



# ***Computer Organization & Assembly Languages***

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## ***Advanced Procedure***

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Adapted from the slides prepared by Kip Irvine for the book,  
Assembly Language for Intel-Based Computers, 5th Ed.



# Chapter Overview

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- **Stack Frames**
- Recursion
- .MODEL Directive
- INVOKE, ADDR, PROC, and PROTO
- Creating Multimodule Programs



# Stack Frames

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- Stack Parameters
- Local Variables
- ENTER and LEAVE Instructions
- LOCAL Directive



# Stack Parameters

---

- More convenient than register parameters
- Two possible ways of calling DumpMem.

Which is easier?

```
pushad
mov esi,OFFSET array
mov ecx,LENGTHOF array
mov ebx,TYPE array
call DumpMem
popad
```

Register-based Method

```
push TYPE array
push LENGTHOF array
push OFFSET array
call DumpMem
```

Stack-based Method



# Stack Frame

---

- Also known as an *activation record*
- Area of the stack set aside for a procedure's return address, passed parameters, saved registers, and local variables
- Created by the following steps:
  - Calling program pushes arguments on the stack and calls the procedure.
  - The called procedure pushes EBP on the stack, and sets EBP to ESP.
  - If local variables are needed, a constant is subtracted from ESP to make room on the stack.



# Explicit Access to Stack Parameters

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- A procedure can explicitly access stack parameters using constant offsets from EBP.
  - Example: `[ebp + 8]`
- EBP is often called the **base pointer** or **frame pointer** because it holds the base address of the stack frame.
- EBP does not change value during the procedure.
- EBP must be restored to its original value when a procedure returns.



# RET Instruction

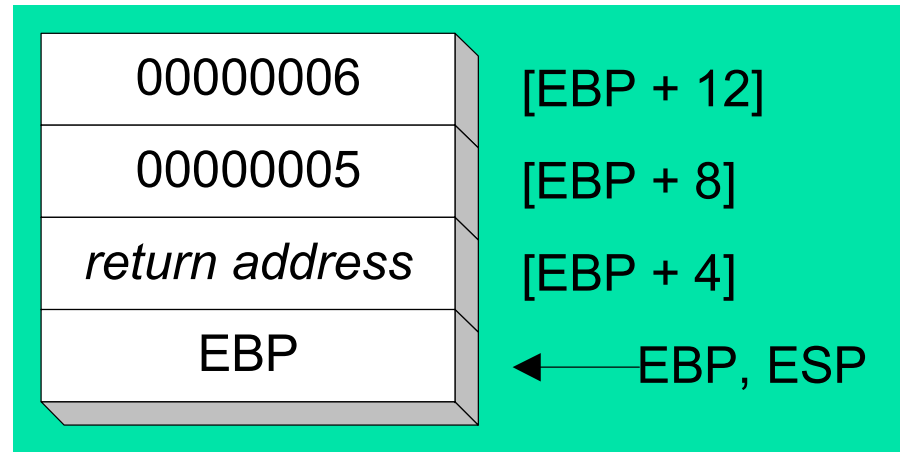
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- *Return from subroutine*
- Pops stack into the instruction pointer (EIP or IP). Control transfers to the target address.
- Syntax:
  - **RET**
  - **RET *n***
- Optional operand *n* causes *n* bytes to be added to the stack pointer after EIP (or IP) is assigned a value.

# Stack Frame Example

```
.data
sum DWORD ?
.code
    push 6           ; second argument
    push 5           ; first argument
    call AddTwo      ; EAX = sum
    mov  sum, eax    ; save the sum
```

```
AddTwo PROC
    push ebp
    mov  ebp, esp
    .
    .
```





# Passing Arguments by Reference

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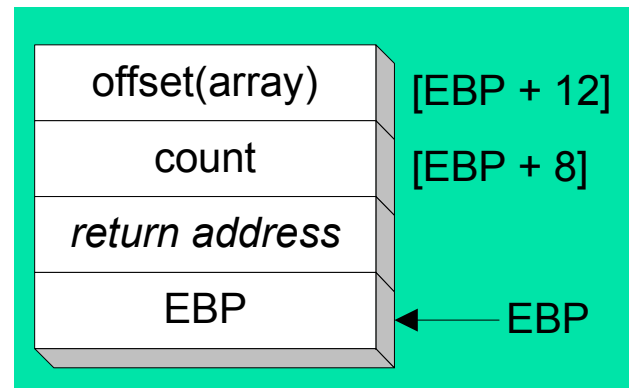
- The `ArrayFill` procedure fills an array with 16-bit random integers
- The calling program passes the address of the array, along with a count of the number of array elements:

```
.data
count = 100
array WORD count DUP(?)
.code
    push OFFSET array
    push COUNT
    call ArrayFill
```

# Passing Arguments by Reference (cont.)

ArrayFill can reference an array without knowing the array's name:

```
ArrayFill PROC
    push ebp
    mov  ebp, esp
    pushad
    mov  esi, [ebp+12]
    mov  ecx, [ebp+8]
    .
    .
```



ESI points to the beginning of the array, so it's easy to use a loop to access each array element.



