



Computer Organization & Assembly Languages

Data Transfers, Addressing & Arithmetic

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Adapted from the slides prepared by Kip Irvine for the book,
Assembly Language for Intel-Based Computers, 5th Ed.



Chapter Overview

- **Data Transfer Instructions**
- Addition and Subtraction
- Data-Related Operators and Directives
- Indirect Addressing
- JMP and LOOP Instructions



Data Transfer Instructions

- Operand Types
- Instruction Operand Notation
- Direct Memory Operands
- MOV Instruction
- Zero & Sign Extension
- XCHG Instruction
- Direct-Offset Instructions



Operand Types

- Three basic types of operands:
 - Immediate – a constant integer (8, 16, or 32 bits)
 - value is encoded within the instruction
 - Register – the name of a register
 - register name is converted to a number and encoded within the instruction
 - Memory – reference to a location in memory
 - memory address is encoded within the instruction, or a register holds the address of a memory location



Instruction Operand Notation

Operand	Description
<i>r8</i>	8-bit general-purpose register: AH, AL, BH, BL, CH, CL, DH, DL
<i>r16</i>	16-bit general-purpose register: AX, BX, CX, DX, SI, DI, SP, BP
<i>r32</i>	32-bit general-purpose register: EAX, EBX, ECX, EDX, ESI, EDI, ESP, EBP
<i>reg</i>	any general-purpose register
<i>sreg</i>	16-bit segment register: CS, DS, SS, ES, FS, GS
<i>imm</i>	8-, 16-, or 32-bit immediate value
<i>imm8</i>	8-bit immediate byte value
<i>imm16</i>	16-bit immediate word value
<i>imm32</i>	32-bit immediate doubleword value
<i>r/m8</i>	8-bit operand which can be an 8-bit general register or memory byte
<i>r/m16</i>	16-bit operand which can be a 16-bit general register or memory word
<i>r/m32</i>	32-bit operand which can be a 32-bit general register or memory doubleword
<i>mem</i>	an 8-, 16-, or 32-bit memory operand



Direct Memory Operands

- A direct memory operand is a named reference to storage in memory
- The named reference (label) is automatically dereferenced by the assembler

```
.data
var1 BYTE 10h
.code
mov al,var1           ; AL = 10h
mov al,[var1]        ; AL = 10h
```

alternate format





MOV Instruction

- Move from source to destination. Syntax:

MOV destination, source

- Source and destination have the same size
- No more than one memory operand permitted
- CS, EIP, and IP cannot be the destination
- No immediate to segment moves

```
.data
count BYTE 100
wVal  WORD 2
.code
    mov bl, count
    mov ax, wVal
    mov count, al

    mov al, wVal      ; error
    mov ax, count    ; error
    mov eax, count   ; error
```



Your Turn . . .

Explain why each of the following MOV statements are invalid:

```
.data
bVal  BYTE  100
bVal2 BYTE  ?
wVal  WORD  2
dVal  DWORD 5
.code
    mov ds,45           immediate move to DS not permitted
    mov esi,wVal       size mismatch
    mov eip,dVal       EIP cannot be the destination
    mov 25,bVal        immediate value cannot be destination
    mov bVal2,bVal     memory-to-memory move not permitted
```



Memory to Memory

```
.data
```

```
var1 WORD ?
```

```
var2 WORD ?
```

```
.code
```

```
mov ax, var1
```

```
mov var2, ax
```



Copy Smaller to Larger

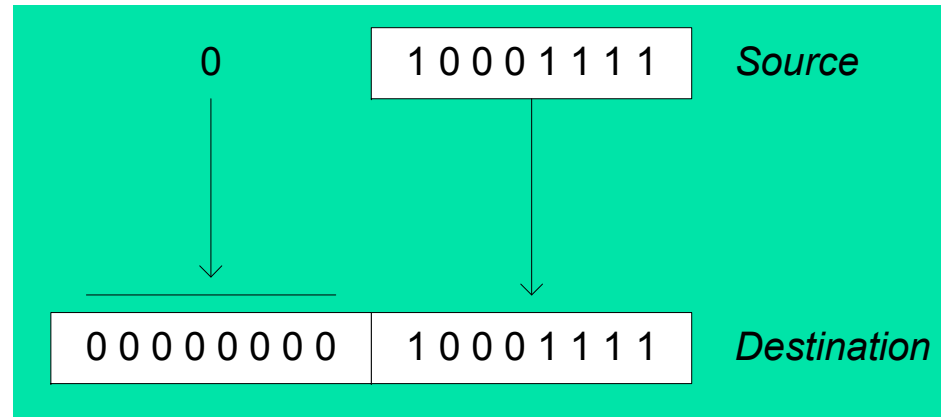
```
.data  
count WORD 1
```

```
.code  
mov ecx, 0  
mov cx, count
```

```
.data  
signedVal SWORD -16 ; FFF0h  
.code  
mov ecx, 0 ; mov ecx, 0FFFFFFFFh  
mov cx, signedVal
```

Zero Extension

When you copy a smaller value into a larger destination, the MOVZX instruction fills (extends) the upper half of the destination with zeros.

