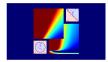
Machine Learning Foundations

(機器學習基石)



Lecture 2: Learning to Answer Yes/No

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Roadmap

1 When Can Machines Learn?

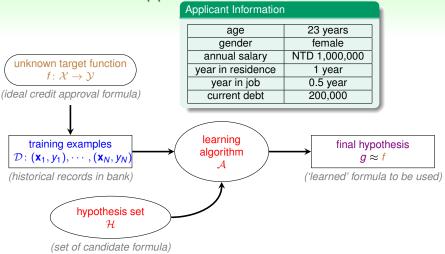
Lecture 1: The Learning Problem

 \mathcal{A} takes \mathcal{D} and \mathcal{H} to get g

Lecture 2: Learning to Answer Yes/No

- Perceptron Hypothesis Set
- Perceptron Learning Algorithm (PLA)
- Guarantee of PLA
- Non-Separable Data
- 2 Why Can Machines Learn?
- 3 How Can Machines Learn?
- 4 How Can Machines Learn Better?

Credit Approval Problem Revisited



what hypothesis set can we use?

A Simple Hypothesis Set: the 'Perceptron'

age	23 years
	NITE 4 000 000
annual salary	NTD 1,000,000
	0.5
year in job	0.5 year
aurrant dabt	200,000
current debt	200,000

• For $\mathbf{x} = (x_1, x_2, \dots, x_d)$ 'features of customer', compute a weighted 'score' and

approve credit if
$$\sum_{i=1}^{d} w_i x_i > \text{threshold}$$

deny credit if $\sum_{i=1}^{d} w_i x_i < \text{threshold}$

• \mathcal{Y} : $\{+1(good), -1(bad)\}$, 0 ignored—linear formula $h \in \mathcal{H}$ are

$$h(\mathbf{x}) = \operatorname{sign}\left(\left(\sum_{i=1}^{d} \mathbf{w}_i x_i\right) - \operatorname{threshold}\right)$$

called 'perceptron' hypothesis historically

Vector Form of Perceptron Hypothesis

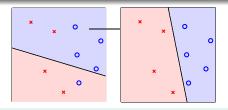
$$\begin{array}{ll} \textit{h}(\mathbf{x}) & = & \text{sign}\left(\left(\sum_{i=1}^{d} \mathbf{w}_{i} x_{i}\right) - \text{threshold}\right) \\ & = & \text{sign}\left(\left(\sum_{i=1}^{d} \mathbf{w}_{i} x_{i}\right) + \underbrace{\left(-\text{threshold}\right) \cdot \left(+1\right)}_{w_{0}}\right) \\ & = & \text{sign}\left(\sum_{i=0}^{d} \mathbf{w}_{i} x_{i}\right) \\ & = & \text{sign}\left(\mathbf{w}^{\mathsf{T}} \mathbf{x}\right) \end{array}$$

 each 'tall' w represents a hypothesis h & is multiplied with 'tall' x —will use tall versions to simplify notation

what do perceptrons h 'look like'?

Perceptrons in \mathbb{R}^2

$$h(\mathbf{x}) = \text{sign}(w_0 + w_1x_1 + w_2x_2)$$



- customer features **x**: points on the plane (or points in \mathbb{R}^d)
- labels y:
 ○ (+1), × (-1)
- hypothesis h: lines (or hyperplanes in \mathbb{R}^d)

 —positive on one side of a line, negative on the other side
- · different line classifies customers differently

perceptrons ⇔ linear (binary) classifiers

Fun Time

Consider using a perceptron to detect spam messages.

Assume that each email is represented by the frequency of keyword occurrence, and output +1 indicates a spam. Which keywords below shall have large positive weights in a **good perceptron** for the task?

- 1 coffee, tea, hamburger, steak
- free, drug, fantastic, deal
- 3 machine, learning, statistics, textbook
- 4 national, Taiwan, university, coursera

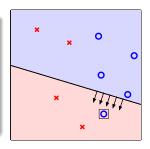
Reference Answer: (2)

The occurrence of keywords with positive weights increase the 'spam score', and hence those keywords should often appear in spams.

Select g from \mathcal{H}

 $\mathcal{H} = \text{all possible perceptrons}, g = ?$

- want: $g \approx f$ (hard when f unknown)
- almost necessary: $g \approx f$ on \mathcal{D} , ideally $g(\mathbf{x}_n) = f(\mathbf{x}_n) = y_n$
- difficult: H is of infinite size
- idea: start from some g_0 , and 'correct' its mistakes on \mathcal{D}



will represent g_0 by its weight vector \mathbf{w}_0

Perceptron Learning Algorithm

start from some \mathbf{w}_0 (say, $\mathbf{0}$), and 'correct' its mistakes on \mathcal{D}

For t = 0, 1, ...