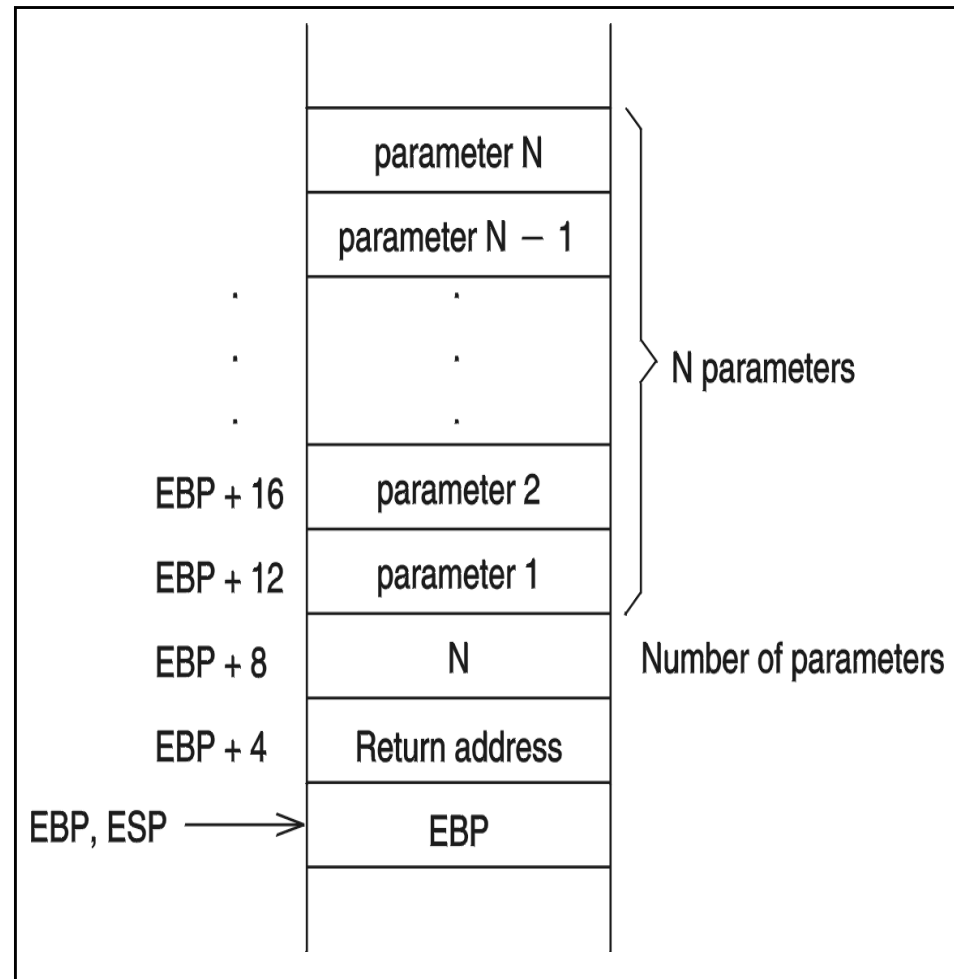
# Variable Number of Parameters

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- For most procedures, the number of parameters is fixed
  - Every time the procedure is called, the same number of parameter values are passed
- In procedures that can have variable number of parameters
  - With each procedure call, the number of parameter values passed can be different
    - C supports procedures with variable number of parameters such as *printf*
  - Easy to support variable number of parameters using the stack method

# Variable Number of Parameters (cont.)

- To implement variable number of parameter passing:
  - Parameter count should be one of the parameters passed
  - This count should be the last parameter pushed onto the stack





# Local Variables

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- To explicitly create local variables, subtract their total size from ESP.
- The following example creates and initializes two 32-bit local variables (we'll call them `locA` and `locB`):

```
MySub PROC
    push ebp
    mov  ebp, esp
    sub  esp, 8
    mov  [ebp-4], 123456h        ; locA
    mov  [ebp-8], 0             ; locB
    .
    .
```