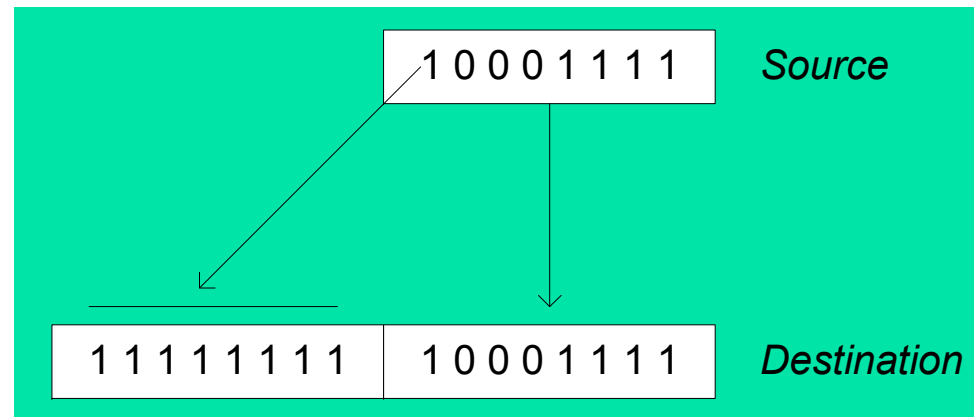


```
mov bl,10001111b  
movzx ax,bl ; zero-extension
```

The destination must be a register.

Sign Extension

The MOVSX instruction fills the upper half of the destination with a copy of the source operand's sign bit.



```
mov bl,10001111b  
movsx ax,bl           ; sign extension
```

The destination must be a register.



MOVZX & MOVSX

From a smaller location to a larger one

```
mov bx, 0A69Bh
```

```
movzx eax, bx ; EAX=0000A69Bh
```

```
movzx edx, bl ; EDX=0000009Bh
```

```
movzx cx, bl ; EAX=009Bh
```

```
mov bx, 0A69Bh
```

```
movsx eax, bx ; EAX=FFFA69Bh
```

```
movsx edx, bl ; EDX=FFFFFF9Bh
```

```
movsx cx, bl ; EAX=FF9Bh
```



LAHF & SAHF

```
.data
```

```
saveflags BYTE ?
```

```
.code
```

```
lahf
```

```
mov saveflags, ah
```

```
...
```

```
mov ah, saveflags
```

```
sahf
```



XCHG Instruction

XCHG exchanges the values of two operands. At least one operand must be a register. No immediate operands are permitted.

```
.data
var1 WORD 1000h
var2 WORD 2000h
.code
xchg ax,bx           ; exchange 16-bit regs
xchg ah,al           ; exchange 8-bit regs
xchg var1,bx         ; exchange mem, reg
xchg eax,ebx         ; exchange 32-bit regs

xchg var1,var2       ; error: two memory operands
```



Direct-Offset Operands

A constant offset is added to a data label to produce an effective address (EA). The address is dereferenced to get the value inside its memory location.

```
.data
arrayB BYTE 10h,20h,30h,40h
.code
mov al,arrayB+1           ; AL = 20h
mov al,[arrayB+1]        ; alternative notation
```

Q: Why doesn't `arrayB+1` produce 11h?



Direct-Offset Operands (cont.)

A constant offset is added to a data label to produce an effective address (EA). The address is dereferenced to get the value inside its memory location.

```
.data
arrayW  WORD 1000h,2000h,3000h
arrayD  DWORD 1,2,3,4
.code
mov ax, [arrayW+2]           ; AX = 2000h
mov ax, [arrayW+4]           ; AX = 3000h
mov eax, [arrayD+4]          ; EAX = 00000002h
```

```
; Will the following statements assemble?
mov ax, [arrayW-2]           ; ??
mov eax, [arrayD+16]         ; ??
```

What will happen when they run?



Your Turn. . .

Write a program that rearranges the values of three doubleword values in the following array as: 3, 1, 2.

```
.data  
arrayD DWORD 1,2,3
```

- Step1: copy the first value into EAX and exchange it with the value in the second position.

```
mov eax,arrayD  
xchg eax,[arrayD+4]
```

- Step 2: Exchange EAX with the third array value and copy the value in EAX to the first array position.

```
xchg eax,[arrayD+8]  
mov arrayD,eax
```



Evaluate This . . .

- We want to write a program that adds the following three bytes:

```
.data  
myBytes BYTE 80h,66h,0A5h
```

- What is your evaluation of the following code?

```
mov al,myBytes  
add al,[myBytes+1]  
add al,[myBytes+2]
```

- What is your evaluation of the following code?

```
mov ax,myBytes  
add ax,[myBytes+1]  
add ax,[myBytes+2]
```

- Any other possibilities?



Evaluate This . . . (cont)

```
.data  
myBytes BYTE 80h,66h,0A5h
```

- How about the following code. Is anything missing?

```
movzx ax,myBytes  
mov    bl,[myBytes+1]  
add    ax,bx  
mov    bl,[myBytes+2]  
add    ax,bx                ; AX = sum
```

Yes: Move zero to BX before the MOVZX instruction.