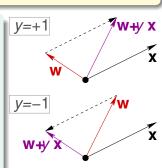
1 find a mistake of \mathbf{w}_t called $(\mathbf{x}_{n(t)}, y_{n(t)})$

$$\operatorname{sign}\left(\mathbf{w}_{t}^{\mathsf{T}}\mathbf{x}_{n(t)}\right) \neq y_{n(t)}$$

(try to) correct the mistake by

$$\mathbf{w}_{t+1} \leftarrow \mathbf{w}_t + y_{n(t)} \mathbf{x}_{n(t)}$$

... until no more mistakes return last \mathbf{w} (called \mathbf{w}_{PLA}) as g



That's it!

—A fault confessed is half redressed. :-)

Practical Implementation of PLA

start from some \mathbf{w}_0 (say, $\mathbf{0}$), and 'correct' its mistakes on \mathcal{D}

Cyclic PLA

For t = 0, 1, ...

1 find the next mistake of \mathbf{w}_t called $(\mathbf{x}_{n(t)}, y_{n(t)})$

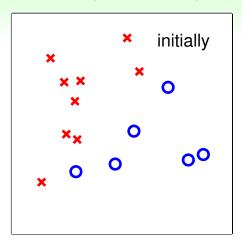
$$sign\left(\mathbf{w}_{t}^{\mathsf{T}}\mathbf{x}_{n(t)}\right) \neq y_{n(t)}$$

2 correct the mistake by

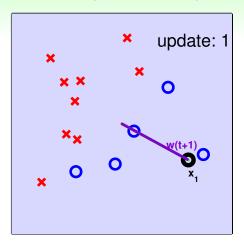
$$\mathbf{w}_{t+1} \leftarrow \mathbf{w}_t + y_{n(t)} \mathbf{x}_{n(t)}$$

... until a full cycle of not encountering mistakes

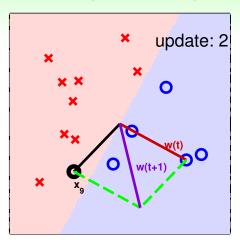
next can follow naïve cycle $(1, \dots, N)$ or precomputed random cycle



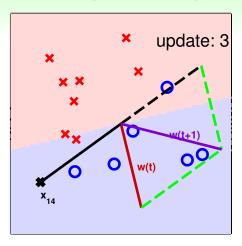
worked like a charm with < 20 lines!!



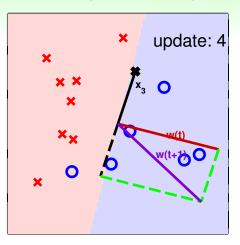
worked like a charm with < 20 lines!! (note: made $x_i \gg x_0 = 1$ for visual purpose)



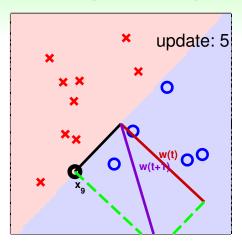
worked like a charm with < 20 lines!!



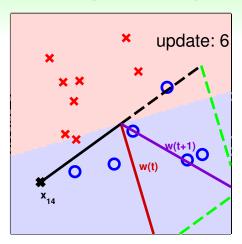
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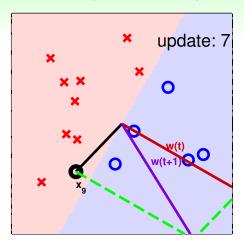
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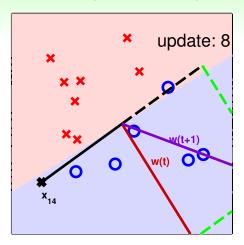
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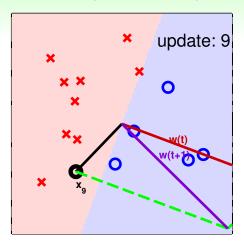
worked like a charm with < 20 lines!!



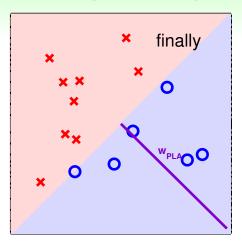
worked like a charm with < 20 lines!!



worked like a charm with < 20 lines!!



worked like a charm with < 20 lines!!



worked like a charm with < 20 lines!!

Some Remaining Issues of PLA

'correct' mistakes on \mathcal{D} until no mistakes

Algorithmic: halt (with no mistake)?

- naïve cyclic: ??
- random cyclic: ??
- other variant: ??

Learning: $g \approx f$?