## Local Variables (cont.)

---

- To clear local variables, set ESP to be EBP

```
MySub PROC
    push ebp
    mov  ebp, esp
    sub  esp, 8
    mov  [ebp-4], 123456h        ; locA
    mov  [ebp-8], 0             ; locB
    .
    .

    mov  esp, ebp
    pop  ebp
    ret
```



# LEA Instruction

---

- The LEA instruction returns offsets of both direct and indirect operands.
  - OFFSET operator can only return constant offsets.
- LEA is required when obtaining the offset of a stack parameter or local variable. For example:

```
CopyString PROC,  
    count:DWORD  
    LOCAL temp[20]:BYTE  
  
    mov edi,OFFSET count        ; invalid operand  
    mov esi,OFFSET temp        ; invalid operand  
    lea edi,count              ; ok  
    lea esi,temp               ; ok
```



# ENTER and LEAVE

---

- ENTER instruction creates stack frame for a called procedure
  - pushes EBP on the stack
  - sets EBP to the base of the stack frame
  - reserves space for local variables
  - Example:
    - MySub PROC
    - enter 8,0
  - Equivalent to:
    - MySub PROC
    - push ebp
    - mov ebp,esp
    - sub esp,8



# LEAVE

---

MySub PROC

```
push ebp
mov  ebp, esp
sub  esp, 8
```

```
mov  eax, val1
add  eax, val2
```

```
leave
ret  8
```

AddTwo ENDP

The LEAVE instruction  
is shorthand for:

```
mov  esp, ebp
pop  ebp
```



# LOCAL Directive

---

- A **local variable** is created, used, and destroyed within a single procedure
- The LOCAL directive declares a list of local variables
  - immediately follows the PROC directive
  - each variable is assigned a type
- Syntax:
  - LOCAL *varlist*

Example:

```
MySub PROC
    LOCAL var1:BYTE, var2:WORD, var3:SDWORD
```



# Using LOCAL

---

## Examples:

```
LOCAL flagVals[20]:BYTE    ; array of bytes
```

```
LOCAL pArray:PTR WORD     ; pointer to an array
```

```
myProc PROC,              ; procedure
    LOCAL t1:BYTE,        ; local variables
           t2:WORD,
           t3:DWORD,
           t4:PTR DWORD
```



# LOCAL Example

---

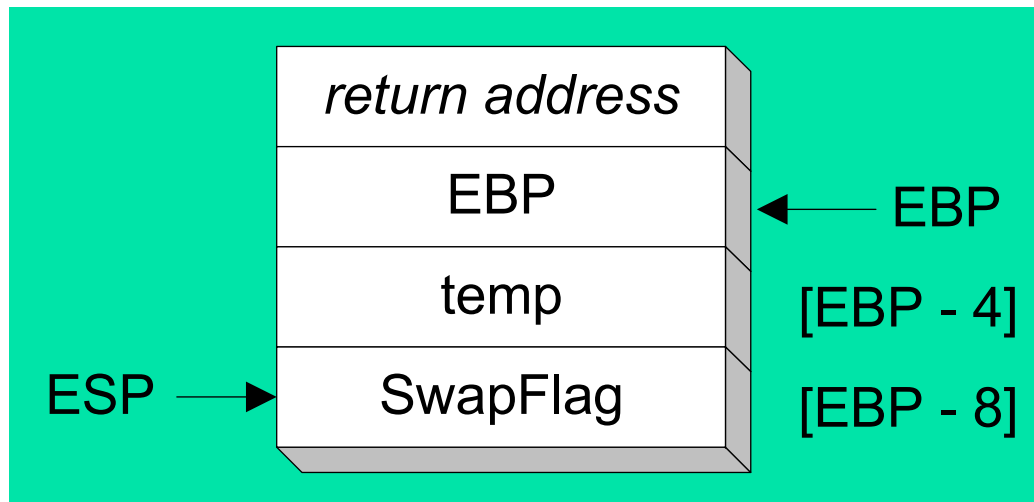
```
BubbleSort PROC
    LOCAL temp:DWORD, SwapFlag:BYTE
    . . .
    ret
BubbleSort ENDP
```

MASM generates the following code:

```
BubbleSort PROC
    push ebp
    mov  ebp,esp
    add  esp,0FFFFFFF8h          ; add -8 to ESP
    . . .
    mov  esp,ebp
    pop  ebp
    ret
BubbleSort ENDP
```