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- **Addition and Subtraction**
- Data-Related Operators and Directives
- Indirect Addressing
- JMP and LOOP Instructions



Addition and Subtraction

- INC and DEC Instructions
- ADD and SUB Instructions
- NEG Instruction
- Implementing Arithmetic Expressions
- Flags Affected by Arithmetic
 - Zero
 - Sign
 - Carry
 - Overflow



INC and DEC Instructions

- Add 1, subtract 1 from destination operand
 - operand may be register or memory
- **INC *destination***
 - Logic: $destination \leftarrow destination + 1$
- **DEC *destination***
 - Logic: $destination \leftarrow destination - 1$



INC and DEC Examples

```
.data
myWord  WORD 1000h
myDword DWORD 10000000h
.code
    inc myWord           ; 1001h
    dec myWord           ; 1000h
    inc myDword          ; 10000001h

    mov ax,00FFh
    inc ax                ; AX = 0100h
    mov ax,00FFh
    inc al                ; AX = 0000h
```



Your Turn...

Show the value of the destination operand after each of the following instructions executes:

```
.data
myByte BYTE 0FFh, 0
.code
    mov al,myByte           ; AL = FFh
    mov ah,[myByte+1]      ; AH = 00h
    dec ah                 ; AH = FFh
    inc al                 ; AL = 00h
    dec ax                 ; AX = FEFF
```



ADD and SUB Instructions

- ADD destination, source
 - Logic: $destination \leftarrow destination + source$
- SUB destination, source
 - Logic: $destination \leftarrow destination - source$
- Same operand rules as for the MOV instruction



ADD and SUB Examples

```
.data
var1 DWORD 10000h
var2 DWORD 20000h
.code
; ---EAX---
    mov eax,var1      ; 00010000h
    add eax,var2      ; 00030000h
    add ax,0FFFFh     ; 0003FFFFh
    add eax,1         ; 00040000h
    sub ax,1          ; 0004FFFFh
```



NEG (negate) Instruction

Reverses the sign of an operand. Operand can be a register or memory operand.

```
.data
valB BYTE -1
valW SWORD +32767
.code
    mov al,valB           ; AL = -1
    neg al                ; AL = +1
    neg valW              ; valW = -32767
```

Suppose AX contains $-32,768$ and we apply NEG to it. Will the result be valid?



Implementing Arithmetic Expressions

HLL compilers translate mathematical expressions into assembly language. You can do it also. For example:

$$Rval = -Xval + (Yval - Zval)$$

```
Rval SDWORD ?
Xval SDWORD 26
Yval SDWORD 30
Zval SDWORD 40
.code
    mov eax,Xval
    neg eax                ; EAX = -26
    mov ebx,Yval
    sub ebx,Zval          ; EBX = -10
    add eax,ebx
    mov Rval,eax          ; -36
```



Your Turn...

Translate the following expression into assembly language.
Do not permit Xval, Yval, or Zval to be modified:

$$Rval = Xval - (-Yval + Zval)$$

Assume that all values are signed doublewords.

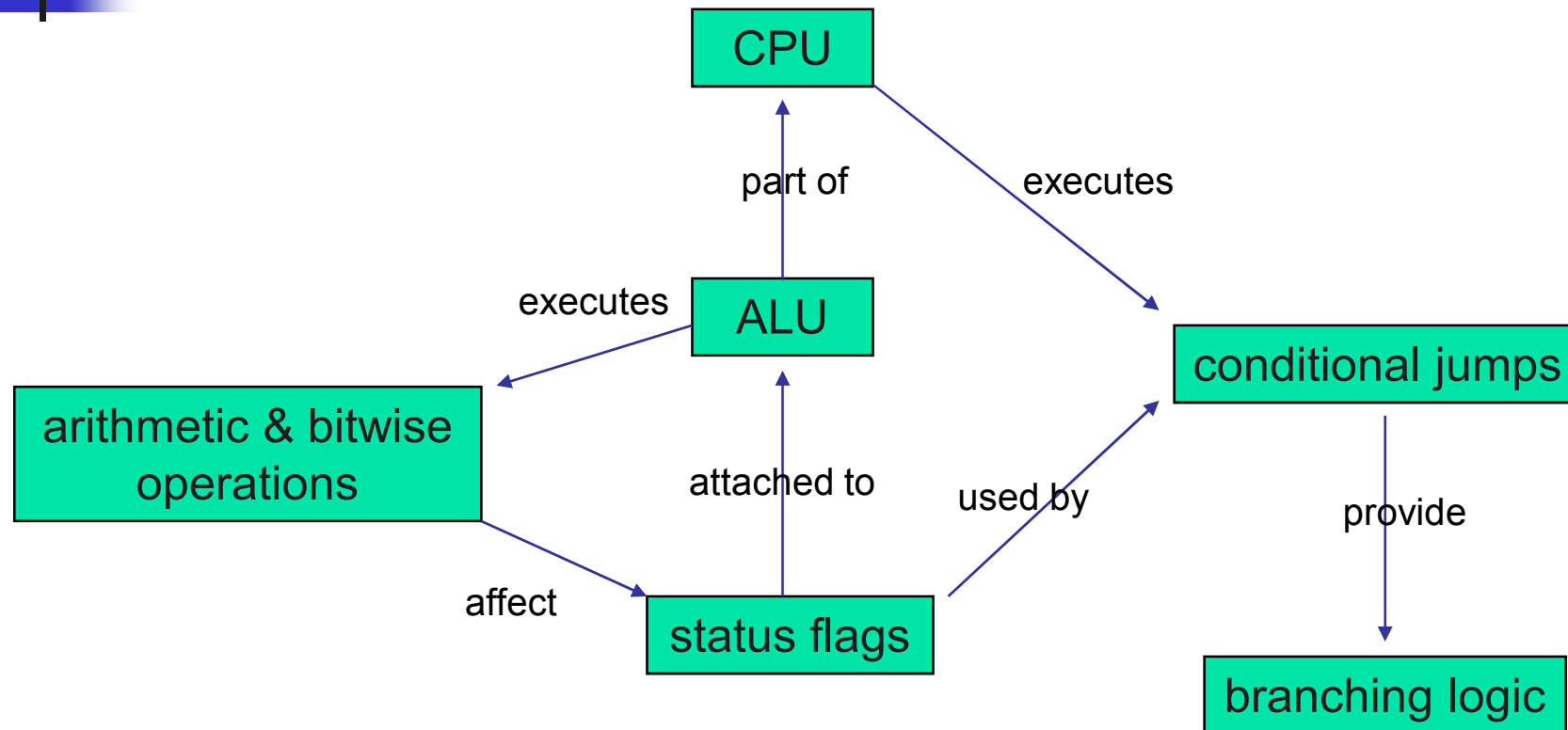
```
mov ebx, Yval
neg ebx
add ebx, Zval
mov eax, Xval
sub eax, ebx
mov Rval, eax
```



