

Finiteness of BinarySearch

1 Question

The following BINARYSEARCH algorithm has been introduced in the class. Prove that the algorithm satisfies the *finiteness* property. That is, the while only runs for a finite number of iterations for an ordered array *arr*.

```

BINARYSEARCH(integer array arr, integer len, integer key)
left ← 0, right ← len - 1
while left ≤ right do
  mid ← floor((left + right)/2)
  if key = arr[mid] then
    return mid
  else if key < arr[mid] then
    right ← mid - 1
  else
    left ← mid + 1
  end if
end while
return NOTFOUND

```

2 Answer

Claim 0: In every iteration, either $key = mid$ and the algorithm returns, or the range $right - left + 1$ strictly decreases.

Proof: If the algorithm does not return in the iteration, then either $right \leftarrow mid - 1$ or $left \leftarrow mid + 1$. Let (l_0, r_0) denote the old pair of *left* and *right*, and (l_1, r_1) denote the new pair. Then, in the former case of $right \leftarrow mid - 1$,

$$\begin{aligned}
 r_1 - l_1 + 1 &= \left\lfloor \frac{l_0 + r_0}{2} \right\rfloor - 1 - l_0 + 1 \\
 &\leq \frac{l_0 + r_0}{2} - l_0 \\
 &= \frac{r_0 - l_0}{2} \\
 &= r_0 - l_0 + 1 - \left(\frac{r_0 - l_0}{2} + 1\right)
 \end{aligned}$$

Because $r_0 - l_0 \geq 0$ (the condition of WHILE), the term $\frac{r_0 - l_0}{2} + 1$ is strictly positive. Thus, $r_1 - l_1 + 1$ strictly decreases.

Claim 1: The algorithm runs for a finite number of iterations.

Proof: The range $right - left + 1$ needs to be a positive integer because the WHILE loop needs $left \leq right$. Initially, $right - left + 1$ is simply *len*. From **Claim 0**, the range cannot remain positive after at most *len* iterations. Thus, the algorithm runs for a finite number of iterations.