## LOCAL Example (cont.)

Diagram of the stack frame for the BubbleSort procedure:





# Non-Doubleword Local Variables

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- Local variables can be different sizes
- How created in the stack by LOCAL directive:
  - 8-bit: assigned to next available byte
  - 16-bit: assigned to next even (word) boundary
  - 32-bit: assigned to next doubleword boundary





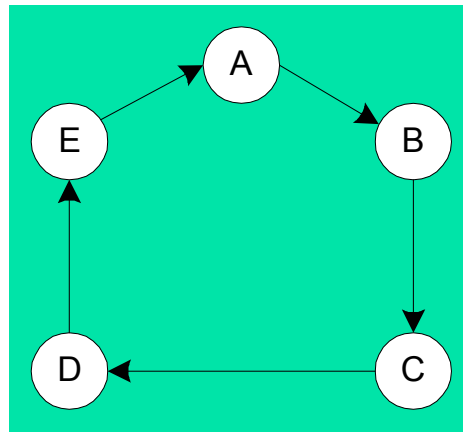
# Recursion

---

- What is recursion?
- Recursively Calculating a Sum
- Calculating a Factorial

# What is Recursion?

- The process created when . . .
  - A procedure calls itself
  - Procedure A calls procedure B, which in turn calls procedure A
- Using a graph in which each node is a procedure and each edge is a procedure call, recursion forms a **cycle**:





# Recursively Calculating a Sum

The CalcSum procedure recursively calculates the sum of an array of integers. Receives: ECX = count. Returns: EAX = sum

```
CalcSum PROC
    cmp ecx,0                ; check counter value
    jz L2                    ; quit if zero
    add eax,ecx              ; otherwise, add to sum
    dec ecx                  ; decrement counter
    call CalcSum             ; recursive call
L2: ret
CalcSum ENDP
```

Stack frame:

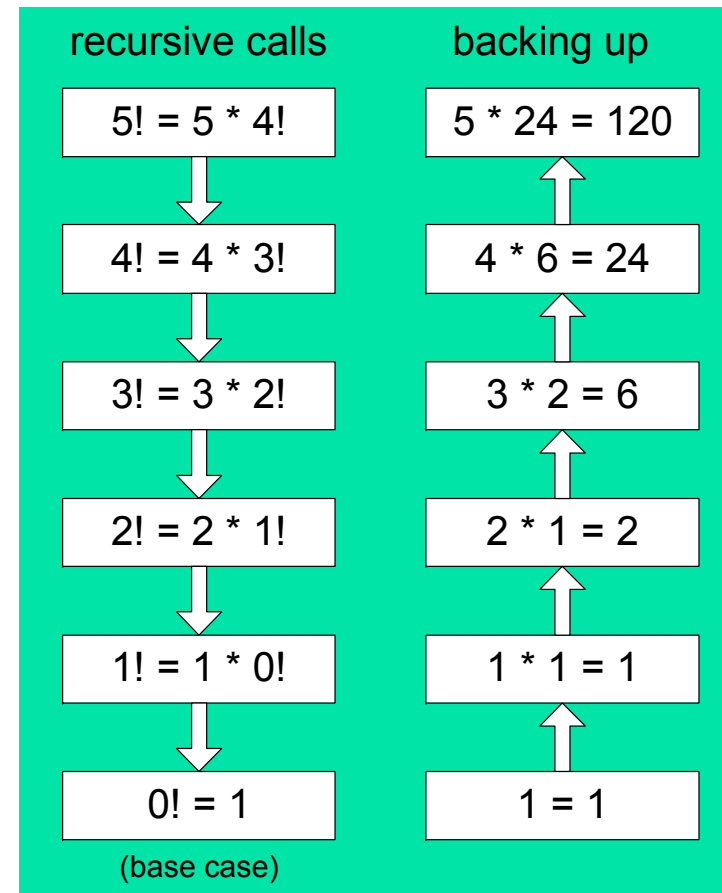
Pushed On Stack	ECX	EAX
L1	5	0
L2	4	5
L2	3	9
L2	2	12
L2	1	14
L2	0	15

# Calculating a Factorial

This function calculates the factorial of integer  $n$ . A new value of  $n$  is saved in each stack frame:

```
int function factorial(int n)
{
    if(n == 0)
        return 1;
    else
        return n * factorial(n-1);
}
```

As each call instance returns, the product it returns is multiplied by the previous value of  $n$ .





