## Theory of Computation

## Homework 4

## Due: 2015/12/08

**Problem 1.** Find all the primitive roots of 5 and all the primitive roots of 7.

Solution.

The primitive roots of 5 are 2 and 3, because  $\phi(5) = 4$  and

$$2^1 \equiv 2 \pmod{5}$$
;  $2^2 \equiv 4 \pmod{5}$ ,  
 $2^3 \equiv 3 \pmod{5}$ ;  $2^4 \equiv 1 \pmod{5}$ ,

and

$$3^1 \equiv 3 \pmod{5}$$
;  $3^2 \equiv 4 \pmod{5}$ ,  
 $3^3 \equiv 2 \pmod{5}$ ;  $3^4 \equiv 1 \pmod{5}$ .

Similarly, the primitive roots of 7 are 3 and 5 because  $\phi(7) = 6$  and

$$3^{1} \equiv 3 \pmod{7} ; \quad 3^{2} \equiv 2 \pmod{7},$$
  

$$3^{3} \equiv 6 \pmod{7} ; \quad 3^{4} \equiv 4 \pmod{7},$$
  

$$3^{5} \equiv 5 \pmod{7} ; \quad 3^{6} \equiv 1 \pmod{7},$$

and

$$5^1 \equiv 5 \pmod{7}$$
;  $5^2 \equiv 4 \pmod{7}$ ,  
 $5^3 \equiv 6 \pmod{7}$ ;  $5^4 \equiv 2 \pmod{7}$ ,  
 $5^5 \equiv 3 \pmod{7}$ ;  $5^6 \equiv 1 \pmod{7}$ .

**Problem 2.** We know that 3-SAT is NP-complete. Show that for n > 3, n-SAT is also NP-complete. (You don't need to show that is in NP.)

## Solution.

We reduce 3-SAT to *n*-SAT as follows. Let  $\phi$  be a 3-SAT boolean expression. For any clause  $(a \lor b \lor c)$ , we replace it with  $(a \lor b \lor \underline{c} \lor \cdots \lor \underline{c})$ . By repeating this process in all the clauses of  $\phi$ , we get an *n*-SAT boolean expression  $\phi'$ . Now, we proceed to show that this is a reduction from 3-SAT to *n*-SAT as follows:

- (⇒) From the construction, we see that if a truth assignment satisfies  $\phi$ , then it must satisfy  $\phi'$ .
- ( $\Leftarrow$ ) Let's note that if a truth assignment satisfy  $\phi'$ , then it must also satisfy  $\phi$ .

From this, we then deduce that  $\phi$  is satisfiable if and only if  $\phi'$  is satisfiable; hence 3-SAT is reducible to *n*-SAT, proving that *n*-SAT is NP-complete.  $\Box$