Flags Affected by Arithmetic

- The ALU has a number of status flags that reflect the outcome of arithmetic (and bitwise) operations
 - based on the contents of the destination operand
- Essential flags:
 - Zero flag – set when destination equals zero
 - Sign flag – set when destination is negative
 - Carry flag – set when unsigned value is out of range
 - Overflow flag – set when signed value is out of range
- The MOV instruction never affects the flags.

Concept Map



You can use diagrams such as these to express the relationships between assembly language concepts.



Zero Flag (ZF)

The Zero flag is set when the result of an operation produces zero in the destination operand.

```
mov cx,1
sub cx,1           ; CX = 0, ZF = 1
mov ax,0FFFFh
inc ax            ; AX = 0, ZF = 1
inc ax            ; AX = 1, ZF = 0
```

Remember...

- A flag is **set** when it equals 1.
- A flag is **clear** when it equals 0.



Sign Flag (SF)

The Sign flag is set when the destination operand is negative. The flag is clear when the destination is positive.

```
mov cx,0
sub cx,1           ; CX = -1, SF = 1
add cx,2           ; CX = 1, SF = 0
```

The sign flag is a copy of the destination's highest bit:

```
mov al,0
sub al,1           ; AL = 11111111b, SF = 1
add al,2           ; AL = 00000001b, SF = 0
```



Signed and Unsigned Integers

A Hardware Viewpoint

- All CPU instructions operate exactly the same on signed and unsigned integers
- The CPU cannot distinguish between signed and unsigned integers
- YOU, the programmer, are solely responsible for using the correct data type with each instruction

Overflow and Carry Flags

A Hardware Viewpoint

- How the **ADD** instruction modifies OF and CF:
 - $OF = (\text{carry out of the MSB}) \text{ XOR } (\text{carry into the MSB})$
 - $CF = (\text{carry out of the MSB})$
- How the **SUB** instruction modifies OF and CF:
 - NEG the source and ADD it to the destination
 - $OF = (\text{carry out of the MSB}) \text{ XOR } (\text{carry into the MSB})$
 - $CF = \text{INVERT } (\text{carry out of the MSB})$

MSB = Most Significant Bit (high-order bit)
XOR = eXclusive-OR operation
NEG = Negate (same as SUB 0,operand)



Carry Flag (CF)

The Carry flag is set when the result of an operation generates an **unsigned** value that is out of range (too big or too small for the destination operand).

```
mov al,0FFh
add al,1           ; CF = 1, AL = 00

; Try to go below zero:

mov al,0
sub al,1          ; CF = 1, AL = FF
```

In the second example, we tried to generate a negative value. Unsigned values cannot be negative, so the Carry flag signaled an error condition.



Your Turn . . .

For each of the following marked entries, show the values of the destination operand and the Sign, Zero, and Carry flags:

```
mov ax, 00FFh
add ax, 1           ; AX= 0100h  SF= 0 ZF= 0 CF= 0
sub ax, 1           ; AX= 00FFh  SF= 0 ZF= 0 CF= 0
add al, 1           ; AL= 00h    SF= 0 ZF= 1 CF= 1
mov bh, 6Ch
add bh, 95h        ; BH= 01h    SF= 0 ZF= 0 CF= 1

mov al, 2
sub al, 3           ; AL= FFh    SF= 1 ZF= 0 CF= 1
```



Overflow Flag (OF)

The Overflow flag is set when the signed result of an operation is invalid or out of range.

```
; Example 1
mov al,+127
add al,1                ; OF = 1,    AL = ??

; Example 2
mov al,7Fh
add al,1                ; OF = 1,    AL = 80h
```

The two examples are identical at the binary level because 7Fh equals +127. To determine the value of the destination operand, it is often easier to calculate in hexadecimal.