# Calculating a Factorial (cont.)

---

Factorial PROC

```
    push ebp
    mov  ebp,esp
    mov  eax,[ebp+8]           ; get n
    cmp  eax,0               ; n < 0?
    ja   L1                  ; yes: continue
    mov  eax,1               ; no: return 1
    jmp  L2
```

```
L1: dec  eax
    push eax                 ; Factorial(n-1)
    call Factorial
```

```
; Instructions from this point on execute when each
; recursive call returns.
```

ReturnFact:

```
    mov  ebx,[ebp+8]         ; get n
    mul  ebx                 ; eax = eax * ebx
```

```
L2: pop  ebp                ; return EAX
    ret  4                   ; clean up stack
```

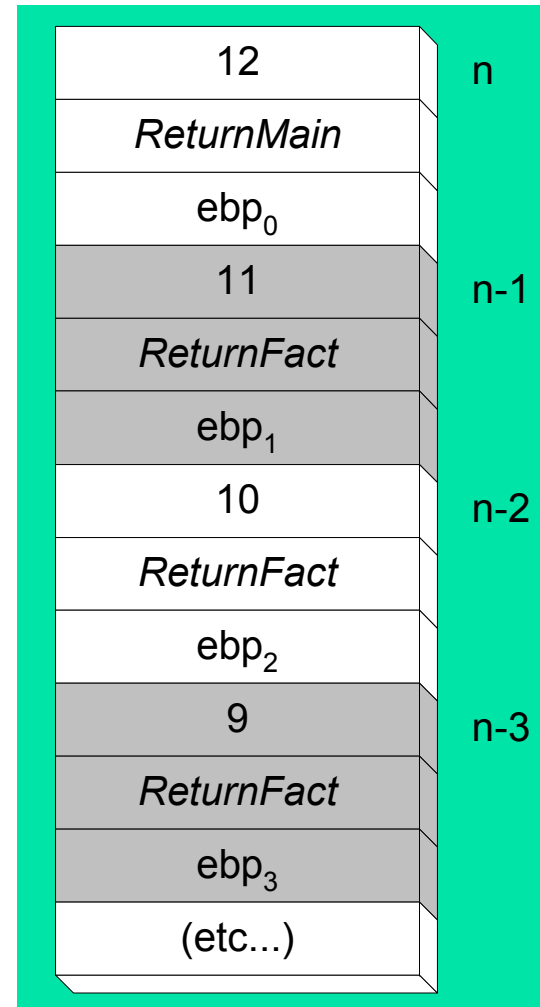
Factorial ENDP

# Calculating a Factorial (cont.)

Suppose we want to calculate 12!

This diagram shows the first few stack frames created by recursive calls to Factorial

Each recursive call uses 12 bytes of stack space.





# Reserving Stack Space

---

- `.stack 4096`
- `Sub1` calls `Sub2`, `Sub2` calls `Sub3`

`Sub1 PROC`

```
LOCAL array1[50]:DWORD ; 200 bytes
```

`Sub2 PROC`

```
LOCAL array2[80]:WORD ; 160 bytes
```

`Sub3 PROC`

```
LOCAL array3[300]:WORD ; 300 bytes
```



# What's Next

---

- Stack Frames
- Recursion
- **.MODEL Directive**
- INVOKE, ADDR, PROC, and PROTO
- Creating Multimodule Programs



## .MODEL Directive

---

- .MODEL directive specifies a program's memory model and model options (language-specifier).
- Syntax:
  - **.MODEL *memorymodel* [,*modeloptions*]**
- *memorymodel* can be one of the following:
  - tiny, small, medium, compact, large, huge, or flat
- *modeloptions* includes the language specifier:
  - procedure naming scheme
  - parameter passing conventions



# Memory Models

---

- A program's memory model determines the number and sizes of code and data segments.
- Real-address mode supports **tiny**, **small**, **medium**, **compact**, **large**, and **huge** models.
- Protected mode supports only the **flat** model.

Small model: code < 64 KB, data (including stack) < 64 KB. All offsets are 16 bits.

Flat model: single segment for code and data, up to 4 GB. All offsets are 32 bits.