- on \mathcal{D} , if halt, yes (no mistake)
- outside D: ??
- if not halting: ??

[to be shown] if (...), after 'enough' corrections, any PLA variant halts

Fun Time

Let's try to think about why PLA may work.

Let n = n(t), according to the rule of PLA below, which formula is true?

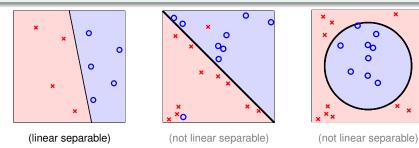
$$sign\left(\mathbf{w}_{t}^{T}\mathbf{x}_{n}\right) \neq y_{n}, \quad \mathbf{w}_{t+1} \leftarrow \mathbf{w}_{t} + y_{n}\mathbf{x}_{n}$$

Reference Answer: (3)

Simply multiply the second part of the rule by $y_n \mathbf{x}_n$. The result shows that **the rule somewhat 'tries to correct the mistake.'**

Linear Separability

- if PLA halts (i.e. no more mistakes),
 (necessary condition) D allows some w to make no mistake
- call such \mathcal{D} linear separable



assume linear separable \mathcal{D} , does PLA always halt?

PLA Fact: w_t Gets More Aligned with w_f

linear separable $\mathcal{D} \Leftrightarrow \text{exists perfect } \mathbf{w}_f \text{ such that } y_n = \text{sign}(\mathbf{w}_f^\mathsf{T} \mathbf{x}_n)$

• \mathbf{w}_f perfect hence every \mathbf{x}_n correctly away from line:

$$y_{n(t)}\mathbf{w}_{f}^{T}\mathbf{x}_{n(t)} \geq \min_{n} y_{n}\mathbf{w}_{f}^{T}\mathbf{x}_{n} > 0$$

• $\mathbf{w}_{t}^{\mathsf{T}}\mathbf{w}_{t}\uparrow$ by updating with any $(\mathbf{x}_{n(t)},y_{n(t)})$

$$\mathbf{w}_{f}^{T}\mathbf{w}_{t+1} = \mathbf{w}_{f}^{T}(\mathbf{w}_{t} + y_{n(t)}\mathbf{x}_{n(t)})$$

$$\geq \mathbf{w}_{f}^{T}\mathbf{w}_{t} + \min_{n} y_{n}\mathbf{w}_{f}^{T}\mathbf{x}_{n}$$

$$> \mathbf{w}_{f}^{T}\mathbf{w}_{t} + \mathbf{0}.$$

 \mathbf{w}_t appears more aligned with \mathbf{w}_t after update (really?)

PLA Fact: **w**_t Does Not Grow Too Fast

\mathbf{w}_t changed only when mistake

$$\Leftrightarrow$$
 sign $(\mathbf{w}_t^T \mathbf{x}_{n(t)}) \neq y_{n(t)} \Leftrightarrow y_{n(t)} \mathbf{w}_t^T \mathbf{x}_{n(t)} \leq 0$

• mistake 'limits' $\|\mathbf{w}_t\|^2$ growth, even when updating with 'longest' \mathbf{x}_n

$$\|\mathbf{w}_{t+1}\|^{2} = \|\mathbf{w}_{t} + y_{n(t)}\mathbf{x}_{n(t)}\|^{2}$$

$$= \|\mathbf{w}_{t}\|^{2} + 2y_{n(t)}\mathbf{w}_{t}^{T}\mathbf{x}_{n(t)} + \|y_{n(t)}\mathbf{x}_{n(t)}\|^{2}$$

$$\leq \|\mathbf{w}_{t}\|^{2} + 0 + \|y_{n(t)}\mathbf{x}_{n(t)}\|^{2}$$

$$\leq \|\mathbf{w}_{t}\|^{2} + \max_{n} \|y_{n}\mathbf{x}_{n}\|^{2}$$

start from $\mathbf{w}_0 = \mathbf{0}$, after T mistake corrections,

$$\frac{\mathbf{w}_{\mathit{f}}^{\mathit{T}}}{\|\mathbf{w}_{\mathit{f}}\|}\frac{\mathbf{w}_{\mathit{T}}}{\|\mathbf{w}_{\mathit{T}}\|} \geq \sqrt{\mathit{T}} \cdot \mathsf{constant}$$

Fun Time

Let's upper-bound T, the number of mistakes that PLA 'corrects'.

Define
$$R^2 = \max_n \|\mathbf{x}_n\|^2$$
 $\rho = \min_n y_n \frac{\mathbf{w}_f^T}{\|\mathbf{w}_f\|} \mathbf{x}_n$

We want to show that $T \leq \square$. Express the upper bound \square by the two terms above.