A Rule of Thumb

- When adding two integers, remember that the Overflow flag is only set when . . .
 - Two positive operands are added and their sum is negative
 - Two negative operands are added and their sum is positive

What will be the values of the Overflow flag?

```
mov al,80h
add al,92h           ; OF = 1
```

```
mov al,-2
add al,+127         ; OF = 0
```



Your Turn . . .

What will be the values of the given flags after each operation?

```
mov al, -128
neg al                ; CF = 0    OF = 1

mov ax, 8000h
add ax, 2            ; CF = 0    OF = 0

mov ax, 0
sub ax, 2            ; CF = 1    OF = 0

mov al, -5
sub al, +125         ; CF = 0    OF = 1
```



What's Next

- Data Transfer Instructions
- Addition and Subtraction
- **Data-Related Operators and Directives**
- Indirect Addressing
- JMP and LOOP Instructions

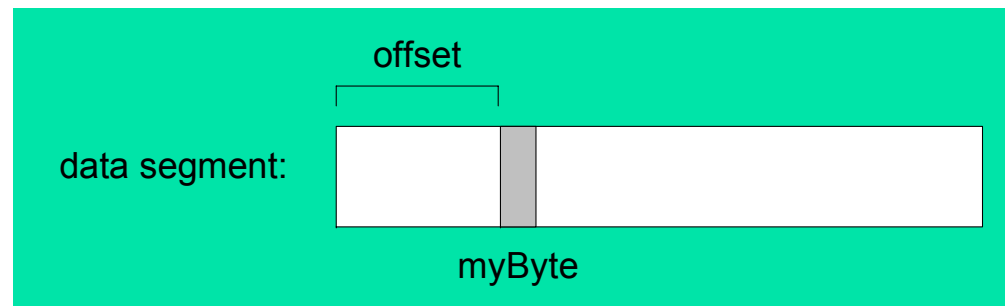


Data-Related Operators and Directives

- OFFSET Operator
- PTR Operator
- TYPE Operator
- LENGTHOF Operator
- SIZEOF Operator
- LABEL Directive

OFFSET Operator

- OFFSET returns the distance in bytes, of a label from the beginning of its enclosing segment
 - Protected mode: 32 bits
 - Real mode: 16 bits



The Protected-mode programs we write only have a single segment (we use the [flat memory model](#)).



OFFSET Examples

Let's assume that the data segment begins at 00404000h:

```
.data
bVal BYTE ?
wVal WORD ?
dVal DWORD ?
dVal2 DWORD ?

.code
mov esi,OFFSET bVal      ; ESI = 00404000
mov esi,OFFSET wVal      ; ESI = 00404001
mov esi,OFFSET dVal      ; ESI = 00404003
mov esi,OFFSET dVal2     ; ESI = 00404007
```



Relating to C/C++

The value returned by OFFSET is a pointer. Compare the following code written for both C++ and assembly language:

```
; C++ version:  
char array[1000];  
char * p = array;
```

```
.data  
array BYTE 1000 DUP(?)  
.code  
mov esi,OFFSET array           ; ESI is p
```



PTR Operator

Overrides the default type of a label (variable).
Provides the flexibility to access part of a variable.

```
.data
myDouble DWORD 12345678h
.code
mov ax,myDouble           ; error - why?

mov ax,WORD PTR myDouble  ; loads 5678h

mov WORD PTR myDouble,4321h ; saves 4321h
```

Recall that **little endian** order is used when storing data in memory (see Section 3.4.9).



Little Endian Order

- Little endian order refers to the way Intel stores integers in memory.
- Multi-byte integers are stored in reverse order, with the least significant byte stored at the lowest address
- For example, the doubleword 12345678h would be stored as:

byte	offset
78	0000
56	0001
34	0002
12	0003

When integers are loaded from memory into registers, the bytes are automatically re-reversed into their correct positions.



PTR Operator Examples

```
.data  
myDouble DWORD 12345678h
```

doubleword	word	byte	offset	
12345678	5678	78	0000	myDouble
		56	0001	myDouble + 1
	1234	34	0002	myDouble + 2
		12	0003	myDouble + 3

```
mov al,BYTE PTR myDouble ; AL = 78h  
mov al,BYTE PTR [myDouble+1] ; AL = 56h  
mov al,BYTE PTR [myDouble+2] ; AL = 34h  
mov ax,WORD PTR myDouble ; AX = 5678h  
mov ax,WORD PTR [myDouble+2] ; AX = 1234h
```



PTR Operator (cont.)

PTR can also be used to combine elements of a smaller data type and move them into a larger operand. The CPU will automatically reverse the bytes.

```
.data
myBytes BYTE 12h,34h,56h,78h

.code
mov ax,WORD PTR [myBytes]           ; AX = 3412h
mov ax,WORD PTR [myBytes+2]        ; AX = 7856h
mov eax,DWORD PTR myBytes           ; EAX = 78563412h
```



Your Turn . . .