# Language Specifiers

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- C

- procedure arguments pushed on stack in reverse order (right to left)
- calling program cleans up the stack

- PASCAL

- procedure arguments pushed in forward order (left to right)
- called procedure cleans up the stack

- STDCALL

- procedure arguments pushed on stack in reverse order (right to left)
- called procedure cleans up the stack



# What's Next

---

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# INVOKE, ADDR, PROC, and PROTO

---

- INVOKE Directive
- ADDR Operator
- PROC Directive
- PROTO Directive
- Parameter Classifications
- Debugging Tips



# INVOKE Directive

---

- The INVOKE directive is a powerful replacement for Intel's CALL instruction that lets you pass multiple arguments
- Syntax:
  - INVOKE *procedureName* [, *argumentList*]
- *ArgumentList* is an optional comma-delimited list of procedure arguments
- Arguments can be:
  - immediate values and integer expressions
  - variable names
  - address and ADDR expressions
  - register names



# INVOKE Examples

---

```
.data
byteVal BYTE 10
wordVal WORD 1000h
.code
    ; direct operands:
    INVOKE Sub1,byteVal,wordVal

    ; address of variable:
    INVOKE Sub2,ADDR byteVal

    ; register name, integer expression:
    INVOKE Sub3,eax,(10 * 20)

    ; address expression (indirect operand):
    INVOKE Sub4,[ebx]
```



# INVOKE Example

---

```
.data
```

```
val1 DWORD 12345h
```

```
val2 DWORD 23456h
```

```
.code
```

```
    INVOKE AddTwo, val1, val2
```

```
push val1
```

```
push val2
```

```
call AddTwo
```



# ADDR Operator

---

- Returns a near or far pointer to a variable, depending on which memory model your program uses:
  - Small model: returns 16-bit offset
  - Large model: returns 32-bit segment/offset
  - Flat model: returns 32-bit offset
- Simple example:

```
.data  
myWord WORD ?  
.code  
INVOKE mySub, ADDR myWord
```



## Your Turn . . .

---

- Create a procedure named **Difference** that subtracts the first argument from the second one. Following is a sample call:

- `push 14` ; first argument
- `push 30` ; second argument
- `call Difference` ; EAX = 16

```
Difference PROC
    push ebp
    mov  ebp, esp
    mov  eax, [ebp + 8]      ; second argument
    sub  eax, [ebp + 12]    ; first argument
    pop  ebp
    ret  8
Difference ENDP
```



# Passing by Value

---

- When a procedure argument is passed by value, a copy of a 16-bit or 32-bit integer is pushed on the stack.

Example:

```
.data
myData WORD 1000h
.code
main PROC
    INVOKE Sub1, myData
```

MASM generates the following code:

```
push myData
call Sub1
```



# Passing by Reference

---

- When an argument is passed by reference, its address is pushed on the stack. Example:

```
.data  
myData WORD 1000h  
.code  
main PROC  
    INVOKE Sub1, ADDR myData
```

MASM generates the following code:

```
push OFFSET myData  
call Sub1
```



# PROC Directive

---

- The PROC directive declares a procedure with an optional list of named parameters.
- Syntax:  
*label PROC paramList*
- *paramList* is a list of parameters separated by commas. Each parameter has the following syntax:  
*paramName : type*

*type* must either be one of the standard ASM types (BYTE, SBYTE, WORD, etc.), or it can be a pointer to one of these types.



## PROC Directive (cont.)

---

- Alternate format permits parameter list to be on one or more separate lines:

*label* PROC, ← comma required  
paramList

- The parameters can be on the same line . . .

*param-1:type-1, param-2:type-2, . . . , param-n:type-n*

- Or they can be on separate lines:

*param-1:type-1,*

*param-2:type-2,*

*. . . ,*

*param-n:type-n*



# PROC Examples

---

FillArray receives a pointer to an array of bytes, a single byte fill value that will be copied to each element of the array, and the size of the array.

```
FillArray PROC,  
    pArray:PTR BYTE, fillVal:BYTE  
    arraySize:DWORD  
  
    mov ecx,arraySize  
    mov esi,pArray  
    mov al,fillVal  
L1: mov [esi],al  
    inc esi  
    loop L1  
    ret  
FillArray ENDP
```



# PROC Examples (cont.)

---

```
Swap PROC,  
    pValX:PTR DWORD,  
    pValY:PTR DWORD  
    . . .  
Swap ENDP
```

```
ReadFile PROC,  
    pBuffer:PTR BYTE  
    LOCAL fileHandle:DWORD  
    . . .  
ReadFile ENDP
```



# PROTO Directive

---

- Creates a procedure prototype
- Syntax:
  - *label PROTO paramList*
- Every procedure called by the INVOKE directive must have a prototype
- A complete procedure definition can also serve as its own prototype



