

Ex. $L = \{w \in \{0,1\}^* : w \text{ has two or three occurrences of } 1, \text{ the first and second of which are not consecutive.}\}$

$$L = \{0\}^* \circ \{1\} \circ \{0\} \circ \{0\}^* \circ \{1\} \circ \{0\}^* \circ (\{1\} \circ \{0\}^* \cup \phi^*)$$

↑
{e}

$$L = 0^* 1 0 0^* 1 0^* (1 0^* \cup \phi^*)$$

Regular expressions:

(1) ϕ and $a \in \Sigma$: regular expression.

(2) α, β : regular exp. $\Rightarrow (\alpha\beta)$: regular exp.

(3) α, β : regular exp. $\Rightarrow (\alpha \cup \beta)$: regular exp.

(4) α : regular exp. $\Rightarrow \alpha^*$: regular exp.

$L(\alpha)$: the language represented by α .

The complement of α is regular. (DFA ^{Later} accepting states \downarrow other states)

$$\alpha \cap \beta = \overline{\overline{\alpha} \cup \overline{\beta}}$$

Ex. $L((a \cup b)^* a) = \{a, b\}^* \{a\}$

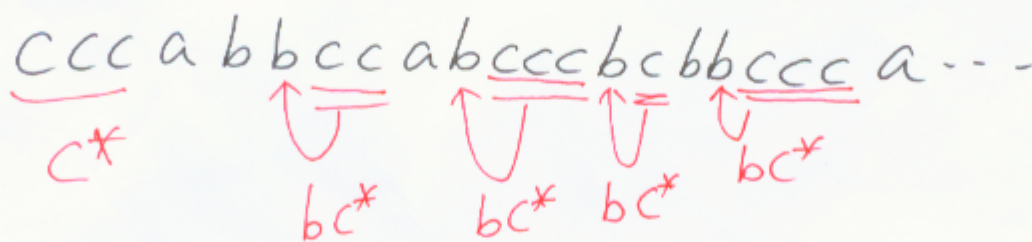
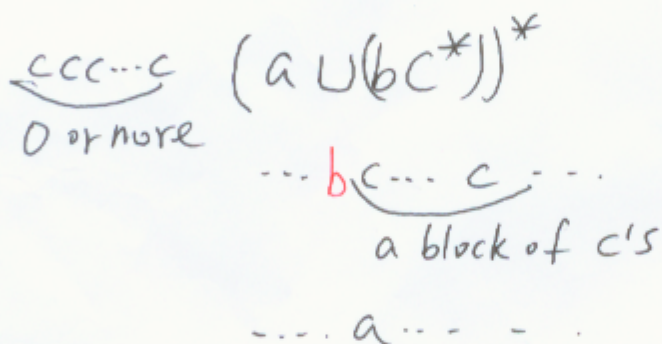
Kim-Ma. Chew
10/9, 2012.

$= \{w \in \{a, b\}^* : w \text{ ends with an } a\}$

Ex. $L(c^*(a \cup (bc^*))^*) = ?$

1. No string in $L(c^*(a \cup (bc^*))^*)$ contains the substring "ac".

2. Any string that does not contain a c



$L(c^*(a \cup (bc^*))^*) = \{w \in \{a, b, c\}^* : w \text{ does not contain the substring } ac\}$.

Cf. $(a^* b \cup c)^* a^*$

$$\text{Ex. } \alpha = (0U1)^* 111 (0U1)^*$$

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Oct. 9, 2012

What is $\mathcal{L}(\alpha)$?

$$\mathcal{L}(\alpha) = \{ w \in \{0,1\}^* : w \text{ has the substring } 111 \}$$

$$\text{Ex. } \beta = (0^* U ((0^* (1U(111))) ((00^*) (1U(111)))^* U^*))$$

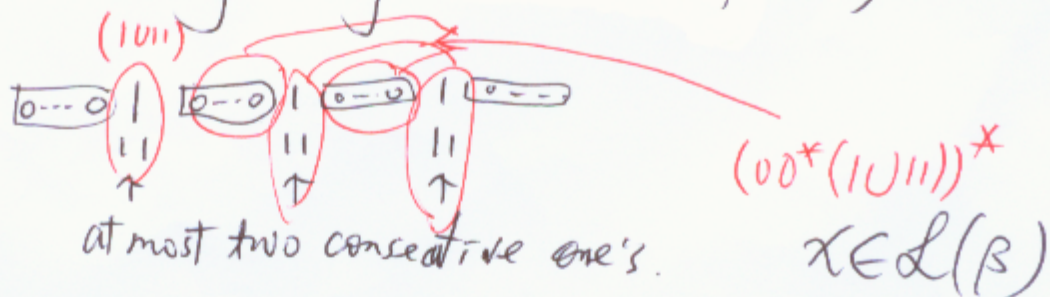
What is $\mathcal{L}(\beta)$?

$$\beta = 0^* U 0^* (1U111) (00^* (1U111))^* 0^*$$

$$\mathcal{L}(\beta) = \{ w \in \{0,1\}^* : w \text{ does not have the substring } 111 \}$$

If $x \in \mathcal{L}(\beta)$, x does not have the substring 111.