Write down the value of each destination operand:

```
.data
varB BYTE 65h,31h,02h,05h
varW WORD 6543h,1202h
varD DWORD 12345678h

.code
mov ax,WORD PTR [varB+2]           ; a. 0502h
mov bl,BYTE PTR varD              ; b. 78h
mov bl,BYTE PTR [varW+2]         ; c. 02h
mov ax,WORD PTR [varD+2]         ; d. 1234h
mov eax,DWORD PTR varW           ; e. 12026543h
```



TYPE Operator

The TYPE operator returns the size, in bytes, of a single element of a data declaration.

```
.data
var1 BYTE ?
var2 WORD ?
var3 DWORD ?
var4 QWORD ?

.code
mov eax,TYPE var1      ; 1
mov eax,TYPE var2      ; 2
mov eax,TYPE var3      ; 4
mov eax,TYPE var4      ; 8
```



LENGTHOF Operator

The LENGTHOF operator counts the number of elements in a single data declaration.

	LENGTHOF
<code>.data</code>	
<code>byte1 BYTE 10,20,30</code>	<code>; 3</code>
<code>array1 WORD 30 DUP(?),0,0</code>	<code>; 32</code>
<code>array2 WORD 5 DUP(3 DUP(?))</code>	<code>; 15</code>
<code>array3 DWORD 1,2,3,4</code>	<code>; 4</code>
<code>digitStr BYTE "12345678",0</code>	<code>; 9</code>
<code>.code</code>	
<code>mov ecx,LENGTHOF array1</code>	<code>; 32</code>



sizeof Operator

The sizeof operator returns a value that is equivalent to multiplying LENGTHOF by TYPE.

	sizeof
<code>.data</code>	
<code>byte1 BYTE 10,20,30</code>	<code>; 3</code>
<code>array1 WORD 30 DUP(?),0,0</code>	<code>; 64</code>
<code>array2 WORD 5 DUP(3 DUP(?))</code>	<code>; 30</code>
<code>array3 DWORD 1,2,3,4</code>	<code>; 16</code>
<code>digitStr BYTE "12345678",0</code>	<code>; 9</code>
<code>.code</code>	
<code>mov ecx, sizeof array1</code>	<code>; 64</code>



Spanning Multiple Lines

A data declaration spans multiple lines if each line (except the last) ends with a comma. The LENGTHOF and SIZEOF operators include all lines belonging to the declaration:

```
.data
array WORD 10,20,
        30,40,
        50,60

.code
mov  eax,LENGTHOF array           ; 6
mov  ebx,SIZEOF array            ; 12
```



Spanning Multiple Lines (cont.)

In the following example, `array` identifies only the first `WORD` declaration. Compare the values returned by `LENGTHOF` and `SIZEOF` here to those in the previous slide:

```
.data
array  WORD 10,20
        WORD 30,40
        WORD 50,60

.code
mov  eax,LENGTHOF array           ; 2
mov  ebx,SIZEOF array            ; 4
```



LABEL Directive

- Assigns an alternate label name and type to an existing storage location
- LABEL does not allocate any storage of its own
- Removes the need for the PTR operator

```
.data
dwList    LABEL DWORD
wordList  LABEL WORD
intList   BYTE 00h,10h,00h,20h
.code
mov  eax,dwList           ; 20001000h
mov  cx,wordList         ; 1000h
mov  dl,intList           ; 00h
```



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Indirect Addressing

- Indirect Operands
- Array Sum Example
- Indexed Operands
- Pointers



Indirect Operands

An indirect operand holds the address of a variable, usually an array or string. It can be **dereferenced** (just like a pointer).

```
.data
val1 BYTE 10h,20h,30h
.code
mov esi,OFFSET val1
mov al,[esi]           ; dereference ESI (AL = 10h)

inc esi
mov al,[esi]           ; AL = 20h

inc esi
mov al,[esi]           ; AL = 30h
```



Indirect Operands (cont.)

Use PTR to clarify the size attribute of a memory operand.

```
.data
myCount WORD 0

.code
mov esi,OFFSET myCount
inc [esi] ; error: ambiguous
inc WORD PTR [esi] ; ok
```

Should PTR be used here?

```
add [esi],20
```

yes, because [esi] could point to a byte, word, or doubleword



Array Sum Example

Indirect operands are ideal for traversing an array. Note that the register in brackets must be incremented by a value that matches the array type.

```
.data
arrayW WORD 1000h,2000h,3000h
.code
    mov esi,OFFSET arrayW
    mov ax,[esi]
    add esi,2           ; or: add esi,TYPE arrayW
    add ax,[esi]
    add esi,2
    add ax,[esi]       ; AX = sum of the array
```

ToDo: Modify this example for an array of doublewords.