# PROTO Directive

---

- Standard configuration: PROTO appears at top of the program listing, INVOKE appears in the code segment, and the procedure implementation occurs later in the program:

```
MySub PROTO                ; procedure prototype

.code
INVOKE MySub               ; procedure call

MySub PROC                 ; procedure implementation
    .
    .
MySub ENDP
```



# PROTO Example

---

- Prototype for the ArraySum procedure, showing its parameter list:

```
ArraySum PROTO,  
    ptrArray:PTR DWORD,      ; points to the array  
    szArray:DWORD            ; array size
```



# WriteStackFrame Procedure

---

- Displays contents of current stack frame

- Prototype:

WriteStackFrame PROTO,

numParam:DWORD, ; number of passed parameters

numLocalVal: DWORD, ; number of DWordLocal variables

numSavedReg: DWORD ; number of saved registers



# WriteStackFrame Example

---

- main PROC
  - mov eax, 0EAEAEAEA h
  - mov ebx, 0EBEBEBEB h
  - INVOKE aProc, 1111h, 2222h
  - exit
- main ENDP
  
- aProc PROC USES eax ebx,
  - x: DWORD, y: DWORD
  - LOCAL a:DWORD, b:DWORD
  - PARAMS = 2
  - LOCALS = 2
  - SAVED\_REGS = 2
  - mov a,0AAAA h
  - mov b,0BBBB h
  - INVOKE WriteStackFrame, PARAMS, LOCALS, SAVED\_REGS



# Parameter Classifications

---

- An **input parameter** is data passed by a calling program to a procedure.
  - The called procedure is not expected to modify the corresponding parameter variable, and even if it does, the modification is confined to the procedure itself.
- An **output parameter** is created by passing a pointer to a variable when a procedure is called.
  - The procedure does not use any existing data from the variable, but it fills in a new value before it returns.
- An **input-output parameter** is a pointer to a variable containing input that will be both used and modified by the procedure.
  - The variable passed by the calling program is modified.



# Example: Exchanging Two Integers

---

The Swap procedure exchanges the values of two 32-bit integers. `pValX` and `pValY` do not change values, but the integers they point to are modified.

```
Swap PROC USES eax esi edi,  
    pValX:PTR DWORD,           ; pointer to first integer  
    pValY:PTR DWORD           ; pointer to second integer  
  
    mov esi,pValX              ; get pointers  
    mov edi,pValY  
    mov eax,[esi]              ; get first integer  
    xchg eax,[edi]             ; exchange with second  
    mov [esi],eax              ; replace first integer  
    ret  
Swap ENDP
```



# Trouble-Shooting Tips

---

- Save and restore registers when they are modified by a procedure.
  - Except a register that returns a function result
- When using INVOKE, be careful to pass a pointer to the correct data type.
  - For example, MASM cannot distinguish between a DWORD argument and a PTR BYTE argument.
- Do not pass an immediate value to a procedure that expects a reference parameter.
  - Dereferencing its address will likely cause a general-protection fault.



# What's Next

---

- Stack Frames
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# Multimodule Programs

---

- A **multimodule program** is a program whose source code has been divided up into separate ASM files.
- Each ASM file (module) is assembled into a separate OBJ file.
- All OBJ files belonging to the same program are linked using the **link** utility into a single EXE file.
  - This process is called **static linking**



# Advantages

---

- Large programs are easier to write, maintain, and debug when divided into separate source code modules.
- When changing a line of code, only its enclosing module needs to be assembled again. Linking assembled modules requires little time.
- A module can be a container for logically related code and data (think object-oriented here...)
  - **encapsulation**: procedures and variables are automatically hidden in a module unless you declare them public