If x is a string having no occurrence of 111,



Here $\mathcal{L}(\beta) = \overline{\mathcal{L}(\alpha)}$.

$$\mathcal{L}(0^* (0^* 11^* (10^*)^*)^* 0^*) = \{0,1\}^*$$

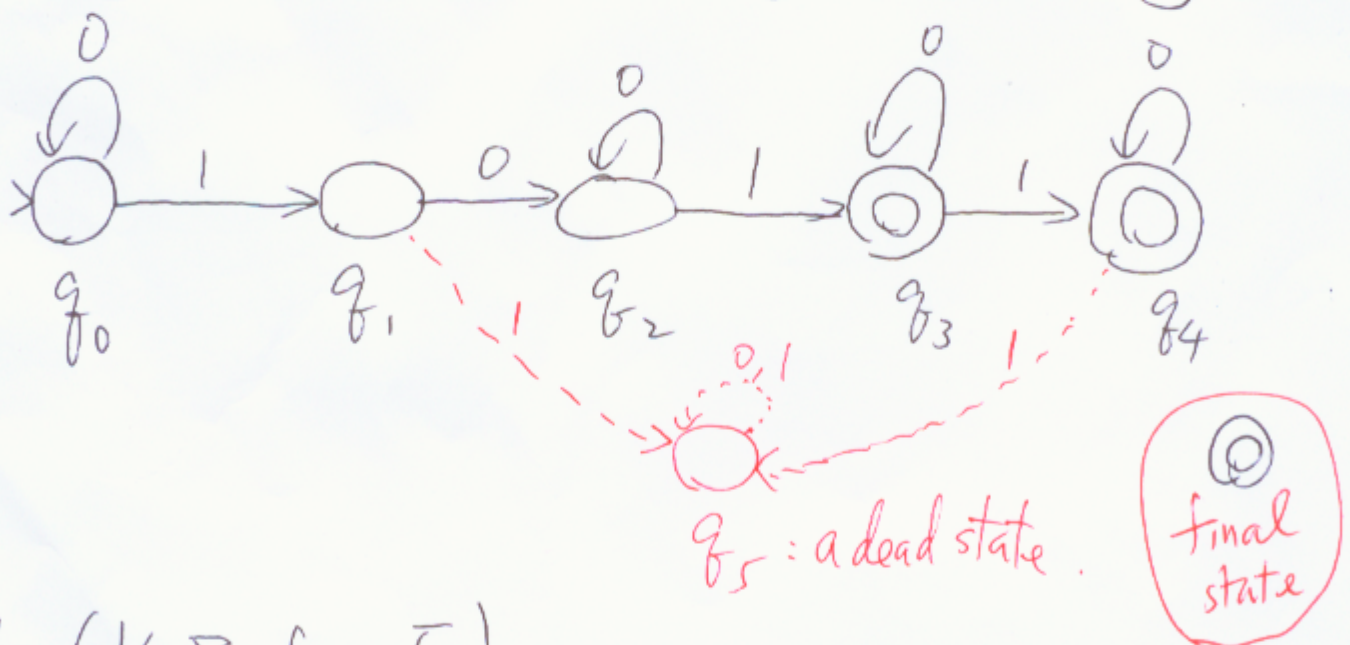
$\frac{000100110100111100001000}{0^* 10^* 10^* 10^* 10^* 10^* 0^*}$

Oct. 9, 2012

$$\alpha = 0^* 1 0 0^* 1 0^* (10^* \cup \emptyset^*)$$

$L(\alpha) = \{w \in \{0,1\}^* : w \text{ has two or three occurrences of } 1, \text{ the first and second of which are not consecutive}\}$

M : a deterministic finite automaton for recognizing $L(\alpha)$. The language recognized by M is denoted by $L(M)$.



$$M = (K, \Sigma, \delta, s, F)$$

$$K = \{q_0, q_1, q_2, q_3, q_4, q_5\}$$

$$\Sigma = \{0, 1\}$$

$$s = q_0$$

$$F = \{q_3, q_4\}$$

$$\delta: \delta(q_0, 0) = q_0$$

$$\delta(q_0, 1) = q_1$$

$$\delta(q_1, 0) = q_2$$

$$\delta(q_1, 1) = q_5$$

Is $0101100100 \in \mathcal{L}(M)$?

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Oct. 9, 2012

$(q_0, 0101100100) \vdash_M (q_0, 101100100)$

$\vdash_M (q_1, 01100100)$

$\vdash_M (q_2, 1100100)$

$\vdash_M (q_3, 100100)$

$\vdash_M (q_4, 00100)$

$\vdash_M (q_4, 0100)$

$\vdash_M (q_4, 100)$

$\vdash_M (q_5, 00)$

$\vdash_M (q_5, 0)$

$\vdash_M (q_5, \epsilon)$ not accepted

Is $0101100 \in \mathcal{L}(M)$?

$(q_0, 0101100) \vdash_M^* (q_4, \epsilon)$

accepted

Input

0 1 0 1 1 0 0 ...

tape
state:

$q_0 \ q_1 \ q_2 \ q_3 \ q_4 \ q_4 \ \dots$