Indexed Operands

An indexed operand adds a constant to a register to generate an effective address. There are two notational forms: `[label + reg]` `label[reg]`

```
.data
arrayW WORD 1000h,2000h,3000h
.code
    mov esi,0
    mov ax,[arrayW + esi]           ; AX = 1000h
    mov ax,arrayW[esi]             ; alternate format
    add esi,2
    add ax,[arrayW + esi]
    etc.
```

ToDo: Modify this example for an array of doublewords.



Index Scaling

You can scale an indirect or indexed operand to the offset of an array element. This is done by multiplying the index by the array's TYPE:

```
.data
arrayB BYTE 0,1,2,3,4,5
arrayW WORD 0,1,2,3,4,5
arrayD DWORD 0,1,2,3,4,5

.code
mov esi,4
mov al,arrayB[esi*TYPE arrayB] ; 04
mov bx,arrayW[esi*TYPE arrayW] ; 0004
mov edx,arrayD[esi*TYPE arrayD] ; 00000004
```



Pointers

You can declare a **pointer variable** that contains the offset of another variable.

```
.data
arrayW WORD 1000h,2000h,3000h
ptrW DWORD arrayW
.code
    mov esi,ptrW
    mov ax,[esi]           ; AX = 1000h
```

Alternate format:

```
ptrW DWORD OFFSET arrayW
```



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- **JMP and LOOP Instructions**



JMP and LOOP Instructions

- JMP Instruction
- LOOP Instruction
- LOOP Example
- Summing an Integer Array
- Copying a String



JMP Instruction

- JMP is an unconditional jump to a label that is usually within the same procedure.
- Syntax: *JMP target*
- Logic: $EIP \leftarrow target$
- Example:

```
top:  
.  
.  
jmp top
```

A jump outside the current procedure must be to a special type of label called a [global label](#) (see Section 5.5.2 for details).



LOOP Instruction

- The LOOP instruction creates a counting loop
- Syntax: `LOOP target`
- Logic:
 - $ECX \leftarrow ECX - 1$
 - if $ECX \neq 0$, jump to *target*
- Implementation:
 - The assembler calculates the distance, in bytes, between the offset of the following instruction and the offset of the target label. It is called the **relative offset**.
 - The relative offset is added to EIP.



LOOP Example

The following loop calculates the sum of the integers
5 + 4 + 3 + 2 + 1:

offset	machine code	source code
00000000	66 B8 0000	mov ax, 0
00000004	B9 00000005	mov ecx, 5
00000009	66 03 C1	L1: add ax, cx
0000000C	E2 FB	loop L1
0000000E		

When LOOP is assembled, the current location = 0000000E (offset of the next instruction). -5 (FBh) is added to the the current location, causing a jump to location 00000009:

$$00000009 \leftarrow 0000000E + FB$$



Your Turn . . .

If the relative offset is encoded in a single signed byte,

(a) what is the largest possible backward jump?

(b) what is the largest possible forward jump?

(a) -128

(b) +127



Your Turn . . .

What will be the final value of AX?

10

```
mov ax, 6
mov ecx, 4
L1:
inc ax
loop L1
```

How many times will the loop execute?

4,294,967,296

```
mov ecx, 0
x2:
inc ax
loop x2
```



Nested Loop

If you need to code a loop within a loop, you must save the outer loop counter's ECX value.

```
.data
count DWORD ?
.code
    mov ecx,100          ; set outer loop count
L1:
    mov count,ecx       ; save outer loop count
    mov ecx,20          ; set inner loop count
L2: .
    .
    loop L2             ; repeat the inner loop
    mov ecx,count       ; restore outer loop count
    loop L1             ; repeat the outer loop
```



Summing an Integer Array

The following code calculates the sum of an array of 16-bit integers.

```
.data
intarray WORD 100h,200h,300h,400h
.code
    mov edi,OFFSET intarray      ; address of intarray
    mov ecx,LENGTHOF intarray   ; loop counter
    mov ax,0                     ; zero the accumulator
L1:
    add ax,[edi]                 ; add an integer
    add edi,TYPE intarray        ; point to next integer
    loop L1                      ; repeat until ECX = 0
```



Your Turn . . .

What changes would you make to the program on the previous slide if you were summing a doubleword array?



Copying a String

The following code copies a string from **source** to **target**:

```
.data
source  BYTE  "This is the source string",0
target  BYTE  sizeof source DUP(0)

.code
    mov  esi,0                ; index register
    mov  ecx, sizeof source  ; loop counter
L1:
    mov  al, source[esi]      ; get char from source
    mov  target[esi], al      ; store it in the target
    inc  esi                  ; move to next character
    loop L1                   ; repeat for entire string
```

good use of
sizeof



Your Turn . . .

Rewrite the program shown in the previous slide, using indirect addressing rather than indexed addressing.