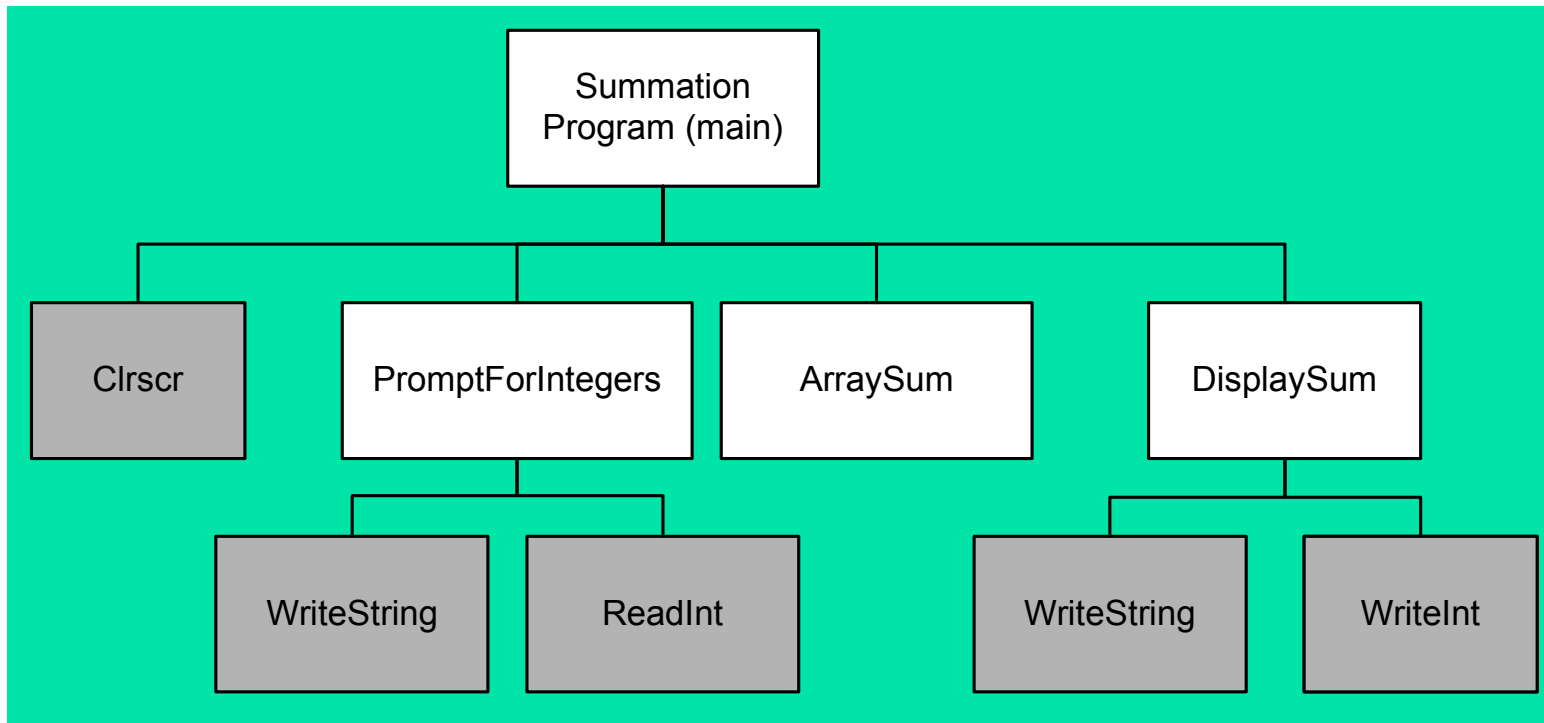
# Creating a Multimodule Program

---

- Here are some basic steps to follow when creating a multimodule program:
  - Create the main module
  - Create a separate source code module for each procedure or set of related procedures
  - Create an include file that contains procedure prototypes for **external procedures** (ones that are called between modules)
  - Use the INCLUDE directive to make your procedure prototypes available to each module

# Example: ArraySum Program

- Let's review the ArraySum program from Ch5.



Each of the four white rectangles will become a module.



# Sample Program output

---

```
Enter a signed integer: -25
```

```
Enter a signed integer: 36
```

```
Enter a signed integer: 42
```

```
The sum of the integers is: +53
```



# INCLUDE File

---

The `sum.inc` file contains prototypes for external functions that are not in the Irvine32 library:

```
INCLUDE Irvine32.inc
```

```
PromptForIntegers PROTO,
```

```
    ptrPrompt:PTR BYTE,           ; prompt string  
    ptrArray:PTR DWORD,          ; points to the array  
    arraySize:DWORD              ; size of the array
```

```
ArraySum PROTO,
```

```
    ptrArray:PTR DWORD,          ; points to the array  
    count:DWORD                  ; size of the array
```

```
DisplaySum PROTO,
```

```
    ptrPrompt:PTR BYTE,          ; prompt string  
    theSum:DWORD                 ; sum of the array
```



# Main.asm

```
TITLE Integer Summation Program

INCLUDE sum.inc

.code
main PROC
    call Clrscr

    INVOKE PromptForIntegers,
        ADDR prompt1,
        ADDR array,
        Count

    ...
    call Crlf
    INVOKE ExitProcess,0
main ENDP
END main
```