- $\mathbf{0} R/\rho$
- **2** R^2/ρ^2
- $3 R/\rho^2$
- Φ^{2}/R^{2}

Reference Answer: (2)

The maximum value of $\frac{\mathbf{w}_t^T}{\|\mathbf{w}_t\|} \frac{\mathbf{w}_t}{\|\mathbf{w}_t\|}$ is 1. Since T mistake corrections **increase the inner product by** $\sqrt{T} \cdot$ **constant**, the maximum number of corrected mistakes is $1/\text{constant}^2$.

More about PLA

Guarantee

as long as linear separable and correct by mistake

- inner product of \mathbf{w}_t and \mathbf{w}_t grows fast; length of \mathbf{w}_t grows slowly
- PLA 'lines' are more and more aligned with w_f ⇒ halts

Pros

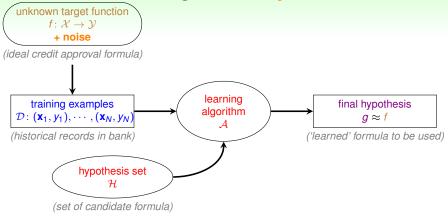
simple to implement, fast, works in any dimension d

Cons

- 'assumes' linear separable $\mathcal D$ to halt
 - —property unknown in advance (no need for PLA if we know \mathbf{w}_f)
- not fully sure how long halting takes (ρ depends on w_f)
 —though practically fast

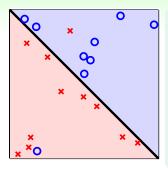
what if \mathcal{D} not linear separable?

Learning with Noisy Data



how to at least get $g \approx f$ on noisy \mathcal{D} ?

Line with Noise Tolerance



- assume 'little' noise: $y_n = f(\mathbf{x}_n)$ usually
- if so, $g \approx f$ on $\mathcal{D} \Leftrightarrow y_n = g(\mathbf{x}_n)$ usually
- how about

$$\mathbf{w}_g \leftarrow \underset{\mathbf{w}}{\operatorname{argmin}} \sum_{n=1}^{N} \left[y_n \neq \operatorname{sign}(\mathbf{w}^T \mathbf{x}_n) \right]$$

—NP-hard to solve, unfortunately

can we modify PLA to get an 'approximately good' *g*?

Pocket Algorithm

modify PLA algorithm (black lines) by keeping best weights in pocket

initialize pocket weights ŵ

For $t = 0, 1, \cdots$

- 1 find a (random) mistake of \mathbf{w}_t called $(\mathbf{x}_{n(t)}, y_{n(t)})$
- (try to) correct the mistake by

$$\mathbf{w}_{t+1} \leftarrow \mathbf{w}_t + y_{n(t)} \mathbf{x}_{n(t)}$$

3 if \mathbf{w}_{t+1} makes fewer mistakes than $\hat{\mathbf{w}}$, replace $\hat{\mathbf{w}}$ by \mathbf{w}_{t+1}

...until enough iterations return $\hat{\mathbf{w}}$ (called \mathbf{w}_{POCKET}) as g

a simple modification of PLA to find (somewhat) 'best' weights

Fun Time

Should we use pocket or PLA?

Since we do not know whether \mathcal{D} is linear separable in advance, we may decide to just go with pocket instead of PLA. If \mathcal{D} is actually linear separable, what's the difference between the two?

- $oldsymbol{0}$ pocket on $\mathcal D$ is slower than PLA
- 2 pocket on \mathcal{D} is faster than PLA
- \odot pocket on \mathcal{D} returns a better g in approximating f than PLA
- 4 pocket on \mathcal{D} returns a worse g in approximating f than PLA

Reference Answer: 1

Because pocket need to check whether \mathbf{w}_{t+1} is better than $\hat{\mathbf{w}}$ in each iteration, it is slower than PLA. On linear separable \mathcal{D} , $\mathbf{w}_{\text{POCKET}}$ is the same as \mathbf{w}_{PLA} , both making no mistakes.

Summary

1 When Can Machines Learn?

Lecture 1: The Learning Problem

Lecture 2: Learning to Answer Yes/No

- Perceptron Hypothesis Set
 hyperplanes/linear classifiers in R^d
- Perceptron Learning Algorithm (PLA)
 correct mistakes and improve iteratively
- Guarantee of PLA
 no mistake eventually if linear separable
- Non-Separable Data
 hold somewhat 'best' weights in pocket
- next: the zoo of learning problems
- 2 Why Can Machines Learn?
- 3 How Can Machines Learn?
- 4 How Can Machines Learn Better?