CMP Instruction (See 6.2.7)

- Compares the destination operand to the source operand (both are unsigned)
- Syntax: *CMP destination, source*
- Example: destination == source

```
mov al,5  
cmp al,5                ; Zero flag set
```

- Example: destination < source

```
mov al,4  
cmp al,5                ; Carry flag set
```



CMP Instruction

- Example: destination > source

```
mov al,6  
cmp al,5 ; ZF = 0, CF = 0
```

(both the Zero and Carry flags are clear)



Jcond Instruction

- A conditional jump instruction branches to a label when specific register or flag conditions are met
- Examples:
 - JB, JC jump to a label if the Carry flag is set
 - JE, JZ jump to a label if the Zero flag is set
 - JS jumps to a label if the Sign flag is set
 - JNE, JNZ jump to a label if the Zero flag is clear
 - JECXZ jumps to a label if ECX equals 0



Jcond Ranges

- Prior to the 386:
 - jump must be within -128 to $+127$ bytes from current location counter
- IA-32 processors:
 - 32-bit offset permits jump anywhere in memory



Jumps Based on Specific Flags

Mnemonic	Description	Flags
JZ	Jump if zero	ZF = 1
JNZ	Jump if not zero	ZF = 0
JC	Jump if carry	CF = 1
JNC	Jump if not carry	CF = 0
JO	Jump if overflow	OF = 1
JNO	Jump if not overflow	OF = 0
JS	Jump if signed	SF = 1
JNS	Jump if not signed	SF = 0
JP	Jump if parity (even)	PF = 1
JNP	Jump if not parity (odd)	PF = 0



Jumps Based on Equality

Mnemonic	Description
JE	Jump if equal (<i>leftOp = rightOp</i>)
JNE	Jump if not equal (<i>leftOp ≠ rightOp</i>)
JCXZ	Jump if CX = 0
JECXZ	Jump if ECX = 0



Jumps Based on Unsigned Comparisons

Mnemonic	Description
JA	Jump if above (if $leftOp > rightOp$)
JNBE	Jump if not below or equal (same as JA)
JAE	Jump if above or equal (if $leftOp \geq rightOp$)
JNB	Jump if not below (same as JAE)
JB	Jump if below (if $leftOp < rightOp$)
JNAE	Jump if not above or equal (same as JB)
JBE	Jump if below or equal (if $leftOp \leq rightOp$)
JNA	Jump if not above (same as JBE)



Jumps Based on Signed Comparisons

Mnemonic	Description
JG	Jump if greater (if $leftOp > rightOp$)
JNLE	Jump if not less than or equal (same as JG)
JGE	Jump if greater than or equal (if $leftOp \geq rightOp$)
JNL	Jump if not less (same as JGE)
JL	Jump if less (if $leftOp < rightOp$)
JNGE	Jump if not greater than or equal (same as JL)
JLE	Jump if less than or equal (if $leftOp \leq rightOp$)
JNG	Jump if not greater (same as JLE)



Block-Structured IF Statements

Assembly language programmers can easily translate logical statements written in C++/Java into assembly language. For example:

```
if( op1 == op2 )
    X = 1;
else
    X = 2;
```

```
mov  eax,op1
cmp  eax,op2
jne  L1
mov  X,1
jmp  L2
L1:  mov  X,2
L2:
```



Your Turn . . .

Implement the following pseudocode in assembly language. All values are unsigned:

```
if( ebx <= ecx )
{
    eax = 5;
    edx = 6;
}
```

```
cmp ebx,ecx
ja next
mov eax,5
mov edx,6
next:
```

(There are multiple correct solutions to this problem.)



WHILE Loops

A WHILE loop is really an IF statement followed by the body of the loop, followed by an unconditional jump to the top of the loop. Consider the following example:

```
while( eax < ebx)
    eax = eax + 1;
```

This is a possible implementation:

```
top: cmp  eax, ebx           ; check loop condition
     jae  next              ; false? exit loop
     inc  eax               ; body of loop
     jmp  top               ; repeat the loop
next:
```



Your Turn . . .

Implement the following loop, using unsigned 32-bit integers:

```
while( ebx <= val1)
{
    ebx = ebx + 5;
    val1 = val1 - 1
}
```

```
top: cmp ebx, val1          ; check loop condition
     ja  next              ; false? exit loop
     add ebx, 5            ; body of loop
     dec val1
     jmp top               ; repeat the loop
next:
```



DIV Instruction (See 7.4.4)

- The DIV (unsigned divide) instruction performs 8-bit, 16-bit, and 32-bit division on unsigned integers
- A single operand is supplied (register or memory operand), which is assumed to be the divisor
- Instruction formats:

DIV *r/m8*

DIV *r/m16*

DIV *r/m32*

Default Operands:

Dividend	Divisor	Quotient	Remainder
AX	<i>r/m8</i>	AL	AH
DX:AX	<i>r/m16</i>	AX	DX
EDX:EAX	<i>r/m32</i>	EAX	EDX



DIV Examples

Divide 8003h by 100h, using 16-bit operands:

```
mov dx,0           ; clear dividend, high
mov ax,8003h       ; dividend, low
mov cx,100h        ; divisor
div cx             ; AX = 0080h, DX = 3
```

Same division, using 32-bit operands:

```
mov edx,0          ; clear dividend, high
mov eax,8003h      ; dividend, low
mov ecx,100h       ; divisor
div ecx            ; EAX = 00000080h, DX = 3
```



Your Turn . . .

What will be the hexadecimal values of DX and AX after the following instructions execute?

```
mov dx,0087h  
mov ax,6000h  
mov bx,100h  
div bx
```

DX = 0000h, AX = 8760h

