so on, and locate them orderly. You can prove the loop invariant by mathematical induction.

Summary

Lecture 1: Algorithm

- definition of algorithm
instructions to complete a task by computer
 - pseudo code of algorithm
communicate alg. efficiently/correctly/effectively
 - criteria of algorithm
input, definite, effective, finite, output
 - correctness proof of algorithm
from (loop) invariants to claims
- **next: 'data structures' and their connections to algorithms**